# All centralising monoids on the set {0*,* 1*,* 2}, including their witnesses

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This data set provides supplementary material for the description of witnesses for all 192 centralising monoids [6, 7] on the three-element set  $A = \{0, 1, 2\}$  as carried out in [1], cp. also [8].

Centralising monoids are the closed sets on one side of the Galois connection between finitary functions in  $\mathcal{O}_A = \bigcup_{n \in \mathbb{N}_+} A^{A^n}$  and  $\mathcal{O}_A^{(1)} = A^A$  induced by the relation of *commutation*: we say that  $f: A^n \to A$  for some  $n \in \mathbb{N}_+$  commutes with some unary function  $s: A \to A$  if and only if  $s: \langle A; f \rangle \to \langle A; f \rangle$  is an endomorphism of the algebra  $\langle A; f \rangle$ , i.e., if  $s \in$  End  $(\langle A; f \rangle)$ , meaning in more detail that  $s(f(x_1, ..., x_n)) = f(s(x_1), ..., s(x_n))$  for all  $x_1, ..., x_n \in A$ . We denote this by  $f \perp s$  or by  $f \in \{s\}^*$  or  $s \in \{f\}^*$  (more precisely  $s \in \{f\}^{*(1)}$ ). A centralising monoid is a set of the form  $F^{*(1)} = \left\{ s \in \mathcal{O}_A^{(1)} \mid \forall f \in F : f \perp s \right\}$ for some  $F \subseteq \mathcal{O}_A$ . Note that  $F^{*(1)} = \bigcap_{f \in F} {\{f\}}^{*(1)}$ . Any set  $F \subseteq \mathcal{O}_A$  describing the monoid  $M = F^{*(1)}$  is called a *witness* of the centralising monoid; witnesses are not unique.

Machida and Rosenberg in [6] presented a strategy to determine all centralising monoids on  $A = \{0, 1, 2\}$  and included a list of the resulting 192 monoids in [7], which is not easily, for example, by a computer, verifiable since there are no witnesses given. In [8, 1], we rephrased the description of all centralising monoids using formal concept analysis  $[4, 3]$ . Namely, the set of centralising monoids can be seen as the set of intents of the formal context  $\mathbb{K}_1 = (\mathcal{O}_A, \mathcal{O}_A^{(1)}, \perp).$ In [8, 1] we demonstrate how to find a finite subcontext  $\mathbb{K}'$  of  $\mathbb{K}_1$  having only 183 objects, but the same set of intents as  $K_1$ . The subcontext  $K'$  can be further reduced to a subcontext  $K''$  with only 51 objects, which is still able to describe all centralising monoids on  $\{0, 1, 2\}$ . The advantage of this approach is twofold: we may apply well-known algorithms (cf. [3, Chapter 2] and [4, Chapter 2], see also the implementations [2, 5]) to automatically enumerate all concepts (and thus in particular all intents) of K′ ; furthermore, the corresponding extents of the concepts provide witnesses that were missing in [7]. In this way we are able to confirm the correctness of the 192 monoids derived by Machida and Rosenberg

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and of their 48 conjugacy types; furthermore, our list of witnesses allows for a quick verification that the 192 monoids we present are indeed centralising monoids. A more detailed discussion of our method can be found in [8] and [1]. Here we only comment on certain aspects that are needed to make sense of the files in this data set (in conjunction with consulting [1] or [8] for the mathematical background); moreover we show how our data shows the correctness of the results of [6, 7].

## **1 Formal contexts**

### **1.1 Object clarification**

The context  $\mathbb{K}_1$  contains many duplicate rows, that is, functions  $f_1 \neq f_2 \in \mathcal{O}_A$ with  ${f_1}^{(1)}={f_2}^{(1)}$ , and it is clearly unnecessary to store duplicate rows in order to give the lattice of intents of  $K_1$ . Keeping only one representative for each duplicate row is called *object clarification* in the language of formal concept analysis. As  $K_1$  contains an infinite number of rows, we have independently written two programmes that construct an object clarified subcontext  $\mathbb{K}'$  =  $(W_A, \mathcal{O}_A^{(1)}, \perp)$  of  $\mathbb{K}_1$ , where  $|W_A| = 183$  (files comm-ternary-unary-3.cxt<sup>1</sup> and context\_of\_commutation.cxt<sup>2</sup>). The contexts  $\mathbb{K}'$  produced by the two programmes contain different sets of witnesses, as the codes differ in the order in which they enumerate certain functions as objects and thus the results of the object clarification process are distinct. However, their object intents (and thus the set of all intents, that is, of centralising monoids) is identical in both cases.

