

2.4.8 Control Instructions

Instructions in this group affect the flow of control.

```

blocktype ::= typeid | valtype?
instr     ::= ...
            | nop
            | unreachable
            | block blocktype instr* end
            | loop blocktype instr* end
            | if blocktype instr* else instr* end
            | br labelidx
            | br_if labelidx
            | br_table vec(labelidx) labelidx
            | return
            | call funcidx
            | call_indirect tableidx typeid

```

The `nop` instruction does nothing.

The `unreachable` instruction causes an unconditional trap.

The `block`, `loop` and `if` instructions are *structured* instructions. They bracket nested sequences of instructions, called *blocks*, terminated with, or separated by, `end` or `else` pseudo-instructions. As the grammar prescribes, they must be well-nested.

A structured instruction can consume *input* and produce *output* on the operand stack according to its annotated *block type*. It is given either as a *type index* that refers to a suitable *function type*, or as an optional *value type* inline, which is a shorthand for the function type `[] → [valtype?]`.

Each structured control instruction introduces an implicit *label*. Labels are targets for branch instructions that reference them with *label indices*. Unlike with other *index spaces*, indexing of labels is relative by nesting depth, that is, label 0 refers to the innermost structured control instruction enclosing the referring branch instruction, while increasing indices refer to those farther out. Consequently, labels can only be referenced from *within* the associated structured control instruction. This also implies that branches can only be directed outwards, “breaking” from the block of the control construct they target. The exact effect depends on that control construct. In case of `block` or `if` it is a *forward jump*, resuming execution after the matching `end`. In case of `loop` it is a *backward jump* to the beginning of the loop.

Note: This enforces *structured control flow*. Intuitively, a branch targeting a `block` or `if` behaves like a break statement in most C-like languages, while a branch targeting a `loop` behaves like a continue statement.

Branch instructions come in several flavors: `br` performs an unconditional branch, `br_if` performs a conditional branch, and `br_table` performs an indirect branch through an operand indexing into the label vector that is an immediate to the instruction, or to a default target if the operand is out of bounds. The `return` instruction is a shortcut for an unconditional branch to the outermost block, which implicitly is the body of the current function. Taking a branch *unwinds* the operand stack up to the height where the targeted structured control instruction was entered. However, branches may additionally consume operands themselves, which they push back on the operand stack after unwinding. Forward branches require operands according to the output of the targeted block’s type, i.e., represent the values produced by the terminated block. Backward branches require operands according to the input of the targeted block’s type, i.e., represent the values consumed by the restarted block.

The `call` instruction invokes another *function*, consuming the necessary arguments from the stack and returning the result values of the call. The `call_indirect` instruction calls a function indirectly through an operand indexing into a *table* that is denoted by a *table index* and must have type *funcref*. Since it may contain functions of heterogeneous type, the callee is dynamically checked against the *function type* indexed by the instruction’s second immediate, and the call is aborted with a trap if it does not match.