

Necessary Conditions for $K/2$ Degrees of Freedom

Recep Gül*, Helmut Bölcskei*, and Shlomo Shamai (Shitz)[†]

*Dept. IT & EE, ETH Zürich, Zürich, Switzerland

[†]Dept. EE, Technion-Israel Institute of Technology, Haifa, Israel

Abstract—Stotz et al., 2016, reported a sufficient (injectivity) condition for each user in a K -user single-antenna constant interference channel to achieve $1/2$ degree of freedom. The present paper proves that this condition is necessary as well and hence provides an equivalence characterization of interference channel matrices allowing full degrees of freedom.

I. INTRODUCTION

Cadambe and Jafar [1], [2] proposed a signaling scheme—known as interference alignment—that exploits time-frequency selectivity to achieve $K/2$ degrees of freedom (DoF) in K -user single-antenna interference channels (ICs). In [3] and [4] it was shown that $K/2$ DoF can also be achieved in ICs with constant channel matrix, i.e., in the absence of selectivity. Wu et al. [5] developed a general formula for the number of DoF in single-antenna ICs, extended to vector ICs in [6]. This formula can, however, be difficult to evaluate as it is expressed in terms of Rényi information dimension [7]. Building on the work by Wu et al. [5] and a recent breakthrough result in fractal geometry by Hochman [8], Stotz and Bölcskei [9] derived a DoF-formula for single-antenna ICs, which is exclusively in terms of Shannon entropy; this formula was then used to develop an explicit sufficient condition for achieving $K/2$ DoF. Stotz et al. [10] later identified an even more general sufficient (injectivity) condition for each user to achieve $1/2$ DoF and hence $K/2$ DoF in total.

The main contribution of the present paper is to establish that the sufficient condition in [10] is also necessary for each user to achieve $1/2$ DoF in fully connected ICs¹. The tools used in the proof of this result are the DoF-formula developed in [9] and the entropy version of the Plünnecke-Ruzsa inequality [11].

II. SYSTEM MODEL

We consider a single-antenna K -user (additive) IC with fully connected channel matrix $\mathbf{H} = (h_{ij})_{1 \leq i, j \leq K} \in \mathbb{R}^{K \times K}$, i.e., $h_{ij} \neq 0$ for all i, j , and input-output relation

$$Y_i = \sqrt{\text{snr}} \sum_{j=1}^K h_{ij} X_j + Z_i, \quad i = 1, \dots, K,$$

where $X_i \in \mathbb{R}$ is the input at the i -th transmitter, $Y_i \in \mathbb{R}$ is the output at the i -th receiver, and $Z_i \in \mathbb{R}$ is noise of

¹Note that the sufficient condition in [10] applies to all ICs, whereas here we restrict ourselves to fully connected ICs.

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absolutely continuous distribution satisfying $h(Z_i) > -\infty$ and $H(\lfloor Z_i \rfloor) < \infty$. The input signals are independent across transmitters, and noise is i.i.d. across users and channel uses.

The channel matrix \mathbf{H} is known perfectly at all transmitters and receivers. We impose the average power constraint

$$\frac{1}{n} \sum_{k=1}^n (x_i^{(k)})^2 \leq 1$$

on codewords $(x_i^{(1)} \dots x_i^{(n)})$ of block-length n transmitted by user $i = 1, \dots, K$. The number of DoF of the channel is

$$\text{DoF}(\mathbf{H}) := \limsup_{\text{snr} \rightarrow \infty} \frac{\overline{C}(\mathbf{H}; \text{snr})}{\frac{1}{2} \log \text{snr}},$$

where $\overline{C}(\mathbf{H}; \text{snr})$ stands for the sum-capacity of the IC.

We note that [10] studies conditions for each user to achieve $1/2$ DoF as opposed to $K/2$ DoF for all users in total. This is done to avoid trivial exceptions induced by particular network topologies reflected by zeros in the IC matrix. Here, we also study conditions for each user to achieve $1/2$ DoF. We remark, however, that topology-induced exceptions cannot occur for fully connected ICs.