## **1.2 Encoding of operations used as objects or attributes**

In order to understand the objects and attributes in the two mentioned context files representing variants of  $\mathbb{K}'$ , we have to define the coding of operations used in these files. The operations appearing in the files are operations of arity at most three, that is, functions of the form  $f: A^n \to A$  where  $A = \{0, 1, 2\}$ and  $n \in \mathbb{N}, 1 \leq n \leq 3$ . In the context files, these functions are represented by finite strings of digits from  $A = \{0, 1, 2\}$  that stand for the value tables of the operations. In order to identify a function with such a linearised value table, we have to specify the order of the arguments used in the string of digits. We represent functions  $f: A \to A$  as the string  $(f(0), f(1), f(2)) \in A^3$  (in the context files, no commas or parentheses are used since each letter is a single symbol, for example, the identity operation  $id_3$  is represented by the string 012). For binary operations  $f: A^2 \to A$  there are two natural representations, *row-wise* and *column-wise* coding. Row-wise coding means that *f* is stored as the string  $(f(0,0), f(0,1), f(0,2), f(1,0), f(1,1), f(1,2), f(2,0), f(2,1), f(2,2)) \in A^9$ , in other words, in row-wise coding a string  $(a_0, \ldots, a_8) \in A^9$  represents the function

<sup>&</sup>lt;sup>1</sup>We have also added the  $c++$ -file commutation.cpp with which this data was produced from the command compute\_ternary\_unary("comm-ternary-unary-3.cxt");. Let us note that the same programme is also capable of outputting the contexts  $(\mathcal{O}_A^{(1)}, \mathcal{O}_A^{(1)}, \perp)$  via  $\texttt{write\_context\_unary\_unary('comm-unary-unary-3. cxt"); and of \big(}\mathcal{O}^{(2)}_A, \mathcal{O}^{(1)}_A, \bot\big) \text{ via write\_}$ context\_binary\_unary("comm-bin-unary-3.cxt"); if these commands are uncommented.

<sup>&</sup>lt;sup>2</sup>This file has been produced by the c++-file context\_of\_commutation\_code.cpp, written by Leon Renkin, cf. [8, Section 7.1].

with value table

$$
\begin{array}{c|cc}\nx \lor y & 0 & 1 & 2 \\
\hline\n0 & a_0 & a_1 & a_2 \\
1 & a_3 & a_4 & a_5 \\
2 & a_6 & a_7 & a_8\n\end{array}
$$
, where  $f(x, y)$  appears in row x, column y.

On the other hand, column-wise coding means that  $f: A^2 \to A$  is stored as the string  $(f(0,0), f(1,0), f(2,0), f(0,1), f(1,1), f(2,1), f(0,2), f(1,2), f(2,2)) \in A^9$ in other words, in column-wise coding a string  $(a_0, \ldots, a_8) \in A^9$  represents the function with the value table

$$
\begin{array}{c|cc}\nx \backslash y & 0 & 1 & 2 \\
\hline\n0 & a_0 & a_3 & a_6 \\
1 & a_1 & a_4 & a_7 \\
2 & a_2 & a_5 & a_8\n\end{array}
$$
, where  $f(x, y)$  appears in row x, column y.

If one were to accidentally confuse the two encodings, i.e., to construct the opposite table out of a row-wise encoded function by mistaking it as a columnwise coding, not much harm would arise, as both function have identical properties with respect to commutation, that is, they commute with the same functions and hence define the same centraliser.

