Goldbach's Conjecture — A Route to the Inconsistency of Arithmetic

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Abstract. This paper proves an inconsistency in Peano arithmetic (PA). The contradiction we derive is based on two properties of a specific set which we use to reformulate a strengthened form of the strong Goldbach conjecture.

Notations. Let $\mathbb N$ denote the natural numbers starting from 1, let $\mathbb N_n$ denote the natural numbers starting from $n > 1$ and let \mathbb{P}_3 denote the prime numbers starting from 3.

Strengthened strong Goldbach conjecture (SSGB): *Every [even](http://en.wikipedia.org/wiki/Even_and_odd_numbers) [integer](http://en.wikipedia.org/wiki/Integer) greater than 6 can be expressed as the sum of two different [primes.](http://en.wikipedia.org/wiki/Prime_number)*

Theorem. *PA is contradictory, i.e. the statement FALSE can be derived*.

Proof. We define the set $S_g := \{ (pk, mk, qk) | k, m \in \mathbb{N}$; p, $q \in \mathbb{P}_3$, p < q; m = (p + q) / 2 }.

SSGB is equivalent to saying that every integer $x \ge 4$ is the arithmetic mean of two different odd primes and so it is equivalent to saying that all integers $x \geq 4$ appear as m in a middle component mk of S_g . So, by the definition of S_g we have

SSGB \leq \forall $x \in \mathbb{N}_4$ \exists (pk, mk, qk) \in S_g $x = m$. $-SSGB \leq 3$ $x \in \mathbb{N}_4$ \forall (pk, mk, qk) \in S_q $x \neq m$.

The set S_g has the following two properties.

First, the whole range of \mathbb{N}_3 can be expressed by the triple components of S_g ("*covering*"), because every integer $x \ge 3$ can be written as some pk with $k = 1$ when x is prime, as some pk with $k \neq 1$ when x is composite and not a power of 2, or as $(3 + 5)k / 2$ when x is a power of 2; $p \in \mathbb{P}_3$, $k \in \mathbb{N}$. So we have

(C) $\forall x \in \mathbb{N}_3$ \exists (pk, mk, qk) \in S_g $x = pk$ \lor $x = mk = 4k$.

A few examples of the covering:

- x = 19: (**19∙1**, 21∙1, 23∙1), (**19∙1**, 60∙1, 101∙1)
- x = 36: (**3∙12**, 7∙12, 11∙12)

x = 38: (**19∙2**, 21∙2, 23∙2)

- x = 42: (**3∙14**, 5∙14, 7∙14), (**7∙6**, 9∙6, 11∙6)
- x = 64: (3∙16, **4∙16**, 5∙16)
- x = 10000: (**5∙2000**, 6∙2000, 7∙2000).

Second, according to the statement SSGB, all pairs (p, q) of distinct odd primes are used in the definition of the set S^g ("*maximality*"). So we have

(M) \forall p, q $\in \mathbb{P}_3$, p $\lt q$ \forall k $\in \mathbb{N}$ (pk, mk, qk) \in S_g, where m = (p + q) / 2.

The proof is motivated by the following view.

There are two possibilities for Sg, exactly one of which must occur: Either there is an $n \in \mathbb{N}_4$ *in addition to all the numbers m defined in* S_q *or there is not. The latter is equivalent to SSGB and the former is equivalent to SSGB.*

Since, due to (C), every n given by \neg *SSGB as well as every multiple nk, k* \in N, *equals a component of some S^g triple that exists by definition, the covering of* 3 *by the S^g triples if n exists (SSGB) is equal to that if n does not exist (SSGB). This causes a contradiction because in the case SSGB the numbers m defined in* S_g *take all integer values* $x \ge 4$ whereas in the case \neg SSGB they don't.

First of all, we note that each of the two properties (C) and (M) is a condition sine qua non for the proof, for the following reasons.

 \neg (C) immediately implies \neg SSGB, since an n \geq 4 different from all S_q triple components pk, mk, qk is in particular different from all m in Sg.

