

VICEPRESIDENCIA PRIMERA DEL GOBIERNO **IISTERIO
ASUNTOS ECONÓMICOS ANSFORMACIÓN DIGITAL**

Towards a smart and efficient telecom infrastructure meeting current and future industry needs (TIMING SP-2) (Ref. TSI-063000-2021-148)

Deliverable D1.1

Year 1 report on component development, integration, testing and validation

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Version 1.5

Date 31 Jan-2024

Distribution PUBLIC (PU)

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REVISION HISTORY

LIST OF AUTHORS

GLOSSARY

EXECUTIVE SUMMARY

The aim of this deliverable is to report about the status of the different components/modules of the overall TIMING architecture in support of E2E services, defined in the framework of SP1 project. In particular, it reports the status of the prototyping of the Wi-Fi and Ethernet TSN nodes, the TSN controller, the TSN Connectivity Manager, the scheduling algorithms for the Wi-Fi node, TSN models to be used by the Digital Twin (DT) to estimate the KPIs, metro network solution, and the industrial application.

Moreover, this deliverable contains a preliminary version of the operational workflow for the service provisioning over the E2E infrastructure, putting emphasis on the interactions among the different components to materialize the required resources configuration to support the E2E services with the required KPIs.

Finally, this deliverable also contains the preliminary integration tests conducted so far, with special emphasis on the data plane interconnection tests between TSN Wi-Fi node and the TSN Ethernet node.

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1 INTRODUCTION

This deliverable reports, firstly, the ongoing status of implementation of the different components/modules that have been defined in the architecture design, reported in the SP-1 deliverables. More details about the overall architecture can be found in deliverable D1.2 [1] and D1.3 [2]. The implementation of the different modules as well as their evaluation represents the required step prior to their integration to materialize the TIMING architecture properly designed to support E2E service provisioning. The different components that compose the architecture can be listed as:

- Wi-Fi-based TSN nodes
- Scheduling algorithms for Wi-Fi TSN
- Ethernet TSN nodes
- TSN SDN controller
- Connectivity Manager
- TSN models/DT
- Optical transport/metro SDN Controller
- Industrial application.

For each component, this deliverable reports the current status of the development, as well some stand-alone tests to validate the implementation. Additionally, this deliverable also reports some preliminary integration tests, mostly at the data plane level, that is, related to data plane connectivity between the Wi-Fi node and the Ethernet TSN node.

The overall integration tests will be reported in the upcoming deliverable D1.2 of SP2 project.

2 TIMING ARCHITECTURE COMPONENTS

2.1 TSN WI-FI NODE

2.1.1 Description

The SW/HW architecture of the WI-FI TSN Node can be divided in three sections as shown [Figure 2-1.](#page-9-3)

- **Wi-Fi TSN Node's hardware (VHDL)**, including key components such as physical communications interfaces (Ethernet TSN and Wi-Fi TSN), a PTP Hardware clock, and a RT traffic translation entity. The Ethernet TSN features at least one physical port, supporting IEEE 802.1 standards such as IEEE 802.11AS for PTP and IEEE 802.1Qbv for scheduled and mixed-critical traffic. The Wi-Fi TSN modem employs a w-SHARP implementation, incorporating time-sensitive communication and synchronization techniques as in the Wired TSN interface.
- **The main processor** who operates a general-purpose operating system, hosting the primary control application, hardware and control access drivers, a TCP/IP network stack, and various applications for Wi-Fi TSN node operation. These applications consist of a PTP stack for IEEE 802.1AS compliance, a controller interface for external TSN controller communication, and a scheduling algorithm interface for wireless scheduling methods.
- **The real-time co-processor** who hosts the Wi-Fi TSN modem driver, RT software program, and components governing precise-time operations of the Wi-Fi TSN modem. The Wireless TSN driver facilitates frame communication with the Wi-Fi TSN modem, offering a control interface for the scheduling algorithm and TSN controller.

Figure 2-1: Block diagram of the Wi-Fi TSN node HW/SW architecture.

With respect to the status reported in SP1 D1.3 [2], work has been done specially in two main topics:

• Provide fast **reconfiguration of the Wi-Fi-TSN nodes**: Every time the SDN-TSN controller or the Wireless Flow Scheduler demands a change in the configuration, the capability to

receive the new specific TSN configuration and be able to forward it to the connected STAs and apply it as fast as possible.

• Provide **statistics/telemetry/capabilities** to the SDN-TSN controller and the Wireless Flow Scheduler.

Along with these, the first integration tests between Ikerlan's WiFi-TSN nodes and Safran's TSN switches are also reported in section 5.

2.1.2 Report on the Current Status

2.1.2.1 Current status of the implementation

SP1 derivable D1.3 [2], section 2.2.1 defined three stages for the development of the new control software:

- A first stage where the different features (publications of statistics, reception of configuration and propagation of configuration) will be implemented through configuration files. The remote access either from the SDN-TSN Controller or the Wireless Flow Scheduler to change the configuration files or read statistics will be done via FTP or SSH connection. The retransmission of the new configuration from the AP to the STAs is also done via SSH and FTP.
- A second stage, where the configuration files from the SDN-TSN Controller and the Wireless Flow Scheduler are received, and the gathered statistics are offered via an API based in an IP/Socket protocol. Besides, the retransmission of the new configuration from the AP to the STAs will be included in the beacon, at the beginning of each superframe, thus speeding un the reconfiguration process.
- A third stage, where the information to/from the SDN-TSN Controller will continue being received/sent via an API based in a IP/Socket protocol, but where the Wireless Flow Scheduler will be integrated in the real-time software of the TSN Wi-Fi node.

At this point, the first stage is practically completed, and the second stage is ongoing.

On the one hand, from the fast reconfiguration of the WiFi-TSN nodes point of view, the development of two APIs, with the aim of receiving the configuration information, is ongoing. One API would be aimed for the SDN-TSN controller, and the other one for the Wireless Flow Scheduler. The API being developed for the SDN-TSN controller is a Flask based solution as its reconfiguration speed requisites are not very demanding. For the Wireless Flow Scheduler an IP/socket based solutions is being developed and tested in order to achieve the highest possible speed.

In addition, some software modifications addressing the possibility of transmitting this configuration information from the AP to the STAs in the beacon are being implemented. In this way, the configuration information would be transmitted at the beginning of each super-frame, offering a higher reconfiguration speed compared to FTP/SSH based transmissions. The new information being sent in the beacon frame can be seen i[n Figure 2-2.](#page-11-0)

```
struct SHARP beacon frame data str
₹
    uint8 t DST ID;
   uint8 t ORIG ID;
   uint8 t reserved;
    uint8 t valid flag;
    uint32 t ID;uint32 t timestamp;
    uint32_t timestamp_s;
   struct superframe info sframe info;
} attribute ((packed));
struct superframe info {
    uint16 t TSN cycle;
   uint8_t STD_period;
   uint8_t DL_start;
   uint8_t N_DL_flows;
    uint8_t N_UL_flows;
    struct Sharp Flow DL_Flows[MAX_N_DL_flows];
    struct Sharp Flow UL Flows[MAX N UL flows];
\} _attribute_((packed));;
 struct Sharp_Flow{
    char TSN_DST[12];
    char TSN_ORIG[12];
    uint8_t_TSN_PCP;
    uint8_t_TSN_VLAN_ID;
    char w_SHARP_header[8];
    uint8_t_length;
    uint8_t MCS;
   uint8 t frame_type;
   uint8_t needs_ACK;
   uint8_t STA;
} attribute ((packed));;
```
Figure 2-2: New beacon data structure, adding the super-frame configuration.

Regarding the provision of statistics/telemetry/capabilities to other devices in the TIMING architecture, some changes are being done in the control software of both processors of the Wi-Fi TSN Node. Most of the statistics are collected in the RT processor so it has been necessary to enable the transmission of this information from this processor to the general processor by using the inter-processor communication module (remoteproc) as shown in [Figure 2-3.](#page-12-0) Once transmitted to the general-purpose processor, they are compiled in a file so that they can be consumed from the API for the Wireless Scheduler and the SDN-TSN Controller.

```
static void SHARP_AP_UL_Data_Rx(void *pvParameters)
      BaseType t xOueue ret:
      udserype_c xquede_rec;<br>uint32_t xQUEUE_buffer_input;<br>struct SHARP_irq_data* SHARP_irq_data_ptr;
      wint32_t* Desc_src_ptr;<br>uint32_t* Desc_src_ptr;<br>uint32_t* Desc_src_ptr;<br>uint3 _t DL_STA_Rx_end_ptr_alloc[8U+DESC_RX_END_BYTES*4U];<br>struct DL_STA_Rx_end_str* DL_STA_Rx_end_ptr;
      static int aux=0;
      DL_STA_Rx_end_ptr = (struct DL_STA_Rx_end_str*) DL_STA_Rx_end_ptr_alloc;
      uint32 t \log[4];
      for (;;)
      ł
           xQueue_ret = xQueueReceive(xQueue_AP_UL_Data_Rx, (uint32_t*)&xQUEUE_buffer_input, portMAX_DELAY);
           if (xQueue net == pdTRUE)SHARP_irq_data_ptr = (struct SHARP_irq_data*) &xQUEUE_buffer_input;
                 Desc_src_ptr = (uint32_t*)(SHARP_DESC_WR_MEM_BASEADDR+SHARP_irq_data_ptr->Desc_end_ADDR);
                 \frac{1}{2} uint32_t No_bytes = 312;
                 SHARP_memcpy_32((uint32_t*)DL_STA_Rx_end_ptr,Desc_src_ptr,No_bytes);
                 if (DL STA Rx end ptr->Desc Rx end.status == 0){
                      ^{\sim}aux++:
                      if (aux==20000){
                            enum rproc_msg_types_e rproc_msg_types = SET_LOG;
                            enum rproc_msg_types_e rproc_msg_set_log_types_e = LOG_RSSI;
                           log[0] = rproc_msg_types;<br>log[1] = rproc_msg_set_log_types_e;<br>log[2] = DL_STA_Rx_end_ptr->Desc_Rx_end.RSSI;
                           log[3] = DL_STA_Rxend_ptr - \triangle Psc_Rxend.MCS;taskENTER_CRITICAL();
                           lept.dest_addr = 0x00000400;
                           rpmsg_send(&lept, log, 16);
                           taskEXIT_CRITICAL();
                           aux=0;
                      \overline{\mathcal{X}}\overline{\phantom{a}}\rightarrow\,\mathcal{F}
```
Figure 2-3: First version of the task of transmitting statistics from the RT processor to the general task processor.