For ternary operations  $f: A^3 \to A$  there are again two major ways of encoding them, coming from listing the arguments lexicographically such that either left-most letters take priority or right-most letters take priority. For brevity, we will call these possibilities again row-wise and column-wise coding. That is, in row-wise coding, a function  $f: A^3 \to A$  is represented as  $(f(0,0,0), f(0,0,1), f(0,0,2), f(0,1,0), f(0,1,1), f(0,1,2), \ldots, f(2,2,2)) \in A^{26}$ i.e.,  $f(x, y, z)$  appears at the index  $3^2x+3y+z$  of the string, while in column-wise coding,  $(f(0,0,0), f(1,0,0), f(2,0,0), f(0,1,0), f(1,1,0), f(2,1,0), \ldots, f(2,2,2))$ in  $A^{26}$  is used to represent *f*, i.e.,  $f(x, y, z)$  appears at the index  $x + 3y + 3^2z$  of the string.

In the programme commutation.cpp and the files comm-unary-unary-3.cxt, comm-bin-unary-3.cxt, comm-ternary-unary-3.cxt produced by it, always row-wise encoding is used. In the programme context\_of\_commutation\_code. cpp, which produces the file context\_of\_commutation.cxt, column-wise encoding is used for the binary and ternary operations appearing as objects.

For a more concise representation of the closures, the objects and attributes of the context context of commutation.cxt have been consecutively numbered in the form  $h_0, \ldots, h_{155}, f_{156}, \ldots, f_{182}$  (for the objects) and  $s_0 = 000, \ldots, s_{26} = 222$ (for the attributes). The precise correspondence between the string representations of the functions and the newly introduced identifiers is given in [8, Section 7.2], the corresponding context with renamed objects and attributes is shown in [8, Section 7.3] and has been added to this data set as file context\_ of\_commutation\_names.cxt.

### **1.3 Object reduction**

Another simplification method that still keeps the set of intents intact is *object reduction* [3, Section 1.4.3, p. 27 et seqq.], also called *row reduction*. In an object reduced subcontext  $\mathbb{K}''$  of  $\mathbb{K}'$  only 51 objects from the 183 functions in  $W_A$  remain; a representation of  $K''$  is given in [1, Table 1], cf. context\_ of\_commutation\_names\_obj\_red.cxt. Applying also the dual operation to the columns, one obtains a fully reduced context  $\mathbb{K}^{\prime\prime\prime}$  from  $\mathbb{K}^{\prime}$  that has 51 objects (rows) and 25 of the 27 attributes (columns) of  $\mathbb{K}_1$  or  $\mathbb{K}'$ . The context  $\mathbb{K}'''$  has an isomorphic lattice of intents, but obviously not the same set of intents as  $K'$  or  $K$ ; it is the standard context belonging to  $\mathbb{K}_1$  and is presented in [8, Section 7.4], cf. context of commutation names red.cxt. The reduction process has been carried out with readily available formal concept analysis software, such as CONEXP<sup>3</sup> or CONEXP-CLJ (cf.  $[2, 5]$ ). Depending on which file representation of  $K'$  we start with, we obtain different object reduced  $(K'')$  or fully reduced  $(K''')$  context files:



For further computation of lists of concepts, intents (cf. [8, Section 7.5]), and intents up to conjugacy (cf. [1]) we have proceeded from the context  $\mathbb{K}'$  as given in context of commutation.cxt or from  $K''$  as given in context of commutation\_names\_obj\_red.cxt.

#### **1.4 Context file format**

All formal contexts in this data set are given in Burmeister format, which can be read by standard formal concept analysis software like ConExp or Conexp-clj [2] and complies with the following specification:

- A line with the letter B (for Burmeister).
- A blank line
- A line with the integer number *m* of objects.
- A line with the integer number *n* of attributes.
- A blank line
- *m* lines with labels for the objects.
- *n* lines with labels for the attributes.
- *m* lines representing the context, that is, the incidence relation between objects and attributes. The *i*-th of these lines corresponds to the *i*-th object  $q_i$  and contains a string of *n* characters **X** or ., **X** standing for  $q_i$ being in relation with the corresponding attribute and . for the opposite. This is an encoding of the characteristic tuple of the object intent given by  $g_i$  where **X** represents 1 and . represents 0.