III. MAIN RESULT

We start by reviewing the sufficient (injectivity) condition for each user to achieve $1/2$ DoF identified in [10]. Let $\check{\mathbf{h}} \in \mathbb{R}^{K(K-1)}$ denote the vector containing the off-diagonal elements of \mathbf{H} . There are

$$\varphi(d) := \binom{K(K-1) + d}{d}$$

monomials² in $K(K-1)$ variables of degree not larger than d , which we enumerate as $f_1, \dots, f_{\varphi(d)}$. We will further need the following sets

$$\mathcal{W}_{N,d} := \left\{ \sum_{i=1}^{\varphi(d)} a_i f_i(\check{\mathbf{h}}) : a_1, \dots, a_{\varphi(d)} \in \{0, \dots, N-1\} \right\}$$

$$\mathcal{W} := \bigcup_{d \geq 0} \bigcup_{N \geq 1} \mathcal{W}_{N,d}. \quad (1)$$

It was shown in [10, Th. 1] that for each user to achieve $1/2$ DoF, it is sufficient that the following condition be satisfied for either \mathbf{H} directly or at least one scaled version of \mathbf{H} , where

²A monomial in k variables x_i is an expression of the form $x_1^{n_1} x_2^{n_2} \dots x_k^{n_k}$, where $n_1, \dots, n_k \in \mathbb{N}$, and the degree of the monomial is $n_1 + \dots + n_k$.

scaling, as defined in [10], refers to a finite number of scalings of individual rows and columns by nonzero constants:

For each $i = 1, \dots, K$, the map

$$\begin{aligned} \mathcal{W} \times \mathcal{W} &\rightarrow \mathcal{W} + h_{ii}\mathcal{W} \\ (w_1, w_2) &\mapsto w_1 + h_{ii}w_2 \end{aligned} \quad (**)$$

is injective.

The main result of the present paper is as follows.

Theorem 1. Let \mathcal{W} be as in (1). Then, for almost all fully connected IC matrices \mathbf{H} , the following condition is necessary for each user to achieve $1/2$ DoF: Either the channel matrix \mathbf{H} itself or at least one scaled version thereof satisfies condition (**).

Proof: See Section V.

With the matching sufficient condition in [10, Th. 1] this yields an equivalence characterization of full-DoF-achieving IC matrices.

Remark 1. Note that thanks to \mathcal{W} containing integer linear combinations of monomials in the off-diagonal elements of \mathbf{H} , the condition (**) is exclusively in terms of channel coefficients.

Remark 2. Almost all channel matrices \mathbf{H} are fully connected. Since an almost all subset of an almost all set is itself almost all, the set of channel matrices \mathbf{H} not covered by Theorem 1 has measure zero.

IV. BALANCING RESULTS

We start with a simple modification of [9, Th. 3], with the proof omitted due to space constraints.

Proposition 1. For almost all fully connected IC matrices \mathbf{H} the following holds: If each user achieves $1/2$ DoF, then for $\varepsilon \in (0, 1/2)$, there exist independent discrete random variables V_1, \dots, V_K of finite entropy such that

$$\frac{1}{2} - \varepsilon \leq \frac{H\left(\sum_{j=1}^K h_{ij}V_j\right) - H\left(\sum_{j \neq i}^K h_{ij}V_j\right)}{\max_{i=1, \dots, K} H\left(\sum_{j=1}^K h_{ij}V_j\right)}, \quad (2)$$

for $i = 1, \dots, K$, with the denominator of (2) nonzero.

The following lemma, based on Proposition 1, states ‘‘balancing’’ results on the entropies of signal and interference contributions and will turn out instrumental in the proof of our main statement.