2.1.2.2 Stand-alone preliminary tests performed

In order to evaluate the current reconfiguration speed, the time between the application of the new configuration to an AP and the reconfiguration of both AP and STA has been measured.

The measured time is around 20 seconds.

Currently the run-time reconfiguration capability of the Wi-Fi TSN Node is limited due to the model based in configuration files and transmission of the configuration via FTP/SSH protocols between AP and STA. Most of the time is taken applying the new configuration file, due to the need to parse the file from the general-purpose processor to the RT processor, and from there, to the HW.

2.2 SCHEDULING ALGORITHMS FOR WI-FI TSN

2.2.1 Description

The Wi-Fi scheduler (WFS) interacts with the Wi-Fi TSN node at regular superframe intervals. On one hand, the Wi-Fi TSN node reports details about the active flows, including comprehensive information such as traffic characteristics, priority type, direction (uplink or downlink), queue size for asynchronous traffic and bit rate for isochronous traffic. Additionally, the Wi-Fi TSN node provides environmental information, particularly regarding the state of the channel. Provided with this information, the scheduler is tasked with making decisions about the structure of Wi-Fi frames. The global scheduler algorithm, integrated into the TSN Wi-Fi node, encompasses several modules, as depicted i[n Figure 2-4](#page-14-1) and summarized below:

- TSN Wi-Fi Windows Designer. At the TSN Wi-Fi node, each isochronous TSN flow is allocated into a periodic time window by the connectivity manager, considering the required throughput and the worst-case channel state. The WFS, depending on the channel state, optimizes the pre-assigned time window in the Wi-Fi link, reducing the size of the time window assigned to the isochronous flows if the channel state allows it and thus freeing slots that can be reused by other non-isochronous and lower priority flows. In this sense, the advanced window designs described in D1.4 (SP1) will not be implemented in SP2 for the sake of compatibility with the Wi-Fi node hardware and the regular TSN architecture adopted in Timing, instead, the TSN traffic corresponding to a single flow is transmitted in the pre-allocated time window at the maximum speed allowed by the channel, which is equivalent to reducing the pre-allocated window to the here named Wi-Fi window.
- DL/UL Splitter for UL and DL distribution of the slots not devoted to isochronous traffic. It computes dynamically the number of slots assigned to DL and UL by applying a reinforcement learning based algorithm which takes decisions depending on the aggregated DL and UL queues sizes.
- DL Asynchronous Traffic Scheduler: Determines the flows served in the DL prioritizing first the asynchronous traffic with quality of service over the BE traffic without QoS.
- UL Asynchronous Traffic Scheduler: Analogously to the DL scheduler, it decides the flows to be served in the uplink considering the established priorities.
- Frame Builder: It builds the frame based on the designed TSN Wi-Fi windows for isochronous traffic and on the DL&UL scheduler decisions.

The flow chart that completely encompasses the scheduler is shown in [Figure 2-4.](#page-14-1)

Figure 2-4: WiFi Scheduler (WFS) flowchart

[Figure 2-5](#page-14-2) shows an example in which in the downlink part of the frame, flows have been allocated into the different slots. The Wi-Fi time windows assigned to isochronous traffic are symbolized in yellow.

Figure 2-5: Frame Design Example

2.2.2 APIs design to integrate the scheduler in the Wi-Fi Node

In the API design to integrate the WFS into the TSN Wi-Fi Node some steps have to be defined or planned previously, as listed below.

- 1. API Design: To integrate the WFS into the TSN node, an API will be designed to manage the interface between both. Progress in programming the API will be carried out in three successive stages.
	- a. In the first stage a socket in Matlab will be programmed. In the first integration test a basic configuration and communications block will work in Matlab

> dynamically providing a frame design every 5 sec. approximately, which is determined by the limitations of the Matlab socket.

- b. In the second stage, the same development will be carried out but programming in C, expecting to obtain a higher periodicity in the frame structure updating process.
- c. In the most advanced level (third stage), the C-programmed API will be compiled into the Linux program, embedding the API locally within the SHARP Wi-Fi TSN card. This approach is intended to achieve the fastest speed in the periodicity of frame design.

The parameters that will be exchanged through the implemented API are those given in tables I and II in deliverable SP-1 D1.3, specifically in Section 2.3.1.

2. Wi-Fi time window design for isochronous traffic: Regarding the isochronous flows, the connectivity manager will orchestrate the route as well as the time window assigned for its transmission in each of the network elements. Based on this design, the TSN Wi-Fi node will inform the WFS of the periodicity and the start and end time stamps of the assigned window, initially designed to support the lowest MCS (channel state worstcase). In the initial phase of the scheduling algorithm, if the channel state is better than the worst-case scenario, the MCS will be increased and the time windows for the Wi-Fi segment will be adjusted. This involves allocating the corresponding flow into a smaller time interval than the initial one, and as a result it will be possible to opportunistically accommodate lower-priority and asynchronous traffic in the free time intervals. The WFS will define the assigned time interval in slots (OFDM symbols), serving as the temporal granularity unit within the Wi-Fi segment.

In any case, it must be guaranteed that the transmission of the assigned flows, in each bridge, does not exceed the time windows assigned by the connectivity manager. If necessary, while Wi-Fi windows are prepared, the information can be buffered in the STA devices (UL) or the AP (DL) without violating the previous condition.

- 3. Real-Time DL/UL split: The Real-Time DL/UL splitter will compute the size of the aggregated DL queue, i.e. the summation of the sizes of all the asynchronous DL queues (traffic with service quality and BE if there are), considering different levels of priority and the size of the aggregated UL queue. Then it will also compute the number of free slots in the frame, once isochronous flows have been allocated in the frame, and by applying a RL based decision it will deliver the number of slots allocated for asynchronous traffic in RT-DL and the number of slots allocated for asynchronous traffic in RT-UL. Alternatively, it will decide the frontier of the RT-DL and RT-UL subframes.
- 4. Scheduler for Real-Time DL frame [DOWNLINK]: The Real-Time DL scheduler allocates slots among the Downlink (DL) flows associated with quality of service using a largestweighted-delayed-first scheduler (LWDF) algorithm. This algorithm leverages channel state information and buffers occupation to determine the number and specific slots assigned to each flow. If there are remaining free slots, they are then distributed among Best Effort (BE) flows using a straightforward Round-Robin algorithm.

- 5. Scheduler for Real-Time UL frame [UPLINK]: In UL, the same scheduler is used, as in DL, but in this case a guard time and a preamble must be inserted before the slots corresponding to the different users or STAs.
- 6. Channel modeling: The channel model for scheduler design and simulation has to be validated providing models for pathloss in terms of distances, pathloss exponent and average received power and for fading profile in terms of PDP, Delay spread, Rayleigh or Rice Line of Sight, Doppler spread and coherence time.
- 7. SHARP Frame definition (RT subframe): RT frame definition for scheduler design and simulation has to be validated, providing values for bandwidth, OFDM symbol length, slot time, number of carriers for data, MCS modes, inter-symbol spacing, and PHY&MAC overhead in UL and in DL.

Model	Parameter	Value
Channel Pathloss	Distance	TBD
	Exponent	TBD
	Average received power	TBD
Multipath Fading	Delay Spread (rms)	89 ns (Ch1) - 29 ns (Ch2)
	NLOS	Yes/No
	Doppler spectrum	TBD
	Coherence Time	30 msec.
Frame Structure	Bandwidth	20 MHz
	OFDM symbol length	$4 s (0.8 s + 3.2 s)$
	Slot Time	OFDM symbol
	Number of data subcarriers	48
	MCS	07
	PHY overhead (DL)	20 _s
	PHY overhead (UL)	4 _s
	MAC overhead (DL)	TBD
	MAC overhead (UL)	TBD
	IFS (Interframe Spacing)	TBD
	Table 2 shows the different parameters needed for both integration and simulation purposes.	
Block	Inputs	Outputs

[Table 1](#page-18-0) shows the input/output variables associated to the different blocks that make up the scheduler an[d](#page-18-1)

Table 1: Input and output variables at the WFS functional blocks

Table 2: Parameter definition for simulation and integration purposes

2.2.3 Report on the Current Status

As explained in D1.3 (SP-1) a simplified version of the WFS component has been programmed in Matlab in order to carry out the first WFS integration test scheduled for the end of February 2024. Recently, some functionalities regarding blocks 1 and 2 i[n Table 1](#page-18-0) have been updated and integrated in the Matlab script. Moreover, blocks 3 and 4 are now only concerned with asynchronous traffic since block 1 is responsible for isochronous traffic allocation. This Matlab script will be used to build the first step of the API to intercommunicate the WFS with the TSN Wi-Fi node as reported in section 2.2.2. [Table 3](#page-19-2) shows the designed Matlab functions devoted to the different WFS blocks.

Table 3: Matlab functions included in the script for first integration test

2.2.4 Report on the stand-alone preliminary tests

The stand-alone preliminary test performed to validate the implementation of the simplified WFS that will be integrated in the first integration test was described in SP-1 D1.3, section 2.3.3. For example, the test was executed in a scenario including 5 DL flows and 3 UL flows (1 packet per flow) of different classes (ISO, ASYN and BE) and were scheduled in one superframe.