<sup>3</sup>http://sourceforge.net/projects/conexp



3 2 f g h P Q X.

B

.X XX

*It represents the Galois connection between* {*f, g, h*} *and* {*P, Q*}*, where f has property P but not Q, g has Q but not P, and h has both. As a binary incidence relation*  $I \subseteq \{f, g, h\} \times \{P, Q\}$  *this would be represented as*  $I = \{(f, P), (g, Q), (h, P), (h, Q)\},$  which commonly is summarised in a tabu*lar form (note the resemblance to the last three lines of the listing) as*

$$
\begin{array}{c|cc}\n & P & Q \\
\hline\nf & \times & \\
g & \times & \\
h & \times & \times\n\end{array}
$$

# **2 Centralising monoids**

One of the main goals of  $[8, 1]$  was to derive a finite context  $K'$ , the intents of which are all centralising monoids on  $\{0, 1, 2\}$ . This has been achieved using the programme context\_of\_commutation\_code.cpp in the files context\_of\_ commutation.cxt and context\_of\_commutation\_names\_obj\_red.cxt. The intents of these contexts are the centralising monoids, the corresponding extents may serve as witnesses. Concepts and intents were computed using [2], see also [5]. The number of concepts/intents computed agrees with the figure 192 reported in [6, 7]. In Section 5 we shall describe that we were also able to confirm the correctness of the lists of 192 monoids given in [7].

The following files were obtained from  $K'$  as represented in context\_of\_ commutation.cxt with the help of  $[2]^4$  and some manual change of the formatting (line breaks, separating symbols, etc.): concepts.txt, the 192 concepts of  $K'$ , i.e., the centralising monoids together with witnesses; intents.txt, the 192 intents of K′ ; intents\_numbers.cppinput, the 192 intents of K′ to be used as an input for the programme reduce\_monoids\_conj.cpp (here the attribute labels abc for unary operations  $s \in \mathcal{O}_A^{(1)}$ , where  $s(0) = a$ ,  $s(1) = b$ ,  $s(2) = c$ , have been replaced by an integer in  $\{0, \ldots, 26\}$  computed as  $3^2 \cdot s(0) + 3 \cdot s(1) + s(2) = 9a + 3b + c$ , and line endings have been marked by the value 30; intents\_numbers.cppinput was obtained from intents.txt with the help of the GNU sed command (version 4.2.2) via sed -f int2numb.sed intents.txt > intents\_numbers.cppinput).

 $^4\rm{Use}$  the commands (def cxt (read-context "path\_to/context\_of\_commutation.cxt")) followed by (concepts cxt) or (intents cxt).

For the presentation of intents (monoids) and witnesses in [1, 8], a different notation (labelling) of the objects and attributes was needed. As explained in Section 1.2 on encoding of operations in the formal contexts, the file context\_ of\_commutation.cxt was changed to context\_of\_commutation\_names.cxt for this purpose, cf. also [8, Sections 7.2, 7.3] for more details. With the help of ConExp/Conexp-clj this file was object reduced, resulting in the context file context of commutation names obj red.cxt representing  $K''$ . Using [2] and some manual formatting as above, the concepts (monoids and their witnesses) of context\_of\_commutation\_names\_obj\_red.cxt were computed and stored in the file concepts\_new\_format.txt (cf. [8, Section 7.5]), and the corresponding intents were stored in intents\_new\_format.txt.

## **3 Centralising monoids up to conjugacy**

For [1, Table 2] we also computed representatives of the 192 centralising monoids on  $A = \{0, 1, 2\}$  up to conjugacy by inner automorphisms  $s \in \text{Sym}(A)$ . In this way we were able to confirm the number 48 of conjugacy types mentioned in [6, 7]. We were also able to verify the correctness of the 48 conjugacy representatives shown in [6, Table 3]; for more details see Section 4.