Lemma 1. For $\varepsilon \in (0, 1/2)$, let V_1, \dots, V_K be the corresponding independent discrete random variables satisfying (2) for the fully connected IC matrix \mathbf{H} . Then,

$$\frac{H\left(\sum_{j \neq i}^K h_{ij}V_j\right)}{H(V_i)} = 1 + \mathcal{O}(\varepsilon), \quad \text{for } i = 1, \dots, K, \quad (3)$$

$$\frac{H(V_i)}{H(V_j)} = 1 + \mathcal{O}(\varepsilon), \quad \text{for } i, j \in \{1, \dots, K\}, i \neq j, \quad (4)$$

$$\frac{H\left(\sum_{j=1}^K h_{ij}V_j\right)}{H(V_i)} = 2 + \mathcal{O}(\varepsilon), \quad \text{for } i = 1, \dots, K. \quad (5)$$

Proof: Starting from (2), we have

$$\frac{1}{2} - \varepsilon \leq \frac{H\left(\sum_{j=1}^K h_{ij}V_j\right) - H\left(\sum_{j \neq i}^K h_{ij}V_j\right)}{H\left(\sum_{j=1}^K h_{ij}V_j\right)}, \quad (6)$$

for $i = 1, \dots, K$. Rearranging terms, we get

$$2H\left(\sum_{j \neq i}^K h_{ij}V_j\right) \leq (1 + 2\varepsilon)H\left(\sum_{j=1}^K h_{ij}V_j\right). \quad (7)$$

Invoking the following inequality for independent discrete random variables X, Y [12, Ex. 2.14]

$$H(X + Y) \leq H(X) + H(Y), \quad (8)$$

on the right-hand side (RHS) of (7) yields

$$(1 - 2\varepsilon)H\left(\sum_{j \neq i}^K h_{ij}V_j\right) \leq (1 + 2\varepsilon)H(V_i). \quad (9)$$

Next, we show that for fully connected ICs

$$H\left(\sum_{j \neq i}^K h_{ij}V_j\right) \geq \frac{(1 - 2\varepsilon)}{(1 + 2\varepsilon)}H(V_i). \quad (10)$$

To this end, w.l.o.g., we assume that $H(V_1) \geq H(V_2) \geq \dots \geq H(V_K)$. Applying [12, Ex. 2.14]

$$H(\alpha X + \beta Y) \geq \max\{H(X), H(Y)\} \quad (11)$$

repeatedly for independent discrete random variables X and Y , and arbitrary $\alpha, \beta \in \mathbb{R} \setminus \{0\}$, with $X = \sum_{j \neq 1, i}^K h_{ij}V_j$, $Y = V_1$, $\alpha = 1$, $\beta = h_{i1}$, for $i = 2, \dots, K$, we obtain

$$(1 - 2\varepsilon)H(V_i) \leq (1 - 2\varepsilon)H(V_1) \leq (1 + 2\varepsilon)H\left(\sum_{j \neq i}^K h_{ij}V_j\right). \quad (12)$$

This proves (10) for $i = 2, \dots, K$. The case $i = 1$ is obtained as follows. First note that

$$(1 - 2\varepsilon)H(V_1) \leq (1 - 2\varepsilon)H\left(\sum_{j \neq i}^K h_{ij}V_j\right) \leq (1 + 2\varepsilon)H(V_i) \quad (13)$$

for all $i \neq 1$, where the first inequality is by (11) and the second by (9). Further, again by (11), we have

$$(1 + 2\varepsilon)H(V_i) \leq (1 + 2\varepsilon)H\left(\sum_{j \neq 1}^K h_{ij}V_j\right),$$

for all $i \neq 1$, and inserting into (13) establishes (10) for $i = 1$. We can now combine (9) and (10) to get

$$\frac{1 - 2\varepsilon}{1 + 2\varepsilon} \leq \frac{H\left(\sum_{j \neq i}^K h_{ij}V_j\right)}{H(V_i)} \leq \frac{1 + 2\varepsilon}{1 - 2\varepsilon} \quad (14)$$

for all i , which yields (3).