The proof would no longer be possible if, for example, we omitted the factor k in the definition of Sg, because then the corresponding (C) could no longer be guaranteed.

Similarly, the property (M) rules out the possibility that there is an $n \geq 4$ different from all m (i.e. \neg SSGB) and n is the arithmetic mean of a pair of primes not used in S_g. Thus (M) excludes the possibility that \neg SSGB applies due to a missing prime number pair. This means that the proof would no longer be possible here either if we left out any prime number pair in the formulation of SSGB and Sg.

We will now show that $((C) \wedge (M))$ leads to a contradiction.

The following proof is independent of the choice of n if there is more than one in the case of SSGB. For example, the minimal such n works.

We split S_g into two complementary subsets: For any $y \in \mathbb{N}_3$, S_g = S_g+(y) ∪ S_g-(y), with $S_g+(y) := \{ (pk, mk, qk) \in S_g | \exists k' \in \mathbb{N} \text{ } pk = yk' \lor mk = yk' \lor qk = yk' \}$ and $S_g(y) := \{ (pk, mk, qk) \in S_g \mid \forall k' \in \mathbb{N} \text{ } pk \neq yk' \land mk \neq yk' \land qk \neq yk' \}.$

Let $n \in \mathbb{N}_4$ be given by \neg SSGB as described above. Then, we have

(*) \neg SSGB => S_g = S_g +(n) ∪ S_g -(n).

More precisely, under the assumption \neg SSGB with the associated n the set S_g can be written as the disjoint union of the following triples.

(i) S_g triples of the form (pk = nk', mk, qk) with $k = k'$ in case n is prime, due to (C)

(ii) S_g triples of the form (pk = nk', mk, qk) with $k \neq k'$ in case n is composite and not a power of 2, due to (C)

(iii) S_q triples of the form $(3k, 4k = nk', 5k)$ in case n is a power of 2, due to (C)

(iv) all remaining S^g triples of the form (pk = nk', mk, qk), (pk, mk = nk', qk) or $(pk, mk, qk = nk')$

and

(v) S_g triples of the form (pk \neq nk', mk \neq nk', qk \neq nk'), i.e. those S_g triples where none of the nk' equals a component.

So, $S_g+(n)$ is the union of the triples of the above types (i) to (iv) and $S_g-(n)$ is the union of the triples of type (v).

Now, we define $S_1 := \{ (pk, mk, qk) \in S_g \mid \neg SSGB holds \}$ $S_2 := \{ (pk, mk, qk) \in S_g \mid SSGB holds \}.$

So, by (*) we obtain

(1) \neg SSGB => S₁ = S_g = S_g+(n) ∪ S_g-(n).

Since $S_g+(n) \cup S_g-(n)$ is independent of n, we can write

(1') $\forall y \in \mathbb{N}$ ₃ \neg SSGB => S₁ = S_g = S_g + (y) ∪ S_g-(y).

Under the assumption SSGB there is no n as above. Therefore, under this assumption, we can choose an arbitrary $y \in \mathbb{N}_3$ such that $S_g = S_g+(y) \cup S_g-(y)$. So, we obtain

(2)
$$
\forall y \in \mathbb{N}_3
$$
 SSGB \Rightarrow S₂ = S_g = S_g+(y) \cup S_g-(y).

We will make use of the following trivial principle.

If two sets of (possibly infinitely many) x-tuples are equal, then the sets of their corresponding i-th components are equal; $1 \le i \le x$.

To this end, for each $k \in \mathbb{N}$ we define

 $M(k) := \{ mk \mid (pk, mk, qk) \in S_g \}$

 $M_1(k) := \{ mk \mid (pk, mk, qk) \in S_1 \}$

 $M_2(k) := \{ mk \mid (pk, mk, qk) \in S_2 \}.$

Then, applying the principle above to the middle component of the triples (pk, mk, qk), (1') and (2) imply

(3) $\forall k \in \mathbb{N}$ $\forall y \in \mathbb{N}$

 $(-SSGB \Rightarrow M_1(k) = M(k) = \{ mk \mid (pk, mk, qk) \in S_g+(y) \cup S_g-(y) \})$

 \wedge

(4) \forall **k** \in \mathbb{N} \forall **y** \in \mathbb{N} ₃

(SSGB => M2(k) = M(k) = { mk | (pk, mk, qk) Sg+(y) ∪ Sg-(y) }).