The input variables definition was as follows:


```
Nflows=Nusers; % number of flows (we consider 1 flow per user)
Nbuffers=Nflows; % number of buffers
PacketLength=1500; % fixed packet length in bits
BufferSize=1000; % buffer size in packets
SystemPars.NumSlots=73; % number of slot per superframe (DL+UL) (INPUT)
DataCarriers=234; % from 802.11ax (Table 27-79)
OFDMSymbolperSlot=1; % 1 slot = 1 OFDM symbol
SystemPars.ChannelUsesperSlot=DataCarriers*OFDMSymbolperSlot; % channel use = 
DataCarriers*OFDMSymbolsperSlot (INPUT)
for n=1:Nbuffers
     buffer_struct(n).BufferSize=BufferSize; %% Number of elements in the 
buffer
    buffer\_struct(n).LastBufferElement = 1; % Position of the last
element in the buffer
    buffer_struct(n).FirstBufferElement = 1; %% Position of the first
element in the buffer
     buffer_struct(n).NumPackets= 1; %% # of packets in buffer
    buffer struct(n).PendingBits = PacketLength*buffer struct(n).NumPackets;
%% Bits pending to be sent (INPUT)
     for jj=1:BufferSize 
         buffer_struct(n).Buffer(jj).SizeOfElement = PacketLength; % Number of 
bits of a buffer element (packets size)
         buffer_struct(n).Buffer(jj).BitsYetTransmitted = 0; % Bits 
transmitted in previous frames
        buffer struct(n).Buffer(jj).GenerationTime = 0; \% Packet
arrival time in seconds (INPUT: not used yet)
     end
end
% DL FLOWS
% Class 1 flows (ISO)
FlowList(1).Direction=1; % 1: DL; 2: UL (INPUT)
FlowList(1).ConnexionIdentifier.MS=1; % STA generating/receiving the flow 
(INPUT)
FlowList(1).ServiceClassIdentifier=1; % Flow class (INPUT)
FlowList(1).FlowBuffer=buffer_struct(1); % Buffer storing the packets of this
flow
FlowList(1).Periodicity=1; % flow periodicity in superframes (integer number) 
(INPUT)
FlowList(1).FirstSlotTXWindow=10; % Slot of the DL frame at which 
transmission starts (INPUT)
% Class 2 flows (ASYN)
for n=2:3
     FlowList(n).Direction=1; 
     FlowList(n).ConnexionIdentifier.MS=n; 
     FlowList(n).ServiceClassIdentifier=2;
     FlowList(n).FlowBuffer=buffer_struct(n);
end
% Class 3 flows (BE)
for n=4:5
     FlowList(n).Direction=1;
```

```
 FlowList(n).ConnexionIdentifier.MS=n; 
     FlowList(n).ServiceClassIdentifier=3;
    FlowList(n).FlowBuffer=buffer struct(n);
end
% UL FLOWS
% Class 1 flows (ISO)
FlowList(6).Direction=2; 
FlowList(6).ConnexionIdentifier.MS=6; 
FlowList(6).ServiceClassIdentifier=1; 
FlowList(6).FlowBuffer=buffer struct(6);
FlowList(6).Periodicity=1; 
FlowList(6).FirstSlotTXWindow=65; 
% Class 2 flows (ASYN)
FlowList(7).Direction=2;
FlowList(7).ConnexionIdentifier.MS=7; 
FlowList(7).ServiceClassIdentifier=2;
FlowList(7).FlowBuffer=buffer_struct(7);
% Class 3 flows (BE)
FlowList(8).Direction=2;
FlowList(8).ConnexionIdentifier.MS=8; 
FlowList(8).ServiceClassIdentifier=3;
FlowList(8).FlowBuffer=buffer struct(8);
% STA Capacity (bits per channel use = bps/Hz = MCS spectral efficiency)
MS(1).Capacity=6*3/4; % MCS=6 (WiFi6) --> 1053 bits/slot
MS(2).Capacity=2*3/4; % MCS=2 (WiFi6) --> 351 bits/slot
MS(3).Capacity=8*5/6; % MCS=9 (WiFi6) 
MS(4).Capacity=1*1/2; % MCS=0 (WiFi6) --> 117 bits/slot
MS(5).Capacity=8*5/6; % MCS=9 (WiFi6) --> 1560 bits/slot
MS(6).Capacity=2*3/4; % MCS=2 (WiFi6) 
MS(7).Capacity=2*3/4; % MCS=2 (WiFi6)
MS(8).Capacity=1*1/2; % MCS=0 (WiFi6)
% ISO traffic
% ServiceClass(1).Name='ISO';
ServiceClass(1).ToleratedJitter=1e-6; % TBD (INPUT: not used yet)
ServiceClass(1).MaximumLatency=1e-3; % TBD (INPUT: not used yet)
ServiceClass(1).MaximumPacketErrorRatio=1e-4; % TBD (INPUT: not used yet)
% ASYN traffic
% ServiceClass(2).Name='ASYN';
ServiceClass(2).ToleratedJitter=1e-4; 
ServiceClass(2).MinimumRate=0; % TBD (INPUT: not used yet)
ServiceClass(2).MaximumLatency=1e-3; 
ServiceClass(2).MaximumPacketErrorRatio=1e-4; 
% BE traffic
% ServiceClass(3).Name='BE';
ServiceClass(3).ToleratedJitter=inf; % TBD (INPUT: not used yet)
ServiceClass(3).MinimumRate=0; % TBD (INPUT: not used yet)
ServiceClass(3).MaximumLatency=inf; % TBD (INPUT: not used yet)
ServiceClass(3).MaximumPacketErrorRatio=1e-3; % TBD (INPUT: not used yet)
```


Obtaining a superframe variable structured as follows:

The length of the SuperFrame.Flow array is the superframe duration in slots. In the example, the superframe is split into 46 slots for the DL (in black) and 27 slots for the UL (in blue). SuperFrame.Flow(3) =2 means that the third slot is assigned to the flow with identifier 2. Isochronous windows are marked in yellow.

2.3 ETHERNET TSN NODES INCLUDING SOFTWARE

2.3.1 Report on the current status

2.3.1.1 The Southbound API between the CNC and the TSN switches

The WRZ-API stands as a specialized interface facilitating communication between the Northbound API and Safran's TSN switches. This API, integrated within the equipment, plays a pivotal role in transmitting network configurations from the CNC to the respective equipment.

We will now delve into the format of the endpoints developed.

2.3.1.1.1 Configuration endpoints

These endpoints are related to the configuration of the TSN network and are categorized into four types: VLAN configuration, TAS configuration, CAM configuration, and QCI configuration. Each of these types encompasses an endpoint through which new configurations can be applied.

- Notes:
	- vlan_prio is an integer between 0-7.
	- The Z16s are divided into 2 quads.
		- For the first quad (wr0-wr3): **dest** and **red_dest** go from 0 to 3.

- For the second quad (wr4-wr7): **dest** and **red_dest** go from 0 to 3.
	- o Example: wr4 is 0 (it is the first element from the second quad)
- Values for parameter **ds** are (used in combination with the protocol *port* number):
	- 1: source port
	- 0: destination port
- Values for parameter **redundant** must be 1/0.
- red_handle is an integer between 0-16. Must be unique per flow.
- vlan_id_cfg and vlan_prio_cfg are unused.

- Notes:
	- **interval_time** is expressed in 16 ns clock cycles. This field would be the result of the following operation: Ceil [T_int (ns) /16].
	- **gate_cfg** is the configuration for each of the TAS queues in one-hot encoding. Example: 0x9 -> 1001 (Q3 and Q0 active).
	- **tv_sec** is the Linux UTC time in which TAS configuration is applied.
	- **tick_granularity** is the size of the step for the interval_time. The default value is 0 which corresponds to 16 ns and it is the recommended value. The other possible values are 1,2,3, and 4, which correspond to 2,3,5 and 9 clock cycles.
	- **framePreemptionStatusTable** indicates the queues that are preemptable in one-hot encoding. Example: 0x8 -> 1000 (Q3 is preemtable). By default, all the queues are express.

- Notes:
	- **vlan_pcp** is the priority and it is an integer between 0-7.
	- The Z16s are divided into 2 quads.
		- For the first quad (wr0-wr3): **rx_dest** goes from 0 to 3.
		- For the second quad (wr4-wr7): **rx_dest** goes from 0 to 3.
			- o Example: wr4 is 0 (it is the first element from the second quad)

- Notes:
	- **EXECT** interval_time is expressed in 16 ns clock cycles. This field would be the result of the following operation: Ceil [T_int (ns) /16].
	- **Example 2 reads r gate_cfg** is the configuration for each of the QCI queues in one-hot encoding. Example: 0x9 -> 1001 (Q3 and Q0 active).
	- **tv_sec** is the Linux UTC time in which QCI configuration is applied.
	- **qci_tick_granularity** is the size of the step for the interval_time. The default value is 0 which corresponds to 16 ns and it is the recommended value. The other possible values are 1,2,3, and 4, which correspond to 2,3,5 and 9 clock cycles.

2.3.1.1.2 Generating reports endpoints

These endpoints aim to generate data about packet traffic within the Z16.

• Notes:

o **frer** must be 1/0 and it is the flag to compensate the offset value to get the sequence number if redundancy headers are present.

2.3.1.2 The Northbound API between the CNC and the TSN Controller

The Northbound API between the CNC and the TSN Controller implements a generic REST API for receiving the optimized operational settings calculated at UPC's TSN Controller.

2.3.1.2.1 Network equipment management

These are the endpoints responsible for managing the settings of the Z16 nodes stored within the CNC's database.

2.3.1.2.2 Network configuration

These endpoints are used to configure the TSN network.

VLAN

▪ **vlan_id_cfg** and **vlan_prio_cfg** are unused.

TAS

CAM

QCI

2.3.1.2.3 Latency and packet loss test generation

These endpoints are used to generate reports on packet traffic in the TSN network.