To prove (4), we again assume, w.l.o.g., that $H(V_1) \geq \dots \geq H(V_K)$, and simply note that thanks to (13)

$$\frac{1 - 2\varepsilon}{1 + 2\varepsilon} \leq \frac{H(V_K)}{H(V_1)} \leq \frac{H(V_i)}{H(V_j)} \leq \frac{H(V_1)}{H(V_K)} \leq \frac{1 + 2\varepsilon}{1 - 2\varepsilon}, \quad (15)$$

for $i, j = 1, \dots, K$.

Finally, to establish (5), we start by noting that

$$\frac{H\left(\sum_{j=1}^K h_{ij}V_j\right)}{H(V_i)} \geq \left(\frac{1-2\varepsilon}{1+2\varepsilon}\right) \frac{H\left(\sum_{j=1}^K h_{ij}V_j\right)}{H\left(\sum_{j \neq i}^K h_{ij}V_j\right)} \quad (16)$$

$$\geq \frac{2(1-2\varepsilon)}{(1+2\varepsilon)^2},$$

owing to (14) and (7). Using (8) and (14), it follows that

$$H\left(\sum_{j=1}^K h_{ij}V_j\right) \leq \left(1 + \frac{1+2\varepsilon}{1-2\varepsilon}\right) H(V_i). \quad (17)$$

Combining (17) with (16) then yields (5). \blacksquare

We conclude this section by recording, for later use, a simple variation of the entropy version of the Plünnecke-Ruzsa inequality [11].

Lemma 2. [A simple variation of [11, Th. 2.8.2]]: For $\varepsilon \in (0, 1/2)$, assume that there exist independent discrete random variables X, Y_1, \dots, Y_m , all of finite entropy, such that,

$$\frac{H(X + Y_i)}{H(X)} = 1 + \mathcal{O}(\varepsilon), \quad (18)$$

for $i = 1, \dots, m$. Then, for finite m ,

$$\frac{H(X + Y_1 + \dots + Y_m)}{H(X)} = 1 + \mathcal{O}(\varepsilon). \quad (19)$$

Proof: In [11, Th. 2.8.2], set $\log K_i = H(X)\mathcal{O}(\varepsilon)$ for $i = 1, \dots, m$, and take the additive group G as \mathbb{R} . This results in $H(X + Y_1 + \dots + Y_m) \leq H(X) + mH(X)\mathcal{O}(\varepsilon)$. Divide this inequality by $H(X)$ and note that owing to (11) the expression on the LHS of (19) is greater than or equal to 1. The proof is concluded by noting that $\mathcal{O}(\varepsilon)m = \mathcal{O}(\varepsilon)$ for finite m . \blacksquare

V. PROOF OF THEOREM 1

For simplicity of exposition and due to space constraints, we detail the proof for the 3-user case only. The proof for the K -user case follows by induction over the number of users with the 3-user case constituting the base case.

Consider the 3-user fully connected IC matrix

$$\mathbf{H} = \begin{pmatrix} h_{11} & h_{12} & h_{13} \\ h_{21} & h_{22} & h_{23} \\ h_{31} & h_{32} & h_{33} \end{pmatrix}.$$

The proof is effected by contradiction, with the contradiction established by induction on the degrees of polynomials in \mathcal{W} as defined in (1). Towards this contradiction, we assume that \mathbf{H} is in the almost all set covered by Theorem 1, while at the same time each user achieves 1/2 DoF and condition (**) is violated for \mathbf{H} and all scaled versions thereof. In particular, condition (**) must also be violated for

$$\tilde{\mathbf{H}} = \begin{pmatrix} g_1 & 1 & 1 \\ 1 & g_2 & 1 \\ 1 & h & g_3 \end{pmatrix}, \quad g_1, g_2, g_3, h \neq 0, \quad (20)$$

which can be obtained from \mathbf{H} by scaling (cf. [5, p. 259]). It follows from a simple modification of the arguments in [3, Lem. 1] that for fully connected ICs, scaling does not change the number of DoF achieved by each user. Hence, each user achieves 1/2 DoF in $\tilde{\mathbf{H}}$ as well. The reduction to $\tilde{\mathbf{H}}$ is for simplicity of exposition as all elements of the corresponding set \mathcal{W} in (1) are polynomials in h .