To clarify the terminology, we define the *conjugate* of an *n*-ary operation  $f: A^n \to A$  by a permutation  $s \in \text{Sym}(A)$  to be the function  $f^s: A^n \to A$  that is given by  $f^{s}(x_1, ..., x_n) := s(f(s^{-1}(x_1), ..., s^{-1}(x_n)))$  for all  $(x_1, ..., x_n) \in A^n$ . Functions  $f, g \in \mathcal{O}_{A}^{(n)}$  are said to be conjugate if there is  $s \in \text{Sym}(A)$  such that  $f = g^s$ , and this is clearly an equivalence relation. We extend the notion of conjugacy element-wise to sets  $F \subseteq \mathcal{O}_A^{(n)}$  by putting  $F^s := \{f^s \mid f \in F\}$  for  $s \in \text{Sym}(A)$ . In fact, we use this definition here only for the case where  $n = 1$  and  $F \subseteq \mathcal{O}_A^{(1)}$ . In this way we can speak of the *conjugate*  $F^s$  by *s* of a transformation monoid  $F \subseteq \mathcal{O}_A^{(1)}$ .

The classification of all centralising monoids on {0*,* 1*,* 2} up to conjugacy was done as follows: the programme reduce monoids conj.cpp reads all 192 centralising monoids on  $A = \{0, 1, 2\}$  as lists of integers in  $\{0, \ldots, 26\}$  from the file intents\_numbers.cppinput and outputs one representative monoid of each conjugacy class to the file intents\_up\_to\_conj.txt. In the programme  $\texttt{reduce\_monoids\_conj}.\texttt{cpp}$  (and its in- and output) unary operations  $s \in \mathcal{O}_A^{(1)}$ are represented by their hash value  $9 \cdot s(0) + 3 \cdot s(1) + s(2) \in \{0, \ldots, 26\}$ . The code is rather straightforward: each monoid (line) of intents\_numbers.cppinput is read in as a set of integers, these 192 sets are stored as a set. One then in compute monoids up to conj loops over each monoid  $F$  in that set and every  $s \in \text{Sym}(A)$  and conjugates all the elements of the monoid *F* by *s*, obtaining  $F^s$ . If the resulting isomorphic monoid  $F^s$  is distinct from the current monoid  $F$  (that is, it is a distinct isomorphic copy in the same conjugacy class), it is removed from the set of monoids, and thus only the representatives up to conjugacy remain. These are then output likewise to intents\_up\_to\_conj.txt; this file has 48 lines.

For the presentation of the representative monoids of the 48 conjugacy types in [1, Table 2], it was necessary to find these 48 intents among the 192 concepts of concepts\_new\_format.txt. For this, the 192 concepts from concepts\_new\_ format.txt were split at the ; into their extent and intent part, and some

delimiting characters such as  $(, \,), \, \{, \}$ ,  $\{$  and  $\}$ , as well as the prefixes s for the attributes of the intents, were removed. This resulted in the files 192\_extents.txt and 192\_intents.txt.

To identify the 48 conjugacy representatives from intents\_up\_to\_conj.txt among the concepts of concepts\_new\_format.txt, that is, among the extents and intents of 192 extents.txt and 192 intents.txt, respectively, the python script finding 48 conjugacy types among 192 intents extents. py was written. The lines of 192\_extents.txt form the entries (sets) of the list extentstrlist, the same lines without the h and f prefixes are the entries of the list extentlist (removing the prefixes is sufficient to identify the objects of the extents uniquely). The lines of 192\_intents.txt are the entries of the list intentlist. The lines of intents\_up\_to\_conj.txt are the entries of the list listoftypes. The procedure print\_sorted\_intents\_extents\_of\_types then searches for the occurrence of each monoid in listoftypes in the list intentlist and outputs the corresponding index in the list, the intent (monoid) and the corresponding extent (witness). The result of running the script finding\_48\_conjugacy\_types\_among\_192\_intents\_extents.py is stored in the file conjugacy\_types\_intents\_extents.txt. From this file the lines containing 'Ext:' have been extracted to the file 48 extents.txt, and the lines containing 'Int:' have been extracted to the file 48\_intents.txt. Based on those files, [1, Table 2] has been constructed. Moreover, from the output file conjugacy\_types\_intents\_extents.txt we also extracted the list 192\_extents\_sorted.txt of the 192 extents corresponding line-wise to the ones in 192\_extents.txt, but where the objects are mentioned in ascending order of their indices. Similarly, we extracted the list 192\_intents\_sorted.txt of the 192 intents (monoids) corresponding line-wise to the ones in 192\_intents. txt, but again the functions in the monoids have been listed in ascending order of their hash values.