• Notes:

o **frer** must be 1/0 and it is the flag to compensate the offset value to get the sequence number if redundancy headers are present.

2.3.1.2.4 Network topology visualization

This endpoint returns a dictionary containing information about the connections of the devices.

2.3.1.2.5 Export Metrics Settings

This endpoint allows for obtaining a list of monitored devices, as well as adding or removing devices.

2.3.1.3 Workflow

In this section we will explain the workflow designed between the TSN Controller, CNC, and the Z16 devices for monitoring and configuring the TSN network. [Figure](#page-35-0) 2-6 illustrates the flow of information exchange among these three components. The downward arrows symbolize the process of network configuration, initiated by the TSN Controller, translated through the CNC API to the WRZ API, and subsequently distributed to the respective devices via the WRZ API. Conversely, the upward arrows signify the network monitoring process, where information is extracted from the devices using the SNMP Exporter, managed by the CNC, and eventually relayed to the TSN Controller.

Figure 2-6: Ethernet TSN workflow

Configure the TSN Network

The steps to be taken to configure the TSN Network are (downward arrows):

Preliminary steps (0.1 and 0.2):

0.1. The CNC API is used to add the devices that are on the TSN network.

Example:

0.2. Configuration rules are created using the CNC

Example – Create TAS rule:

If we want the following GCL specification:

■ Slot #0: Duration 1 ms, Gate Settings: Q3 & Q0 open.

The parameters we should put are:

▪ **interval_time:** result of the operation Ceil [T_int (ns) /16]

- □ In this case, Ceil $[1000000 / 16] = 62500$
- **e** gate_cfg: Q3 and Q0 open is 1001 -> 0x9

The request will be:


```
The rule created has id = 1.
```
Example - Create VLAN rule:

If we want the following VLAN configuration:

- WR0: VID: 1; PCP: 2; MAC_DST: 0xcafebabe0102; MAC_ADDR: 0xaabbccddeeff; HAS_DEST:1; DEST:1.
- WR1: VID: 1; PCP: 2; MAC_DST: 0xcafebabe0102; MAC_ADDR: 0xaabbccddeeff.

The parameters we should put are:

- wr0:
	- "addr": "0xcafebabe0102"
	- \blacksquare "vlan_id": 1,
	- "vlan_prio": 2,
	- "mac_addr": "0xaabbccddeeff",
	- □ "dest": 1,
	- "has_dest": 1
- $wr1$:
	- "addr": "0xcafebabe0102"
	- "vlan_id": 1,
	- "vlan_prio": 2,
	- "mac_addr": "0xaabbccddeeff",

The requests will be:

• Rule for wr0

{"id":1, "vlan_id":1, "vlan_prio":2, "filter_en":0, "mac_addr":"0xaabbccddeeff", "proto":0, "ip":0, "vlan_id_cfg":0, "port": 0,"vlan_prio_cfg": 0, "dscp": 0,"dest": 1,"has_dest": 1,"ds": 0,"redundant": 0,"red_dest": 0,"red_handle": 0}

• Rule for wr1

wr0's rule has **id = 1** and wr1's rule has **id = 2**.

1. TSN Controller decides which network configuration is going to be used and the corresponding rules are applied with the CNC API.

Example – Apply TAS rule:

In rules we put the *id* of the rule we have previously created.

Parameter **tv_sec** is the Linux UTC time in which TAS configuration is applied and should be calculated in the following way:

■ <Current Network Epoch> + Additional offset to make up for configuration time.

Example – Apply VLAN rule:

In rules we put the *id* of the rule we have previously created.

The following steps are done automatically, without user control.

2. The CNC API makes a call to the WRZ API of the corresponding device to apply the network configuration.

3. WRZ API applies the chosen configuration to the devices.

Monitor the TSN Network

To monitor the TSN Network the workflow is:

1. SNMP collects the metrics defined in the CNC.

2. CNC saves the collected metrics and exports them to an external database so that TSN Controller can access them.

3. TSN Controller analyses the collected data and, based on it, decides which network configuration to apply.

2.4 SDN-TSN CONTROLLER

2.4.1 Description

The Software Defined Networking (SDN)- Time-Sensitive Networking (TSN) controller (SDN-TSN) is the entity responsible for the configuration of connectivity services within the TSN-based network infrastructures, and the corresponding management of said infrastructure. To do so, it engages on one hand with the Connectivity Manager (CM), which is the entity that receives service provisioning requests, which ultimately translate to connectivity requests towards the SDN-TSN controller; and, on the other hand, it engages with the underlying TSN-capable network nodes (WiFi and Ethernet) at the data plane, to configure them according to the

characteristics of the requested connectivity. [Figure 2-7](#page-39-0) depicts a schematic of the functional architecture of the module and the interactions with other modules, namely, the CM and the TSN-enabled data plane, as well as the main data structures treated within each of the individual functionalities. The communications are achieved thanks to a pair of interfaces: the Northbound Interface (NBI) and the Southbound Interface (SBI).

Figure 2-7: SDN-TSN Controller architecture.

The core of the functionalities is as listed below:

- Topology Manager (TM): responsible for keeping a graph representation of the underlying data plane, the network resources and their capabilities. It extracts this information via the SBI.
- Inventory Manager (IM): contains the inventory of the established data flows, according to the outputs of the connectivity provisioning operations performed at the Provisioning Manager. It exposes this information to the CM via the Provisioning Manager through the NBI.
- Provisioning Manager (PM): centralizes the provisioning of data flows to satisfy the requirements of connectivity services coming from the NBI. The low level configurations are decided thanks to the information stored at the Topology and Inventory Managers, and executes the computed configuration by means of the SBI.
- Monitoring Manager (MM): it collects the current status of the resources at the TSNcapable data plane as well as potential alarms via the SBI. The collected information is exposed to the CM thanks to the NBI.

2.4.2 Report on the Current Status

The advancements made for the current version of the SDN-TSN are mainly in two fronts: on one hand, the core functionalities of both the provisioning and the topology managers have been further developed to include the main classes, data structures and logic for their operation; on the other hand, both NBI and SBI have been expanded and validated for the operations of flow creation and deletion according to connectivity services requests coming from the CM (NBI) and the de-configuration of data plane nodes as a response of a flow deletion (SBI).

Starting with the NBI, [Figure 2-8](#page-40-0) depicts a schematic of the data structure of the flow of an application/service, and the relationship between sub-structures. Using such structure, the operations for the creation and deletion of flows have been defined[. Table 4](#page-41-0) an[d Table 5](#page-41-1) provide the details of the operations and the parameters of the requests.

Figure 2-8: Data structure tree of a service request.

Table 4: Description of the operation for creating a data flow

Table 5: Description of the operation for deleting a data flow

The remote access to the presented NBI operations has been validated with the tests reported below. The current tests expand what have been reported already in deliverable SP1 - D1.3, which focused on the flow and topology status dissemination. In addition, let us mention that the time from flow creation/deletion request to data plane configuration/de-configuration operation, that is, the time that the SDN-TSN requires for processing upper layer requests, transform them into specific configurations and issue them towards the data plane via the SBI, has been measured and is in the order of several tens to few hundreds of milliseconds, depending on the current load of the server in which the SDN-TSN controller is running. With this being said, the tests are reported below, for which we showcase Wireshark captures of the successful exchanges between a test REST client and the NBI of the controller.

Test #002

Test Id TSN-C_NBI_deleteFlow

Description

Test related to the Northbound Interface (NBI) of the SDN-TSN controller to verify the correct reception and processing of data deletion creation requests.

Expected Results

The NBI accepts the operation and processes the body of the request.

Output

The operation is successful (200 OK)

 $\overline{17}$ ntroller/flows/c384c16 172.26.3...HT... 179 DELETE tsnContro
172.26.3...HT... 60 HTTP/1.0 200 OK 39 172 26 37 218 rame 312: 179 bytes on wire (1432 bits), 179 bytes captured (1432 bits) on interface any, id 0 inux cooked capture v1 nternet Protocol Version 4, Src: 172.26.37.218 , Dst: 172.26.37.218
|ransmission Control Protocol, Src Port: 52080, Dst Port: 16485, Seq: 1, Ack: 1, Len: 111

In regards to the SBI, the SDN-TSN Controller implements the API reported in Section [2.3.1.2](#page-26-0) of this same document. For the sake of conciseness, we refer to the description of the operations and parameters already described. The SBI has been further developed to consider an expanded set of operation with respect what was previously reported in SP1 - D1.3, essentially, the configuration of the CAM as well as the deletion of VLAN, TAS and CAM configurations. In below

we report several experiments that validate the implemented operations. To this end, a mockup REST server emulating the NBI of CNC has been employed so as to validate that the SDN-TSN Controller issues properly the operation, thus, the SBI works as expected.

Frame 5704: 194 bytes on wire (1552 bits), 194 bytes captured (1552 bits) on interface any, id 0 Linux cooked capture v1
Internet Protocol Version 4, Src: 172.26.37.218, Dst: 172.26.37.218
Iransmission Control Protocol, Sr

2.5 CONNECTIVITY MANAGER

2.5.1 Description

The Time-Sensitive Networking (TSN) Connectivity Manager (CM) is a key component in the TIMING architecture, playing as the cornerstone role in managing and orchestrating timesensitive networks. In the context of the TIMING project, the TSN CM leverages the rest of the TIMING control plane component to ensure efficient, reliable, and timely data transport across the network. The TSN CM is strategically positioned at the top of the TIMING architecture as depicted in [Figure 2-9.](#page-45-0) It is designed to interact directly with all TSN and Metro SDN (Software-Defined Networking) Controllers, as well as the Digital Twin (DT). This central placement allows the TSN CM to have an overarching view of the network, enabling it to perform its functions effectively.

Figure 2-9: TSN Connectivity Manager Architecture

To fulfil its role, the TSN connectivity manager does so by leveraging a set of functionalities tailored to manage and optimize time-sensitive data flows across the network, which are summarized below.