We proceed by assuming that the map in (**) is not injective for user $i = 1$ so that condition (**) is violated. This implies the existence of $w_1, w_2, \tilde{w}_1, \tilde{w}_2 \in \mathcal{W}$ such that $w_1 \neq \tilde{w}_1$, $w_2 \neq \tilde{w}_2$, and

$$w_1 + g_1 w_2 = \tilde{w}_1 + g_1 \tilde{w}_2.$$

We can hence write

$$g_1 = \frac{w_1 - \tilde{w}_1}{\tilde{w}_2 - w_2} = \frac{\sum_{p=0}^{d_1} \hat{a}_p h^p}{\sum_{p=0}^{d_2} \hat{b}_p h^p}, \quad (21)$$

where d_1 and d_2 are finite and the coefficients \hat{a}_p, \hat{b}_p are integers. Since $w_1 \neq \tilde{w}_1$ and $w_2 \neq \tilde{w}_2$, the polynomials on the RHS of (21) are nonzero.

Since $\tilde{\mathbf{H}}$ is fully connected, we know from (5) that, for $i = 1$,

$$\frac{H\left(\sum_{p=0}^{d_1} \hat{a}_p h^p V_1 + V_2 + V_3\right)}{H(V_1)} = 2 + \mathcal{O}(\varepsilon). \quad (22)$$

We shall show that this leads to a contradiction, by proving that (2) implies

$$\frac{H\left(\sum_{p=0}^{d_1} \hat{a}_p h^p V_1 + V_2 + V_3\right)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon) \quad (23)$$

for every nonzero finite-degree polynomial in h with integer coefficients a_p and b_p , in particular \hat{a}_p and \hat{b}_p . The implication (2) \implies (23) is established by induction over $d = \max\{d_1, d_2\}$.

Base case ($d = 0$): When $d = 0$, we need to show that

$$\frac{H\left(a_0 V_1 + b_0(V_2 + V_3)\right)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (24)$$

From (3) with $i = 1$ we have

$$\frac{H(V_2 + V_3)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (25)$$

Likewise, (3) with $i = 2$, upon using (4) with $i = 1, j = 2$ yields

$$\frac{H(V_1 + V_3)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (26)$$

As the polynomials in (23) are nonzero, it suffices to prove the statement for $a_0, b_0 \in \mathbb{Z} \setminus \{0\}$. Using (4) with $i = 1, j = 3$, we can replace the denominators of (25) and (26) by $H(V_3)$ and then apply Lemma 2 to get

$$\frac{H(V_1 + V_2 + V_3)}{H(V_3)} = 1 + \mathcal{O}(\varepsilon). \quad (27)$$

Replacing $H(V_3)$ in (27) by $H(V_1)$, which is possible thanks

to (4) with $i = 1, j = 3$, we obtain

$$\frac{H(V_1 + V_2 + V_3)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (28)$$

Applying [5, Th. 14] with $p = a_0, q = b_0, X = V_1$, and $Y = V_2 + V_3$, and dividing the result thereof by $H(V_1)$ yields

$$\frac{H(a_0 V_1 + b_0(V_2 + V_3))}{H(V_1)} - \frac{H(V_1 + V_2 + V_3)}{H(V_1)} \leq \quad (29)$$

$$\tau_{a_0, b_0} \left(\frac{(2H(V_1 + V_2 + V_3) - H(V_1) - H(V_2 + V_3))}{H(V_1)} \right),$$

where $\tau_{a_0, b_0} = 7\lfloor \log |a_0| \rfloor + 7\lfloor \log |b_0| \rfloor + 2$. By (25) and (28), the RHS of (29) equals $\mathcal{O}(\varepsilon)$. We therefore have

$$1 \leq \frac{H(a_0 V_1 + b_0(V_2 + V_3))}{H(V_1)} \leq 1 + \mathcal{O}(\varepsilon), \quad (30)$$

where the first inequality follows from (11). In summary, we have

$$\frac{H(a_0 V_1 + b_0(V_2 + V_3))}{H(V_1)} = 1 + \mathcal{O}(\varepsilon), \quad (31)$$

which establishes (24) as desired.