# **4 Comparison with results from [6]**

We went to compare our results (in this data set and those mentioned in [1]) with the data available from [6]. The figures 192 and 48 reported in [6, Proposition 3.5] are already confirmed by the numbers of intents in intents.txt and intents\_up\_to\_conj.txt, respectively. The other central piece of data is Table 3 from [6, p. 280]. We captured this table in MachidaRosenberg2012 Table3.txt, and, in order to better compare it with our results, reordered the columns of this table such that the operations  $s \in \mathcal{O}_A^{(1)}$  are listed in ascending order of their hash values  $9 \cdot s(0) + 3 \cdot s(1) + s(2)$  in  $\{0, \ldots, 26\}$ . That is, after reordering, the first column corresponds to the operation  $s_0$  and the last one to  $s_{26}$ as in context\_of\_commutation.cxt and context\_of\_commutation\_names. cxt. The reordered table can be found in the file MachidaRosenberg2012\_ Table3\_ordered\_0-26.txt, and it is stored as MR2012tbl\_reordered in the file finding\_intents\_up\_to\_conj\_as\_conjugates\_of\_MR2012\_Table3. py. We used the function print\_converted\_list\_to\_integers(MR2012tbl\_ reordered) to print the 48 monoids from [6, Table 3] as 48 lines of integer hash values in  $\{0, \ldots, 26\}$ , keeping the order of the table rows. The result is stored in the file MachidaRosenberg2012\_Table3\_numbers.txt. We then used the function call find\_conjugacy\_types\_as\_conjugates\_of\_

MR2012Table3(listoftypes,MR2012tbl\_reordered,"human") in the script finding\_intents\_up\_to\_conj\_as\_conjugates\_of\_MR2012\_Table3.py to find for each of our 48 representative monoids *M* from intents\_up\_to\_conj.txt a monoid *F* listed in some row of [6, Table 3] and a permutation  $s \in Sym(3)$ such that  $M = F^s$ . The result of this computation has been stored in the file intents\_up\_to\_conj\_as\_conjugates\_of\_MachidaRosenberg2012\_ Table3.txt and identifies the 48 representatives from intents\_up\_to\_conj.txt in a one-to-one fashion with the representatives given in [6, Table 3]. This identification is also mentioned in  $[1,$  Table 2, and, since all 48 representatives from  $[6,$ Table 3] occur there, *the conjugacy classes computed in [6] and those represented by the monoids in* intents\_up\_to\_conj.txt *coincide*; moreover intents\_up\_ to\_conj\_as\_conjugates\_of\_MachidaRosenberg2012\_Table3.txt shows how the two presentations can be translated into each other.

# **5 Comparison with results from [7]**

We also compared our computational results with the data available from [7, Tables 2–6, pp. 61–67]. To this end we first collected these five tables as accurately as possible in the files MachidaRosenberg2013\_Table2.txt up to MachidaRosenberg2013 Table6.txt. We then joined these tables into one large table in the file MachidaRosenberg2013 Tables2-6.txt. As the next step we reordered the columns of this table (as in Section 4) in ascending order of the hash values  $9 \cdot s(0) + 3 \cdot s(1) + s(2)$  of the unary operations  $s \in \mathcal{O}_A^{(1)}$  that form the monoids represented in the table. The reordered table can be found (with additional information that was added later) in the file MachidaRosenberg2013\_ Tables2-6\_ordered\_0-26.txt.