- *Network Resource Discovery and Management:* One of the key functionalities of the TSN CM is to discover the resources and capacities available within the TSN network. This involves identifying available bandwidth, latency characteristics, and other critical network parameters.
- *Configuration and Optimization:* Post discovery, the TSN CM summarizes the main changes needed in the network and communicates these to the underlying controllers. This process involves deploying new configurations to optimize network performance, especially for time-sensitive data flows.
- *Integration with Digital Twin for Predictive Analysis:* The TSN CM collaborates with the Digital Twin (DT) to simulate network conditions and predict the performance of various network configurations. This predictive modeling is essential for making informed decisions about network configuration changes.
- *Optimization Algorithms:* Utilizing advanced algorithms such as the Dijkstra algorithm, the TSN CM efficiently manages the routing of data, ensuring the most optimal paths are chosen for time-sensitive traffic.

The TSN CM utilizes a range of interfaces to interact with different components of the TIMING architecture. The TSN CM makes use of three main APIs: Southbound Interface (SBI), Northbound Interface (NBI), and East/West Interface (E/WBI).

- *SouthBound Interface (SBI)*: This interface is used for reading underlying topological information, which is crucial for understanding the current state of the network.
- *NorthBound Interface (NBI)*: It facilitates interactions with other TIMING modules and higher-level applications. This interface is vital for global and automated management of the architecture.
- *East/WestBound Interface (E/WBI)*: This interface enables the TSN CM to communicate with the Digital Twin, sending and receiving information necessary for simulating network conditions.

2.5.2 Report on the Current Status

The overall progress of the development of the TSN Connectivity Manager is on-track and with solid advances, specially, in the definition of internal models and in the Core and ENP functional blocks, as presented in D1.4 (SP1) and in this document, see later sections.

For future steps, it is planned to develop the external client-side modules (SBI and E/WBI) to allow the expected interaction with the controllers (via SBI) and with the Digital Twin (E/WBI). In addition, integration tests will be performed between all components involved. Potentially, as a consequence of the integration phase, some refinements to the TSN CM data model may be carried out.

2.5.3 NorthBound Interface Implementation

One of the major advances of the TSN Connectivity Manager has been on the definition and implementation of the NorthBound interface, that allows vertical and upper-layer applications to communicate with it. The revised architecture of the Northbound Interface (NBI) adopts a modular design, partitioning various functions into separate, distinct endpoints. As depicted in **Error! Reference source not found.**, this architecture is comprised of three principal categories: controllers, flows, and topologies. Each category is engineered to fulfill a specific role within the broader context of network management, ensuring a more streamlined and efficient operational framework.

Figure 2-10: TSN Connectivity Manager NBI endpoints structure


```
 }
 },
 "bandwidth": 0,
 "latency": 0
}
OUTPUT:
{
 "flow_id": "string",
 "name": "string",
 "description": "string",
 "topology_id": "string",
 "hostA": "string",
 "portA": "string",
 "hostB": "string",
 "portB": "string",
 "switchA": "string",
 "switchB": "string",
   "path": [
     "string"
 ],
 "trafficMix": {
 "video": {
      "scale": 0,
       "pattern": {
 "period": 0,
 "unit": "string",
        "pattern": [
        ],
        "pattern_unit": "string"
      }
 },
 "gaming": {
       "scale": 0,
       "pattern": {
 "period": 0,
 "unit": "string",
        "pattern": [
 ],
 "pattern_unit": "string"
      }
 },
 "internet": {
      "scale": 0,
      "pattern": {
        "period": 0,
        "unit": "string",
        "pattern": [
 0
         ],
        "pattern_unit": "string"
      }
    }
 },
 "bandwidth": 0,
   "latency": 0
}
```



```
Request URL: GET https://{CM-IP}:8003/topologies/{ID}
Example schemas (JSON-based):
INPUT:
-
OUTPUT:
  {
    "topology_id": "string",
    "source_controller": "string",
    "network": [
     {
 "name": "string",
 "description": "string",
 "network_id": "string",
       "network_types": {
          "layer": "string",
          "type": "string" 
 },
 "node": [
 {
          "node_id": "string",
 "name": "string",
 "description": "string",
 "type": "string",
 "ports": {
            "P1": {
             "name": "string",
 "port_type": "string",
 "layer": "string",
 "speed": 0,
            "transmission_delay": 0
 }
          }
        }
 ],
 "link": [
 {
          "link_id": "string",
 "name": "string",
 "description": "string",
 "nodeA": "string",
 "portA": "string",
 "nodeB": "string",
 "portB": "string",
 "layer": "string",
          "capacity": 0
         }
 ],
 "l2_topology_attributes": {},
       "l3_topology_attributes": {}
     }
    ]
  }
```


2.5.4 Stand-alone preliminary tests

To facilitate the user interaction with the Connectivity Manager NBI, we have developed a graphical user interface (see [Figure 2-11](#page-59-0)**Error! Reference source not found.**) , publicly available, within the TIMING project scope, in the following URL: [https://nbi-connectivity-manager.e](https://nbi-connectivity-manager.e-lighthouse.com/docs)[lighthouse.com/docs.](https://nbi-connectivity-manager.e-lighthouse.com/docs)

TSN Connectivity Manager - NorthBound Interface API CONDING

This API manages the connection with higher-application level and GUI Contact - Website

Provisioning

/ flows / Returns a list with all the enrolled TSN flows GET	\checkmark
POST /flows/ Create a new fow	\checkmark
/flows/ Deletes all the registered flows DELETE	\checkmark
/flows/{flow_id} Returns updated information of a TSN flow by id GET	\checkmark
/flows/{flow_id} Delates an existing TSN flow by id DELETE	\checkmark
Controllers	∼
/controllers/ Return all registered controllers GET	$\check{~}$
/controllers/ PostControlars POST	\checkmark
DELETE /controllers/ Deletes all the registered controller by id	\checkmark
GET /controllers/tsn Returns the controllers with type tsn	\checkmark
GET /controllers/metro Returns the controllers with type matro	\checkmark
GET /controllers/{controller_id} Returns the parameters of a controller by id	\checkmark
/controllers/{controller_id} Deletes the registered controller by id DELETE	\checkmark
Topologies	∼
/topologies/all Return all the topologies GET	\checkmark
/topologies/tsn Ratum all TSN topologies GET	\checkmark
/topologies/metro Ratum all Matro topologies GET	\checkmark
GET /topologies/{topology_id} Ratum a topology by id	\checkmark

Figure 2-11: TSN Connectivity Manager NBI GUI

In addition, to validate the remote access to the functionalities provided by this API, a wide range of tests have been executed, classified according to their functional group. These tests have been performed using the POSTMAN software from a personal computer remotely located from the server on which the TSN CM is deployed.

• *Controllers*:

 \sim


```
Accept: */*
Postman-Token: 4fec6d7c-62c4-4b9b-971e-ec788614bfe9
Host: nbi-connectivity-manager.e-lighthouse.com
Accept-Encoding: gzip, deflate, br
Connection: keep-alive
Content-Length: 184
{
"name": "tsn_controller",
"description": "tsn_controller",
"url": "https://192.0.22.1",
"port": 8080,
"username": "user",
"password": "pass123",
"type": "tsn"
}
RESPONSE
HTTP/1.1 200 OK
Content-Length: 202
Content-Type: application/json
Date: Fri, 12 Jan 2024 07:39:54 GMT
Server: uvicorn
{"controller_id":"42f98df2-b54f-495f-8f31-
5e76b91f1fef","name":"tsn_controller","description":"tsn
_controller","url":"https://192.0.22.1","port":8080,"use
rname":"user","password":"pass123","type":"tsn"}
Test 2: Register a Metro controller:
REQUEST
POST /controllers/ HTTP/1.1
Content-Type: application/json
User-Agent: PostmanRuntime/7.36.1
Accept: */*
Postman-Token: 4b04e30e-659c-4820-b369-4b485f970699
Host: nbi-connectivity-manager.e-lighthouse.com
Accept-Encoding: gzip, deflate, br
Connection: keep-alive
Content-Length: 188
{
"name": "metro_controller",
"description": "metro_controller",
"url": "https://192.0.46.1",
"port": 80,
"username": "user",
"password": "pass123",
"type": "metro"
}
RESPONSE
HTTP/1.1 200 OK
Content-Length: 206
Content-Type: application/json
Date: Fri, 12 Jan 2024 07:40:25 GMT
Server: uvicorn
```


• *Topologies*:

• Flows:

2.6 TSN MODELS/DT

2.6.1 Description

The TSN Digital Twin includes the following basic modules (see [Figure 2-12\)](#page-76-0):

- a manager module configuring and supervising the operation of the rest of the modules.
- a few modules that include algorithm, models, and the interface with the TSN connectivity manager.
- a Redis DB that is used in publish-subscribe mode to communicate the different modules among them.

Figure 2-12 Internal architecture of the Digital Twin

In this deliverable, we report the work that has been done from the status reported in TIMING SP-1 D1.3, especially regarding integration. In particular:

- Definition of common traffic profiles based on a set of defined service profiles and traffic patterns that will be adopted by other systems in the TIMING architecture, specifically in the TSN Connectivity Manager.
- Advances on the development of a preliminary version of the REST-API that now is integrated with the other internal modules of the DT.
- Development of a *Simulation Composer* module that creates the simulation scenario based on the topology and KPI evaluation request received.

2.6.2 Traffic Profiles

Two descriptors have been identified to define traffic profiles: service profiles and traffic patterns.

As for service profiles, they are described in terms of statistical distributions of the following random variables (see SP-1 D1.2):

- Inter-packet time
- Packet size
- Inter-burst time
- **Burst size**

Each random variable is defined then, as a probability distribution with some parameters, e.g., mean and standard deviation for the normal distribution, the units and a scale factor, usually 1. The following table reproduces an example for a video service.