Induction step. We assume that (23) holds for $d = m - 1$, with $m \geq 1$, and show that this implies (23) for $d = m$. The following lemma contains the central technical result in the induction step.

Lemma 3. For $\varepsilon \in (0, 1/2)$, let V_1, V_2 , and V_3 be the corresponding independent discrete random variables satisfying (2) for the IC matrix $\tilde{\mathbf{H}}$ in (20). Assume further that V_1, V_2 , and V_3 satisfy (23) for $d = \max\{d_1, d_2\} = m - 1$, where $m \geq 1$. Let V_1^* be an independent copy of V_1 , and \tilde{V}_2, \hat{V}_2 and \tilde{V}_3, \hat{V}_3 independent copies of V_2 and V_3 , respectively. Then,

$$\frac{H(h(V_2 + V_3) + a_m h^m V_1^* + b_m h^m (\tilde{V}_2 + \tilde{V}_3))}{H(V_1)} = 1 + \mathcal{O}(\varepsilon) \quad (32)$$

$$\frac{H(h(V_2 + V_3) + \sum_{i=0}^{m-1} a_i h^i V_1 + \sum_{i=0}^{m-1} b_i h^i (\hat{V}_2 + \hat{V}_3))}{H(V_1)} = 1 + \mathcal{O}(\varepsilon) \quad (33)$$

$$\frac{H(\alpha_1 V_1^* + \alpha_2 V_1 + \beta_1 (\tilde{V}_2 + \tilde{V}_3) + \beta_2 (V_2 + V_3))}{H(V_1)} + \mathcal{O}(\varepsilon) \geq \quad (34)$$

$$\frac{H((\alpha_1 + \alpha_2)V_1 + (\beta_1 + \beta_2)(V_2 + V_3))}{H(V_1)},$$

for all $\alpha_1, \alpha_2, \beta_1, \beta_2 \in \mathbb{R}$.

Proof: To establish (32), we apply Lemma 2 with $X = a_m h^m V_1^*$, $Y_1 = h(V_2 + V_3)$, and $Y_2 = b_m h^m (\tilde{V}_2 + \tilde{V}_3)$. The corresponding conditions (18) hold as dividing $X + Y_1$ and $X + Y_2$ by h yields (23), which is satisfied by the induction hypothesis.

For (33) we have to distinguish the cases $m > 1$ and $m = 1$. For $m > 1$, one applies Lemma 2 with $X = \sum_{i=0}^{m-1} a_i h^i V_1$,

$Y_1 = h(V_2 + V_3)$, and $Y_2 = \sum_{i=0}^{m-1} b_i h^i (\hat{V}_2 + \hat{V}_3)$. Again, the corresponding conditions (18) of Lemma 2 are satisfied thanks to the induction hypothesis. For $m = 1$, (33) reduces to

$$\frac{H(h(V_2 + V_3) + a_0 V_1 + b_0 (\hat{V}_2 + \hat{V}_3))}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (35)$$

Using (11) and (28), we get

$$\frac{H(V_2)}{H(V_1)} \leq \frac{H(h\hat{V}_2 + h\hat{V}_3)}{H(V_1)} \leq \frac{H(V_1 + V_2 + V_3)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (36)$$

Thanks to (4), the leftmost term in (36) equals $1 + \mathcal{O}(\varepsilon)$. Hence,

$$\frac{H(h\hat{V}_2 + h\hat{V}_3)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (37)$$