We then inserted the contents of this table as MR2013tbl\_reordered into the python script compute\_hash\_values\_of\_192\_monoids.py. With the call print converted list to integers(MR2013tbl reordered) we produced a list of the 192 monoids listed in [7, Tables 2–6] as sequences of integers in {0*, . . . ,* 26} (keeping the order of the monoids as listed in MachidaRosenberg2013 Tables2-6.txt). The result of this computation is contained the file MachidaRosenberg2013\_Tables2-6\_numbers.txt. With the help of the call print\_hash\_values\_for\_list\_of\_characteristic\_ tuples(MR2013tbl\_reordered) we computed hash values for the 192 monoids from [7, Tables 2–6]; the result is shown in the csv-file MachidaRosenberg2013\_ Tables2-6\_ordered\_0-26.csv (with ; as a delimiter) and has also been added to the table MachidaRosenberg2013\_Tables2-6\_ordered\_0-26.txt. We further included the list of the 192 intents as shown in the file 192\_intents.txt in the list intentlist. We then called print hash values for list of sets of integers at most m(intentlist, 26) in order to produce hash values for our 192 intents; the result is contained in the file 192\_intent\_hashes.txt and has been combined with the contents of 192\_extents\_sorted.txt and 192\_intents\_sorted.txt in the file 192\_concepts\_with\_hashes.csv. We then pasted the files MachidaRosenberg2013\_Tables2-6\_ordered\_0-26.csv and 192\_concepts\_with\_hashes.csv in order to create the open document spreadsheet file 192\_monoids\_and\_concepts.ods.

We sorted each part of the pasted table (in 192 monoids and concepts.ods) individually by ascending order of the hash values of the monoids; after that in

each table row the two hash values coincided, thus we identified which monoid from MachidaRosenberg2013\_Tables2-6\_ordered\_0-26.txt corresponds to which of the 192 concepts as listed in 192\_extents\_sorted.txt and 192\_ intents\_sorted.txt or concepts\_new\_format.txt. The result is that the monoids we computed in 192\_intents.txt *are in a one-to-one correspondence* with the monoids in [7, Tables 2–6]. We thus prepended the suitable labels used for the monoids in [7, Tables  $2-6$ ] to the corresponding monoids in 192 concepts\_with\_hashes.csv, resulting in the file 192\_concepts\_with\_hashes\_ and\_MR2013labels.csv. We then sorted the whole table in 192\_monoids\_ and\_concepts.ods once in the order of the lines of 192\_intents\_sorted.txt (or of concepts\_new\_format.txt); this produced the file 192\_monoids\_and\_ concepts\_reordered\_by\_concepts.csv. We also sorted the whole table in the order of the monoids as listed in [7, Tables 2–6]; this produced the file 192\_ monoids\_and\_concepts\_reordered\_by\_MRmonoids.csv. Both files witness the fact that the monoids we computed via [2] and the context produced by context\_ of commutation code.cpp are indeed the same as those shown by Machida and Rosenberg in [7].

We have also used the file 192\_monoids\_and\_concepts\_reordered\_by\_ MRmonoids.csv to update MachidaRosenberg2013 Tables2-6 ordered 0-26. txt with the appropriate extents that witness the monoids from [7].

# **6 Overview of files with their uses**







MachidaRosenberg2012\_Table3\_ numbers.txt

finding\_intents\_up\_to\_conj\_as\_ conjugates\_of\_MR2012\_Table3.py

intents\_up\_to\_conj\_ as\_conjugates\_of\_ MachidaRosenberg2012\_Table3.txt

MachidaRosenberg2013\_Table5.txt a faithful representation of [7, Table 5] MachidaRosenberg2013\_Table6.txt a faithful representation of [7, Table 6] MachidaRosenberg2013\_Tables2-6. txt