"video":{ "interBurstRate": {"dist": "normal", "params": {"mean":0.25, "std":2.54e-5}, "unit":"s-1", "scale":1}, "burstSize": {"dist": "normal", "params": {"mean":4034485, "std":1266989}, "unit": "bytes", "scale":1}, "packetSize": { "dist":"constant", "params": {"mean":1500}, "unit": "bytes", "scale":1}, "interPacketTime": {"dist": "normal", "params": {"mean":0.0001, "std":0.00002}, "unit": "s", "scale": 1} }

Regarding traffic patterns, they are used to described the temporal behavior of a given service. They define, the periodicity and the specific pattern as values in the continuous range [0-1].

```
"daily": {"period": 86400, "unit": "s", "pattern": [0.25, 0.3, 0.32, 0.35, 0.39, .....], 
         "pattern_unit": "min"},
"periodical": {"period": 5, "unit": "ms", "pattern": [1,0,0,0,0], "pattern_unit": "ms"}
```
Finally, traffic profiles are defined as a mix of services with different traffic patterns and scale factors. The following table shows an example of traffic profile.

```
"trafficMix": {
         "video": {"scale": 100, "pattern": "daily-1"},
         "gaming": {"scale": 300, "pattern": "daily-2"},
         "internet": {"scale": 500, "pattern": "daily-3"}
}
```
2.6.3 REST API Interface

The current version of the DT includes a REST API which implements a preliminary interface. Specifically, this version of the REST API offers several endpoints that provide methods to interact with the Topology DB and to issue the evaluation of the KPI of a selected flow. The REST API offers the following endpoints:

- NetworkDB: Allows GET, POST and DELETE methods. Implements the creation, deletion and retrieval of the Topology DB in the Digital Twin.
- Topology/Node: Allows GET, POST and DELETE methods. Implements the creation, deletion and retrieval of the Node's information in the Topology DB.
- Topology/Link: Allows POST and DELETE methods. Implements the creation and deletion of the Link's information in the Topology DB.
- Flow: Allows GET, POST and DELETE methods. Implements the creation, deletion and deletion of the Flow's information in the Topology DB.
- KPIEvaluation: Allows GET method. Implements the KPI evaluation of a selected flow of the Topology DB.

The following table describes the end-points and operations and provides examples.

2.6.4 Report on Current Status

A preliminary REST API is already implemented and fully functional. In addition, this version of the DT includes preliminary algorithms to compute the delay in a set of queues that support a traffic flow. Initial versions of traffic definitions are also included.

2.6.5 Stand-alone preliminary tests

We have developed a REST API client that performs requests to the digital twin. The client can be configured by defining a JSON file with the requests to be submitted through the REST API interface and prints the results. Specifically, a JSON file named *Tests.json* has been defined that follows the following sequence:

- Create a network topology with 4 Hosts and 3 Switches in a ring.
- Add two flows as background traffic
- Request for RPI evaluation for a new flow

The results of the test are listed in the next table.

INFO: -- 200 OK - {'link': {'DPA': 'H1', 'PortA': 'E1', 'DPB': 'R1', 'PortB': 'A1'}, 'result': 'OK'} INFO: POST http://127.0.0.1:8080/Topology/Link/ {'DPA': 'H2', 'PortA': 'E1', 'DPB': 'R2', 'PortB': 'A1'} INFO: -- 200 OK - {'link': {'DPA': 'H2', 'PortA': 'E1', 'DPB': 'R2', 'PortB': 'A1'}, 'result': 'OK'} INFO: POST http://127.0.0.1:8080/Topology/Link/ {'DPA': 'R3', 'PortA': 'A1', 'DPB': 'H3', 'PortB': 'E1'} INFO: -- 200 OK - {'link': {'DPA': 'R3', 'PortA': 'A1', 'DPB': 'H3', 'PortB': 'E1'}, 'result': 'OK'} INFO: POST http://127.0.0.1:8080/Topology/Link/ {'DPA': 'R3', 'PortA': 'A2', 'DPB': 'H4', 'PortB': 'E1'} INFO: -- 200 OK - {'link': {'DPA': 'R3', 'PortA': 'A2', 'DPB': 'H4', 'PortB': 'E1'}, 'result': 'OK'} INFO: ---TOPOLOGY LINKS--- INFO: POST http://127.0.0.1:8080/Topology/Link/ {'DPA': 'R1', 'PortA': 'T1', 'DPB': 'R3', 'PortB': 'T1'} INFO: -- 200 OK - {'link': {'DPA': 'R1', 'PortA': 'T1', 'DPB': 'R3', 'PortB': 'T1'}, 'result': 'OK'} INFO: POST http://127.0.0.1:8080/Topology/Link/ {'DPA': 'R1', 'PortA': 'T2', 'DPB': 'R2', 'PortB': 'T1'} INFO: -- 200 OK - {'link': {'DPA': 'R1', 'PortA': 'T2', 'DPB': 'R2', 'PortB': 'T1'}, 'result': 'OK'} INFO: POST http://127.0.0.1:8080/Topology/Link/ {'DPA': 'R2', 'PortA': 'T2', 'DPB': 'R3', 'PortB': 'T2'} INFO: -- 200 OK - {'link': {'DPA': 'R2', 'PortA': 'T2', 'DPB': 'R3', 'PortB': 'T2'}, 'result': 'OK'} INFO: ---FLOWS-- INFO: POST http://127.0.0.1:8080/Flow/ {'flowId': 'BG_Traffic1', 'data': {'src': 'H1', 'srcPort': 'E1', 'dest': 'H3', 'destPort': 'E1', 'srcSwitch': 'R1', 'destSwitch': 'R3', 'path': ['R1', 'R3'], 'trafficMix': {'video': {'scale': 100, 'pattern': 'daily-1'}, 'gaming': {'scale': 300, 'pattern': 'daily-2'}, 'internet': {'scale': 500, 'pattern': 'daily-3'}}}} INFO: -- 201 Created - {'BG_Traffic1': "{'src': 'H1', 'srcPort': 'E1', 'dest': 'H3', 'destPort': 'E1', 'srcSwitch': 'R1', 'destSwitch': 'R3', 'path': ['R1', 'R3'], 'trafficMix': {'video': {'scale': 100, 'pattern': 'daily-1'}, 'gaming': {'scale': 300, 'pattern': 'daily-2'}, 'internet': {'scale': 500, 'pattern': 'daily-3'}}, 'flowId': 'BG_Traffic1'}", 'result': 'OK'} INFO: POST http://127.0.0.1:8080/Flow/ {'flowId': 'BG_Traffic2', 'data': {'src': 'H2', 'srcPort': 'E1', 'dest': 'H4', 'destPort': 'E1', 'srcSwitch': 'R2', 'destSwitch': 'R3', 'path': ['R2', 'R3'], 'trafficMix': {'video': {'scale': 100, 'pattern': 'daily-1'}, 'gaming': {'scale': 300, 'pattern': 'daily-2'}, 'internet': {'scale': 500, 'pattern': 'daily-3'}}}} INFO: -- 201 Created - {'BG_Traffic2': "{'src': 'H2', 'srcPort': 'E1', 'dest': 'H4', 'destPort': 'E1', 'srcSwitch': 'R2', 'destSwitch': 'R3', 'path': ['R2', 'R3'], 'trafficMix': {'video': {'scale': 100, 'pattern': 'daily-1'}, 'gaming': {'scale': 300, 'pattern': 'daily-2'}, 'internet': {'scale': 500, 'pattern': 'daily-3'}}, 'flowId': 'BG_Traffic2'}", 'result': 'OK'} INFO: ---GETs-- INFO: GET http://127.0.0.1:8080/NetworkDB/ INFO: -- 200 OK - {'nodeList': ['R1', 'R2', 'R3'], 'hostList': ['H1', 'H2', 'H3', 'H4'], 'linkList': [{'DPA': 'H1', 'PortA': 'E1', 'DPB': 'R1', 'PortB': 'A1'}, {'DPA': 'H2', 'PortA': 'E1', 'DPB': 'R2', 'PortB': 'A1'}, {'DPA': 'R3', 'PortA': 'A1', 'DPB': 'H3', 'PortB': 'E1'}, {'DPA': 'R3', 'PortA': 'A2', 'DPB': 'H4', 'PortB': 'E1'}, {'DPA': 'R1', 'PortA': 'T1', 'DPB': 'R3', 'PortB': 'T1'}, {'DPA': 'R1', 'PortA': 'T2', 'DPB': 'R2', 'PortB': 'T1'}, {'DPA': 'R2', 'PortA': 'T2', 'DPB': 'R3', 'PortB': 'T2'}], 'nodes': {'H1': {'DPId': 'H1', 'type': 'Host', 'ports': {'E1': {'DPId': 'H1', 'portId': 'E1', 'speed': 100, 'type': 'Host', 'neighbour': {'DPId': 'R1', 'portId': 'A1'}, 'inFlows': [], 'outFlows': ['BG_Traffic1']}}, 'Neighbours2Ports': {'R1': 'E1'}, 'flows': ['BG_Traffic1']}, 'H2': {'DPId': 'H2', 'type': 'Host', 'ports': {'E1': {'DPId': 'H2', 'portId': 'E1', 'speed': 100, 'type': 'Host', 'neighbour': {'DPId': 'R2', 'portId': 'A1'}, 'inFlows': [], 'outFlows': ['BG_Traffic2']}}, 'Neighbours2Ports': {'R2': 'E1'}, 'flows': ['BG_Traffic2']}, 'H3': {'DPId': 'H3', 'type': 'Host', 'ports': {'E1': {'DPId': 'H3', 'portId': 'E1', 'speed': 100, 'type': 'Host', 'neighbour': {'DPId': 'R3', 'portId': 'A1'}, 'inFlows': ['BG_Traffic1'], 'outFlows': []}}, 'Neighbours2Ports': {'R3': 'E1'}, 'flows': ['BG_Traffic1']}, 'H4': {'DPId': 'H4', 'type': 'Host', 'ports': {'E1': {'DPId': 'H4', 'portId': 'E1', 'speed': 100, 'type': 'Host', 'neighbour': {'DPId': 'R3', 'portId': 'A2'}, 'inFlows': ['BG_Traffic2'], 'outFlows': []}}, 'Neighbours2Ports': {'R3': 'E1'}, 'flows': ['BG_Traffic2']}, 'R1': {'DPId': 'R1', 'type': 'Switch', 'ports': {'A1': {'DPId': 'R1', 'portId': 'A1', 'speed': 100, 'type': 'Access', 'neighbour': {'DPId': 'H1', 'portId': 'E1'}, 'inFlows': ['BG_Traffic1'], 'outFlows': []}, 'A2': {'DPId': 'R1', 'portId': 'A2', 'speed': 100, 'type': 'Access', 'neighbour': None, 'inFlows': [], 'outFlows': []}, 'T1': {'DPId': 'R1', 'portId': 'T1', 'speed': 100, 'type': 'Trunk', 'neighbour': {'DPId': 'R3', 'portId': 'T1'}, 'inFlows': [], 'outFlows': ['BG_Traffic1'], 'transmissionDelay': 50.0}, 'T2': {'DPId': 'R1', 'portId': 'T2', 'speed': 100, 'type': 'Trunk', 'neighbour': {'DPId': 'R2', 'portId': 'T1'}, 'inFlows': [], 'outFlows': [], 'transmissionDelay': 50.0}}, 'Neighbours2Ports': {'H1': 'A1', 'R3': 'T1', 'R2': 'T2'}, 'flows': ['BG_Traffic1']}, 'R2': {'DPId': 'R2', 'type': 'Switch', 'ports': {'A1': {'DPId': 'R2', 'portId': 'A1', 'speed': 100, 'type': 'Access', 'neighbour': {'DPId': 'H2', 'portId': 'E1'}, 'inFlows': ['BG_Traffic2'], 'outFlows': []}, 'T1': {'DPId': 'R2', 'portId': 'T1', 'speed': 100, 'type': 'Trunk', 'neighbour': {'DPId': 'R1', 'portId': 'T2'}, 'inFlows': [], 'outFlows': [], 'transmissionDelay': 50.0}, 'T2': {'DPId': 'R2', 'portId': 'T2', 'speed': 100, 'type': 'Trunk', 'neighbour': {'DPId': 'R3', 'portId': 'T2'}, 'inFlows': [], 'outFlows': ['BG_Traffic2'], 'transmissionDelay': 50.0}}, 'Neighbours2Ports': {'H2': 'A1', 'R1': 'T1', 'R3': 'T2'}, 'flows': ['BG_Traffic2']}, 'R3': {'DPId': 'R3', 'type': 'Switch', 'ports': {'A1': {'DPId': 'R3', 'portId': 'A1', 'speed': 100, 'type': 'Access', 'neighbour': {'DPId': 'H3', 'portId': 'E1'}, 'inFlows': [], 'outFlows': ['BG_Traffic1']}, 'A2': {'DPId': 'R3', 'portId': 'A2', 'speed': 100, 'type': 'Access', 'neighbour': {'DPId': 'H4', 'portId': 'E1'}, 'inFlows': [], 'outFlows': ['BG_Traffic2']}, 'A3': {'DPId': 'R3', 'portId': 'A3', 'speed': 100, 'type': 'Access', 'neighbour': None, 'inFlows': [], 'outFlows': []}, 'T1': {'DPId': 'R3', 'portId': 'T1', 'speed': 100, 'type': 'Trunk', 'neighbour': {'DPId': 'R1', 'portId': 'T1'}, 'inFlows': ['BG_Traffic1'], 'outFlows': [], 'transmissionDelay': 50.0}, 'T2': {'DPId': 'R3', 'portId': 'T2', 'speed': 100, 'type': 'Trunk', 'neighbour': {'DPId': 'R2', 'portId': 'T2'}, 'inFlows': ['BG_Traffic2'], 'outFlows': [], 'transmissionDelay': 50.0}}, 'Neighbours2Ports': {'H3': 'A1', 'H4': 'A2', 'R1': 'T1', 'R2': 'T2'}, 'flows': ['BG_Traffic2',