Furthermore, applying (3) with $i = 3$, we have

$$\frac{H(h\hat{V}_2 + V_1)}{H(V_3)} = 1 + \mathcal{O}(\varepsilon), \quad (38)$$

which thanks to (4) with $i = 1, j = 3$ yields

$$\frac{H(h\hat{V}_2 + V_1)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (39)$$

Combining (37) and (39), and applying Lemma 2 with $X = h\hat{V}_2$, $Y_1 = V_1$, and $Y_2 = h\hat{V}_3$, we obtain

$$\frac{H(V_1 + h(\hat{V}_2 + \hat{V}_3))}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (40)$$

Repeating the steps leading from (28) to (31) now results in

$$\frac{H(a_0 V_1 + b_0 h(\hat{V}_2 + \hat{V}_3))}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (41)$$

Using (41) and applying Lemma 2 with $X = a_0 V_1$, $Y_1 = b_0 h(\hat{V}_2 + \hat{V}_3)$, and $Y_2 = h(V_2 + V_3)$, we get

$$\frac{H(h(V_2 + V_3) + a_0 V_1 + b_0 h(\hat{V}_2 + \hat{V}_3))}{H(V_1)} = 1 + \mathcal{O}(\varepsilon), \quad (42)$$

which establishes (35), and thereby completes the proof of (33).

To prove (34), we first apply Lemma 2 with $X = V_1$, $Y_1 = V_2 + V_3$ and $Y_2 = \tilde{V}_2 + \tilde{V}_3$, resulting in

$$\frac{H(V_1 + V_2 + V_3 + \tilde{V}_2 + \tilde{V}_3)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (43)$$

The corresponding conditions (18) of Lemma 2 are satisfied thanks to (28), and upon noting that \tilde{V}_2, \tilde{V}_3 are independent copies of V_2, V_3 . Thanks to (25) and (11), we have

$$1 + \mathcal{O}(\varepsilon) = \frac{H(V_2 + V_3)}{H(V_1)} \leq \frac{H(V_2 + V_3 + \tilde{V}_2 + \tilde{V}_3)}{H(V_1)} \quad (44)$$

$$\leq \frac{H(V_1 + V_2 + V_3 + \tilde{V}_2 + \tilde{V}_3)}{H(V_1)}, \quad (45)$$

which, when combined with (43), yields

$$\frac{H(V_2 + V_3 + \tilde{V}_2 + \tilde{V}_3)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (46)$$

The same line of reasoning, with Lemma 2 applied with $X = V_3$, $Y_1 = V_1$, and $Y_2 = V_1^*$, upon invoking (26), delivers

$$\frac{H(V_1 + V_1^*)}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (47)$$

We next use [5, Lem. 18] with $Z = \beta_1(\tilde{V}_2 + \tilde{V}_3) + \beta_2(V_2 + V_3)$, $X = V_1$, $X' = V_1^*$, $r = \alpha_1$, and $p = \alpha_1 + \alpha_2$ to get³

$$\begin{aligned} & \frac{H((\alpha_1 + \alpha_2)V_1 + \beta_1(\tilde{V}_2 + \tilde{V}_3) + \beta_2(V_2 + V_3))}{H(V_1)} - \frac{\Delta(V_1, V_1^*)}{H(V_1)} \\ & \leq \frac{H(\alpha_1 V_1^* + \alpha_2 V_1 + \beta_1(\tilde{V}_2 + \tilde{V}_3) + \beta_2(V_2 + V_3))}{H(V_1)}, \end{aligned} \quad (48)$$

where $\Delta(V, W) = H(V - W) - \frac{1}{2}H(V) - \frac{1}{2}H(W)$. Note that due to (47), we have