MachidaRosenberg2013\_Tables2-6\_ ordered\_0-26.txt

a representation of the content of [6, Table 3] by listing the unary operations  $s_i$  in each monoid from MachidaRosenberg2012\_Table3\_ ordered\_0-26.txt in ascending order of their hash values  $j$ ; each line represents a monoid (row) of [6, Table 3], the order of the rows has been left unchanged. The file has been computed via calling print converted list\_to\_integers(MR2012tbl\_reordered)

in finding\_intents\_up\_to\_conj\_as\_ conjugates\_of\_MR2012\_Table3.py

a python script finding the 48 conjugacy types from intents\_up\_to\_conj.txt as conjugates of the 48 representative monoids in [6, Table 3] as they are represented in MachidaRosenberg2012\_ Table3\_ordered\_0-26.txt. Calls in this script have produced the files MachidaRosenberg2012 Table3 numbers.txt and intents\_up\_to\_conj\_as\_conjugates\_ of\_MachidaRosenberg2012\_Table3.txt. output of the script finding\_intents\_up\_to\_ conj\_as\_conjugates\_of\_MR2012\_Table3.py via the call find\_conjugacy\_types\_as\_ conjugates\_of\_MR2012Table3(listoftypes, MR2012tbl\_reordered,"human"). MachidaRosenberg2013\_Table2.txt a faithful representation of [7, Table 2]

MachidaRosenberg2013\_Table3.txt a faithful representation of [7, Table 3]

- MachidaRosenberg2013\_Table4.txt a faithful representation of [7, Table 4]
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a faithful representation of [7, Tables 2–6]; obtained by concatenating the tables in the files MachidaRosenberg2013\_Table2.txt, MachidaRosenberg2013 Table3.txt, MachidaRosenberg2013\_Table4.txt, MachidaRosenberg2013\_Table5.txt, MachidaRosenberg2013 Table6.txt in this order.

a representation of the table in MachidaRosenberg2013\_Tables2-6.txt after re-ordering the columns such that the unary operations are listed in the order  $s_0 = 000, \ldots, s_{26} = 222$ ; moreover hash values from MachidaRosenberg2013\_Tables2-6\_ ordered\_0-26.csv and witnesses from 192\_monoids\_and\_concepts\_reordered\_by\_ MRmonoids.csv have been added to this file.



192\_monoids\_and\_concepts.ods using the command paste -d ';' on MachidaRosenberg2013\_Tables2-6\_ordered\_ 0-26.csv and 192\_concepts\_with\_hashes. csv we produced a combined table with 192 rows; we used spreadsheet software to sort both parts of the combined table according to the column containing the hash values; in each row we found identical hash values, that is, matching monoids; we then sorted the whole table once according to the order given by the lines of MachidaRosenberg2013\_<br>Tables2-6\_ordered\_0-26.csv and once Tables2-6 ordered 0-26.csv and once according to the order given by the lines of 192\_concepts\_with\_hashes.csv, the former resulting in 192 monoids and concepts reordered\_by\_MRmonoids.csv and the latter in 192\_monoids\_and\_concepts\_reordered\_ by\_concepts.csv; both tables are also contained in the open document spreadsheet 192 monoids and concepts.ods. We also prepended the monoid labels used in [7, Tables 2–6] to the corresponding monoids in 192 concepts with hashes.csv, resulting in the file 192\_concepts\_with\_hashes\_and\_ MR2013labels.csv.

the 192 concepts listed in 192\_concepts\_with\_ hashes.csv together with the appropriate identifiers for each monoid (intent) used in [7, Tables 2–6]. The file was obtained by a paste  $-d$  '; ' of the column of identifiers extracted from 192 monoids and concepts reordered by concepts.csv and the file 192 concepts\_with\_hashes.csv

the 192 centralising monoids in the order of the file 192 concepts with hashes.csv (i.e., the order of the file 192\_intents.txt,  $\equiv$  as in concepts\_new\_format.txt) with the corresponding monoid from [7, Tables 2–6] mentioned in the same line; this file was produced by exporting data from a sheet in 192\_monoids\_and\_ concepts.ods.

192\_concepts\_with\_hashes\_and\_ MR2013labels.csv

192\_monoids\_and\_concepts\_ reordered\_by\_concepts.csv



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