2.7 METRO SDN CONTROLLER

The TSN domains are expected to be interconnected by means of a Metro Area Network, based on packet-switching technology, i.e., and IP network.

Two aspects should be considered:

- The specific network nodes forming the packet-switching infrastructure.
- The SDN controller used to configure the paths tailored to the characteristics required for satisfying the interconnection of the TSN domains. The goal is to leverage on the concept of network slicing so that service guarantees can be provided to certain flows.

At the time of designing the use case, the initial approach was the creation of an IP network infrastructure for interconnecting the TSN domains, being this IP infrastructure being controlled by an SDN Controller. While being realistic, this option represents a challenge from the point of view of the alignment of existing solutions, for both control and data plane, with the support of the network slicing concept being currently defined in IETF. Note that the purpose of the scenario is to validate the effectiveness of the transport slice approach, rather to resolve the integration issues of current state of the art.

Thus, a new strategy is being analzed, as work in progress to complete the scenario. Such strategy is as follows.

Taking advantage of the fact that the TSN switches in the setup support as much as 1 Gbps bit rate interfaces, a way of emulating the behavior of an IP network from the perspective of the data plane is to use instrumentation devices able to introduce impairments on the network links. That is, apply a given behavior in terms of bandwidth, latency, jitter and packet loss to the interfaces of interest. In this manner, different behaviors can be applied to different flows differentiated e.g. by vlans. This is sufficient to emulate the behavior of network slices in the Metro network.

On the other hand, the control elements of TIMING need to interact with the emulated IP network, such interaction following the network slice model as defined in draft-ietf-teas-ietfnetwork-slice-nbi-yang. To that end, the purpose is to develop a control module able to handle the networks slicing NBI in its interaction with the control elements of TIMING, and program the instrumentation equipment for the applications of specific impairmenta to the flows received from the TSN domains.

As said, the interaction with the control elements is expected to follow the IETF NBI YANG model for slicing, so the same YANG model is assumed to be supported. For the interaction with the instrumentation vendor, the idea is to leverage on a programmable API from the instrumentation device.

This is a work in progress, not yet prototyped.

2.8 INDUSTRIAL APPLICATIONS

2.8.1 Description

The architecture of the AGV control software is shown in [Figure 2-13.](#page-83-0)

Figure 2-13: AGV control architecture

Each AGV is equipped with an embedded computer to run all software components based on Linux. This hardware is connected to the AGV by CAN bus. Each AGV sends the information of

its sensors to the coordinator. To do it, inside each AGV there is a component that is connected to the CAN bus and filters the CAN frames with the sensor information. In addition, each AGV is equipped with a LiDAR. The LiDAR driver receives the point cloud information and send it to the coordinator. The coordinator considering the sensor information, generates the control actions needed to govern the movement of the AGVs. In this control problem, the control actions are the longitudinal speed and the steering angle. Finally, there is a PTP inside each AGV to synchronize the internal time by a PTP server.

2.8.1.1 Coordinator

Sensor information Sensor Sensor information information Observer Controller Control actions Control actions conf **PTP Client** PTP protocol

The coordinator is composed by an observer and a controller [\(Figure 2-14\)](#page-84-0).

Figure 2-14: AGV coordinator SW modules

The observer receives the sensor information (speed, and point cloud from the LiDARs) and estimates the current position of each AGV and its speed $(v_1, w_1, x_1, y_1, \theta_1, v_2, w_2, x_2, y_2, \theta_2)$. This information is provided to the controller.

The controller uses this information to generate the speed and steering commands for the AGVs $(v_{r_1}, \gamma_1, v_{r_2}, \gamma_2)$ required to follow the speed profile stablished in the configuration file.

2.8.1.2 AGV software

The AGV is equipped with an external hardware to run the software components. This hardware is connected to the AGV by a CAN bus. The software components in this hardware are shown in [Figure 2-15.](#page-85-0)

Figure 2-15: AGV software modules

A component manages the communication with the LiDAR and acts as a driver sending this information by UDP. Another component listens the CAN frames, filter the frames related with the sensors in the AGV, and send this information by UDP. A component receives the control actions form UDP frames and translate them to CAN frames. In addition, a PTP client manages the time synchronization in the system.

2.8.2 Report on the Current Status

The following modules have been (or are being) developed/implemented:

- **Observer (OBS)**: It will be implemented in python. It requires UDP communication (library socket).
- **Controller (ACON)**: It will be implemented in python. It requires UDP communication (library socket).
- **CAN-UDP**: Implemented in python. It requires UDP communication (library socket). It requires CAN communication (library python-can). Tested without TSN. Used to capture traffic.
- **UDP-CAN**: Implemented in python. It requires UDP communication (library socket). It requires CAN communication (library python-can). Tested without TSN. Testing with TSN pending.
- **LiDAR driver**: An initial version has been implemented in C#. A new version will be implemented in python or C++.It requires TCP and UDP communication (library socket). Tested without TSN. Used to capture traffic. Testing with TSN pending.
- **PTP client**: The software component is ready in [http://git.code.sf.net/p/linuxptp/code.](http://git.code.sf.net/p/linuxptp/code) Testing pending of BeagleBone reception.

In the following tables, the details of the preliminary tests that have been conducted are reported. In particular, for each test, the description, the steps and the preliminary results are detailed.

Test #1

This test validates that the time is correctly synchronized in the AGV and the coordinator, in the event of a change in the system time. All elements are connected in a network with a PTP server.

Preconditions

The AGV is on, and the PTP client in AGVs and coordinator are running.

Execution Steps

1. Manually change the system time to force a difference respect the time in the server

Expected Results

The time must synchronize with the PTP server. Measure the time needed to achieve the synchronization.