We set M := $M(1)$, M₁ := $M_1(1)$ and M₂ := $M_2(1)$. So we get

(3') \forall y ∈ \mathbb{N}_3 (\neg SSGB => M₁ = M = { m | (p, m, q) ∈ S_g+(y) ∪ S_g-(y) })

 \wedge

(4') \forall y ∈ \mathbb{N}_3 (SSGB => M₂ = M = { m | (p, m, q) ∈ S_g+(y) ∪ S_g-(y) }).

Since for every $y \in \mathbb{N}_3$ S_g+(y) ∪ S_g-(y) equals S_g, there is a set X such that for every $y \in \mathbb{N}_3$ ${m \mid (p, m, q) \in S_{g+(y) \cup S_{g-(y)}}$ equals X. So, from (3') and (4') we obtain

(5) $(\neg SSGB \Rightarrow M_1 = M = X) \land (SSGB \Rightarrow M_2 = M = X)$.

The set X is a free variable in (5) that is either equal to \mathbb{N}_4 or to some non-empty proper subset Y of \mathbb{N}_4 . Therefore, (5) splits as follows.

(5.1) ($(\neg SSGB \Rightarrow M_1 = M = N_4)$ \land (SSGB => $M_2 = M = N_4$)) \vee

(5.2) $((\neg SSGB \Rightarrow M_1 = M = Y \neq N_4) \land (SSGB \Rightarrow M_2 = M = Y \neq N_4)).$

On the other hand, under the assumption SSGB the numbers m defined in S_g take all integer values $x \ge 4$ whereas under \neg SSGB they don't. Therefore, we have

(6.1) SSGB => $M_2 = M = N_4$

(6.2) \neg SSGB => M₁ = M \neq N₄.

 $((5.1) \vee (5.2))$ together with (6.1) and (6.2) yields

$$
((7.11) \text{ -SSGB} \implies M_1 = M = \mathbb{N}_4
$$

\n
$$
\wedge
$$

\n(7.12) \text{ -SSGB} \implies M_1 = M \neq \mathbb{N}_4
\n(7) ∨
\n((7.21) SSGB \implies M_2 = M \neq \mathbb{N}_4

 \hat{I} (7.22) SSGB => M₂ = M = \mathbb{N}_4).

Since

 \neg SSGB => M₁ = M and SSGB => M₁ = { } \neq M

and

SSGB => M_2 = M and \neg SSGB => M_2 = { } \neq M,

from ((7.11) \wedge (7.12)) we get that $M = \mathbb{N}_4$ and $M \neq \mathbb{N}_4$ if \neg SSGB actually is true, and from ((7.21) \wedge (7.22)) we get that M = \mathbb{N}_4 and M \neq \mathbb{N}_4 if SSGB actually is true.

For two statements P and Q, in contrast to the conditional statement

if P is assumed to be true, then $(Q \land \neg Q)$

which symbolically is ($P \Rightarrow (Q \land \neg Q)$),

the statement

if P is actually true, then $(Q \land \neg Q)$ is symbolically (P \land (P => (Q $\land \neg Q$))).

Therefore, from (7) we obtain

(8.1) $(\neg SSGB \land (\neg SSGB \Rightarrow (M = N_4 \land M \neq N_4)))$ \vee

(8.2) (SSGB \wedge (SSGB => $(M = \mathbb{N}_4 \wedge M \neq \mathbb{N}_4)$).

Then, ($(8.1) \vee (8.2)$) yields (FALSE \vee FALSE), which is equivalent to FALSE.