Asynchronous Many-Task systems for Exascale 2023

# Malleable APGAS Programs and their support in Batch Job Scheduler

#### <u>Patrick Finnerty<sup>1</sup></u>, Leo Takaoka<sup>1</sup>, Takuma Kanzaki<sup>1</sup>, Jonas Posner<sup>2</sup>

Kobe University, Kobe, Japan
 University of Kassel, Kassel, Germany

#### Parallel Job Scheduling



Jobs submitted by users

Possible schedule on parallel machines

## Malleable Job Potential

- Malleable Job?
  - → can change nb of nodes used (grow/shrink) during execution following scheduler instruction



- Benefits?
  - $\rightarrow$  More effective use of resources
  - $\rightarrow$  Increased throughput and/or reduced energy consumption

#### Hurdles

- Support in RMS (Slurm, Torque, Open PBS, ...)
  - Interactions between program/RMS necessary
  - Some experimental implementations
- Support in libraries and Programming models
  - MPI w/ Checkpoint/Restart
  - Charm++
  - (A)PGAS languages
- Compatible programs
  - Requires effort to convert/create applications

Our work

#### Contributions

- Malleable implementation of APGAS for Java
  - Simple abstractions for programmers
  - Modularity for adaptation to future schedulers
- Refactoring of an AMT work-stealing scheme
  - Simplified through our abstractions
  - Can respond to scheduler directives
- Demonstrate the benefits of malleability on a Beowulf cluster with a custom scheduler

# A word about (A)PGAS

- Partitioned Global Address Space
  - The memory is partitioned between processes
  - That memory can be accessed from remote processes
- Asynchronous PGAS
  - Can spawn and control termination of tasks on the processes
- In X10 and APGAS for Java:
  - "Place," finish at async



#### Malleable APGAS?

Add and remove Places!



#### Malleable APGAS architecture



## Application example

- Lifeline-based Global Load Balancer
  - Each place has some tasks
  - When a place runs out of tasks, steal through preferential channels, the "lifelines"



# Application example

- Grow order
  - Before: do nothing
  - After: integrate the new place(s) into the lifeline network, they steal some work and start working



Note:

Computation continues while new places are added to the runtime!

# Application example

- Shrink order
  - Before: Disconnect place(s) from lifeline network, relocate any work to remaining places
  - After: do nothing



Note:

Computation continues while places are removed from the runtime!

#### **Evaluation - environment**

- Beowulf cluster
  - 12 nodes
  - Password-less
    SSH authentication
- Simplistic malleable job scheduler
  - First-Come First-Served
  - If possible, shrink jobs to allow the next one to start
  - Otherwise, grow running malleable jobs
- 30-job batch
  - Some MPI, some APGAS
  - Malleable programs have a min/max nb of nodes
  - Compare "100% fixed" with "half malleable/half fixed" batch

#### Evaluation - results



Workload:	100% rigid	Half malleable half rigid	Higher
Makespan (m)	21.4	18.2 (-15%)	throughput!
Avg cluster utilization	72.3%	83.4% (+15%)	Malleable jobs use fewer nodes for longer
Avg wait time (m)	9.4	7.9 (-15%)	
Avg exec time (m)	0.96	1.40 (+45%)	
Avg response time (m)	10.36	9.34 (-10%)	

#### Conclusion

- Presented a practical, extensible, malleable implementation of APGAS for Java
- Demonstrated performance benefits on our small Beowulf, even with a simplistic scheduling strategy
- Perspectives
  - Support for evolving jobs? Other applications?
    - Change in resource allotment initiated by the program
  - Evaluation on larger clusters?
    - Possible within a certain degree ...
    - Scheduling algorithms best studied through simulation

#### Conclusion

- All of our source code is freely available:
  - Malleable APGAS runtime
    <u>https://github.com/projectwagomu/apgas</u>
  - Malleable lifeline-based Global Load Balancer
    <a href="https://github.com/projectwagomu/lifelineglb">https://github.com/projectwagomu/lifelineglb</a>
  - Custom scheduler <u>https://github.com/projectwagomu/FIFO-Malleable-Job-Scheduler</u>

# Thank you for listening!