Output

ptp4l, phc2sys should report the time offset between PHC and System Clock, which determines if the clocks are synchronized

Example:

Execution

Pending of beagle bone reception

- 3. Set a speed 0 in AGV 2.
- 4. Check the output of OBS
- 5. Set a speed 0.3 in AGV 1.

- 6. Check the output of OBS
- 7. Set a speed 0.3 in AGV 2.
- 8. Check the output of OBS

Expected Results

The observer should correctly estimate the profile speed of the AGVs.

Output

OBS generates a file with the speed estimated from the AGVs.

Execution

CAN-UDP correctly sends the information, processing in observer is pending.

OBS generates a file with the position estimated from the AGVs. This file will show the timestamp and the position.

Execution

Lidar driver has been tested. Testing of observer is pending.

3 OPERATIONAL WORKFLOW FOR E2E SERVICE PROVISIONING

In this section, we present a preliminary version of the operational workflow for the provisioning of end-to-end (E2E) services across the multiple considered technological domains. The workflow will be refined as soon as the integration work advances and changes, if any, will be reported in the next deliverables. Figure 3-1 below depicts a diagram of the sequence of involved operations and actors. Namely, the E2E service provisioning entails the following operations:

- 1. Initially, a service provisioning request arrives to the northbound of the Connectivity Manager (CM). This request may come from a client system or an end-user, and includes the necessary details to set-up and configure the service. In particular, the request contains: the set of endpoints between which the service should be set-up; the service type; a traffic profile that characterizes the communication requirements, such as the required bandwidth; and a collection of Service Level Agreements (SLAs) that state other characteristics of the service, such as the expected degree of reliability and the bounded Key Performance Indicators (KPIs).
- 2. Upon reception, the CM contacts the TSN Controller so as to fetch the most updated data in regards to the infrastructure topology and capabilities. This information is necessary in order to determine which are the viable network paths from a resource availability and connectivity perspective.
- 3. The TSN controller replies with the information.
- 4. Similarly, the CM contacts the Metro controller to fetch information about the underlying Metro network infrastructure.
- 5. The Metro controller replies with the information.
- 6. With the updated information, the CM computes the candidate network paths.
- 7. Once computed, the CM interacts with the Digital Twin (DT) so as to assess the KPIs of the computed path, since they have to be both compliant with what the service request, and the selection of the path should not affect negatively to already established E2E services.
- 8. The DT employs as information to estimate the expected KPIs the requested service information and the currently active flows over the candidate paths. Once determined, the KPIs are returned to the CM.
- 9. With the assessed KPIs, the CM takes a final decision of which is the network path that is going to be configured across the interconnected technological domains.
- 10. The CM contacts with TSN controller with a Connectivity configuration request according to the decided network path, which may include nodes at the multiple technological domains, such as WiFi Access Points (AP) or wired TSN-enabled Ethernet switches.
- 11. The TSN controller sends the specific configuration details to the domain controller or API of the WiFi nodes included in the computed path.
- 12. Similar to the previous step, the TSN controller sends the specific configuration details to the domain controller or API of the Ethernet nodes included in the computed path.
- 13. In the same fashion, the CM contacts the Metro controller with a Connectivity configuration request so as to modify if needed any underlying configuration at the Metro network so as to satisfy the requirements of the computed path.

14. As a last step, the Metro controller applies the received configurations to the physical equipment. At this point, the E2E connectivity is fully provisioned and ready to use by upper layer services.

User	CM	TSN controller	Metro controller	WiFi DT	Eth	Metro
1 E2E service request						
	Path and KPI calculation					
	2 Topology and capabilities deatils request					
	3 Topology and capabilities details reply					
	4 Topology and capabilities deatils request					
	5 Topology and capabilities details reply					
	6 E2E path computation					
	7 KPI assesment request					
	8 Calculated KPIs					
	9 Path decision					
	Data plane configuration 10 Connectivity configuration request					
			11 Network configuration			
			12 Node configuration			
	13 Connectivity configuration request					
				14 Node configuration		
User						
	CM	TSN controller	Metro controller	WiFi DT	Eth	Metro

Figure 3-1: E2E service provisioning workflow.

4 PRELIMINARY INTEGRATIONS

In this section, we report some preliminary integration tests that have been already done and involving some of the TIMING architecture components. In particular, here it is reported the integration works and results of data plane devices interconnection, that is, the TSN Ethernet switch and the Wi-Fi node.

4.1 PRELIMINARY DATA PLANE INTERCONNECTION TESTS

4.1.1 PTP Synchronization tests

Precision Time Protocol (PTP) synchronization plays a pivotal role in Time-Sensitive Networking (TSN) technology as it ensures that devices within a TSN network are tightly synchronized, allowing for accurate and predictable transmission of data. Therefore, the first tests that have been carried out in the data plane interconnection have been in charge of testing the interoperability of the different PTP daemons running in the TSN WiFi Access Point and in the Ethernet TSN switch.

On this purpose, a WiFi TSN AP and an Ethernet TSN switch have been connected together via Ethernet and the synchronization precision that they achieve has been measured for 180 seconds comparing the Pulse-Per-Second (PPS) signal that both devices produce with a Keysight 53230A frequency counter. Each device produces a PPS signal in the zero-crossing of their internal timer-counter. These timer-counters are the ones that get synchronized with PTP. In the tested setup the TSN switch acts as PTP master while the WiFi-TSN AP acts as PTP slave.

Besides, the synchronization information that the PTP daemon offers in the WiFi-TSN AP has also been gathered. The PTP daemon calculates this information based on the different timestamps that it receives both from itself and from the TSN switch.

The picture and block diagram below depict the tested scenario.

Figure 4-1: Block diagram of the PTP synchronization measurement setup

Figure 4-2: PTP synchronization measurement setup

The obtained results from the PPS measurement can be seen in the following figures:

Figure 4-3: PTP synchronization error vs time

Figure 4-4: PTP synchronization error histogram

As can be observed, both devices get synchronized within a +-180 ns jitter that evolves during time. That is the expected result with a PTP synchronization. However, the mean synchronization error is expected to be around 0 while in the performed tests a 2.7 us error can be observed. This is something that needs to be further investigated.

The synchronization log that has been obtained from the PTP daemon in the WiFi-TSN AP is the following:

Figure 4-5: Log information obtained from the PTP daemon in the WIFI-TSN AP

In the "master offset" column the synchronization error (in ns) that the PTP daemon has measured every second can be observed. Similar to the error measured with the PPS signals it is bounded to a +-180 nm margin. In this case, the mean is around 0 so the PTP daemon seems not to be aware of the aforementioned 2.7 us error.

Additionally, we wanted to compare the performance impact by using a Z16 initially, and then using a modified system with enhanced oscillators to reduce jitter: the Z16 Low Jitter (Z16-LJ).

In the histogram of [Figure 4-6,](#page-95-0) we can see that the differences in PPS between TSN WiFi Access Point and Low Jitter are localized in a smaller range and follow a normal distribution, while the differences in PPS between the regular Z16 and TSN WiFi AP exhibit more dispersion and a more erratic distribution.

Figure 4-6: PPS Differences histogram

In [Figure 4-7,](#page-96-0) we compare the calculated values of mean, maximum and minimum differences, range of values, and standard deviation between results using regular Z16 and Z16 Low Jitter. We can observe that the values of PPS differences are higher using Z16 Low Jitter than using regular Z16 (both the mean differences and maximum and minimum values). This aligns with the information obtained from the histogram, where we can see that the values using Z16 Low Jitter are shifted further to the right on the x-axis. However, by examining the range and standard deviation, we can see that the dispersion of differences using Z16 Low Jitter is much lower, just as in the histogram where the data is situated in a smaller range.

Figure 4-7: PPS Differences

The obtained values can be seen in the table below.

We can use the Allan deviation plot in [Figure 4-8](#page-97-0) to further quantify the differences between both devices – the regular Z16 and the enhanced Z16-LJ node – in terms of frequency distribution stability in the short term. Initially, both sets of measurements show increasing deviation with increasing τ, which is typical as longer time intervals can introduce more variation. However, the Z16 LJ graph shows a decrease in deviation at a certain point, indicating a period where the PPS differences are less variable or more stable. This implies that the PPS differences stabilize at a certain point for Z16 LJ, whereas we can see that the same does not happen with Z16.

Figure 4-8: Allan Deviation

The main takeaway from these tests is that we have shown that the PTP stacks available on Safran's TSN switches and the IKERLAN TSN WiFi AP prototype are essentially compatible and capable of propagating a unified time reference between the respective network segments handled by Safran and IKERLAN in the demonstrator scenario of TIMING. These PTP tests show an early, preliminary integration that showcases how this time reference can be distributed between the Safran and IKERLAN network domains with different levels of performance as a function of the available hardware when the default profile of PTP over UDP with unicast connections and end-to-end (E2E) transport is used. Other profiles could potentially work as well and be considered for this integration. Nonetheless, this aspect will be further investigated in future integration tests in the upcoming stages of the project.

5 CONCLUSIONS

This deliverable reports the progress on the implementation of the different software components/modules that compose the TIMING architecture. The validation of the different modules in a stand-alone scenario is the prior step of the proper integration; some preliminary integration tests also reported. The results reported show a significant and solid progress in the implementation work that well position the project towards the final integration and validation.

Moreover, an operational workflow for E2E service provisioning has been also defined and reported. From it, the different steps to materialize the provisioning of a service with specific KPIs requirements over the E2E TIMING infrastructure have been identified as well as the different interactions among the software components of the architecture. This is a significant step achieved towards pushing the overall integration, prior to the architecture validation to be performed in the framework of SP3.

6 REFERENCES

[1] TIMING Project, Deliverable SP1 D1.2 "Year 1 Report on Requirements, Architecture and Performance Evaluation", 2023.

[2] TIMING Project, Deliverable SP1 D1.3 "Preliminary version of software components", 2023

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