$$\frac{\Delta(V_1, -V_1^*)}{H(V_1)} = \mathcal{O}(\varepsilon). \quad (49)$$

Thanks to [13, Th. 3.5] this implies

$$\frac{\Delta(V_1, V_1^*)}{H(V_1)} = \mathcal{O}(\varepsilon). \quad (50)$$

We next apply [5, Lem. 18] with $Z = (\alpha_1 + \alpha_2)V_1$, $X = V_2 + V_3$, $X' = \tilde{V}_2 + \tilde{V}_3$, $p = \beta_1 + \beta_2$, and $r = \beta_2$. Noting that thanks to (46), (49) continues to hold if we replace $\Delta(V_1, -V_1^*)$ by $\Delta(V_2 + V_3, -(\tilde{V}_2 + \tilde{V}_3))$, we obtain

$$\begin{aligned} & \frac{H((\alpha_1 + \alpha_2)V_1 + (\beta_1 + \beta_2)(V_2 + V_3))}{H(V_1)} \leq \mathcal{O}(\varepsilon) \\ & + \frac{H((\alpha_1 + \alpha_2)V_1 + \beta_1(V_2 + V_3) + \beta_2(\tilde{V}_2 + \tilde{V}_3))}{H(V_1)}. \end{aligned} \quad (51)$$

Finally, combining (48) and (51), we get (34). \blacksquare

We now use Lemma 3 to finalize the induction argument. First, let

$$\begin{aligned} p(h) & := h(V_2 + V_3) + \sum_{i=0}^{m-1} a_i h^i \hat{V}_1 + \sum_{i=0}^{m-1} b_i h^i (\hat{V}_2 + \hat{V}_3) \\ & + a_m h^m V_1^* + b_m h^m (\tilde{V}_2 + \tilde{V}_3), \end{aligned}$$

where \tilde{V}_2, \hat{V}_2 and \tilde{V}_3, \hat{V}_3 are independent copies of V_2 and V_3 , respectively.

Applying Lemma 2 with $X = h(V_2 + V_3)$, $Y_1 = a_m h^m V_1^* + b_m h^m (\tilde{V}_2 + \tilde{V}_3)$, and $Y_2 = \sum_{i=0}^{m-1} a_i h^i \hat{V}_1 + \sum_{i=0}^{m-1} b_i h^i (\hat{V}_2 + \hat{V}_3)$, noting that the corresponding conditions (18) are satisfied thanks to (32) and (33), and invoking (26), we get

$$\frac{H(p(h))}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (52)$$

³Note that [5, Lem. 18] assumes p and r to be nonzero integers. A closer inspection of the proof reveals, however, that the result holds for all $p, r \in \mathbb{R}$.

Now define

$$\begin{aligned} \tilde{p}(h) & := \sum_{i=0}^{m-1} a_i h^i \hat{V}_1 + \sum_{i=0}^{m-1} b_i h^i (\hat{V}_2 + \hat{V}_3) \\ & + a_m h^m V_1^* + b_m h^m (\tilde{V}_2 + \tilde{V}_3). \end{aligned}$$

Thanks to (11) and (52), we have

$$1 + \mathcal{O}(\varepsilon) = \frac{H(p(h))}{H(V_1)} \geq \frac{H(\tilde{p}(h))}{H(V_1)} \geq 1. \quad (53)$$

Applying (34) with $\alpha_1 = \sum_{i=0}^{m-1} a_i h^i$, $\alpha_2 = a_m h^m$, $\beta_1 = \sum_{i=0}^{m-1} b_i h^i$, and $\beta_2 = b_m h^m$, yields

$$\frac{H(\tilde{p}(h))}{H(V_1)} + \mathcal{O}(\varepsilon) \geq \frac{H(\sum_{i=0}^m a_i h^i V_1 + \sum_{i=0}^m b_i h^i (V_2 + V_3))}{H(V_1)}. \quad (54)$$

Thanks to (11), the RHS of (54) cannot be smaller than one. Combining (53) and (54), we hence get

$$\frac{H(\sum_{i=0}^m a_i h^i V_1 + \sum_{i=0}^m b_i h^i (V_2 + V_3))}{H(V_1)} = 1 + \mathcal{O}(\varepsilon). \quad (55)$$

This concludes the induction argument and thereby the proof.

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