



# A Cooperative Switching Algorithm for Multi-Agent Foraging



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## ABSTRACT

The foraging task is one of the canonical testbeds for cooperative robotics, in which a collection of robots have to find and transport one or more objects to specific storagepoints. Efficiency in foraging can be improved with coordinated team of robots. Swarm robotics investigates the bio-inspired behaviors of simple species, that provide complex behaviors at the group level. We present a multiagent foraging algorithm named Cooperative Switching Algorithm for Foraging (C-SAF) inspired from the classical ant system. It provides a quick search, optimal homing paths and quick exploitation of food. A qualitative comparison between some foraging related works and the proposed algorithm is given here, as well as a quantitative comparison which shows that our algorithm outperforms the reference c-marking algorithm across a range of different scenarios.

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## 1. Introduction

Multi-Agent Systems are a suitable approach to develop many multi-robot distributed applications such as: mine detecting (Acar et al., 2003; Gage, 1995), search in damaged buildings (Kantor et al., 2006; Jennings et al., 1997), fire fighting (Marjovi et al., 2009), and exploration of spaces (Landis, 2003; Schilling and Jungius, 1996), where neither a map, nor a Global Positioning System (GPS) are available (Batalin and Sukhatme, 2002). The efficiency of the group of robots in this case can be dramatically improved through coordination. Coordination issues (Yan et al., 2013) of multi-robot systems are considering more robots and more complex tasks, sometimes including robots and humans together. Swarm robotics investigates multi-robot coordination in a distributed way. It is interested in the implementation of systems which are composed of thousands of simple robots rather than one single complex robot (Vaughan, 2008). The challenge is to develop a group of robots with limited perception and computing capabilities to resolve complex tasks in a collective and distributed manner. This would allow for reduction of costs with respect to heavyweight approaches based on powerful and expensive

cognitive robots/agents. In particular, stigmergic-based coordination mechanisms have been used in many robotics problems (aggregation formation and flocking WeiXing et al., 2006, patrolling Pasqualetti et al., 2010, localization and mapping Stipes et al., 2006, exploration and fire searching Marjovi et al., 2009), where agents adopt an indirect communication by depositing pheromone in their environment. This mechanism takes its inspiration from social insects (ants, bees, termites) which provide collectively intelligent systems (Momen, 2013) that, in spite of the simplicity of their individuals, present a highly structured social organization. As a result of this organization, ant colonies can accomplish complex tasks that in some cases far exceed the individual capacities of a single ant (Dorigo et al., 2000).

Foraging is a complex task that involves the coordination of multiple subtasks each constituting a difficult task (searching, harvesting, homing and unloading). It lends itself to multi-robot systems, even if the task can be achieved by one single robot, it is profitable to use multiple robots with careful design of cooperation and coordination strategies (Winfield, 2009). The sophisticated foraging behavior observed in social insects, provides inspirations to produce simple individuals (like ants) that use simple coordination rules and provide more complex (emergent) behaviors as a whole.

We investigate in this paper the Multi-Agent Foraging problem. Many approaches to such problem have been proposed in the literature so far (Momen, 2013; Hoff et al., 2010; Meng et al., 2012; Hoff

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et al., 2013; Lee and Ahn, 2011; Pitonakova et al., 2014). Here, we propose therefore the Cooperative Switching Algorithm for Foraging (C-SAF), an extended version of the Multi-Agent Foraging algorithms in Zedadra et al. (2015a) and Zedadra et al. (2015b). In this paper, we apply it to large-scale foraging systems with hundreds and thousands of agents, where we consider a wider range of system sizes than the ones in literature works (Hoff et al., 2013; Simonin et al., 2014; Shell et al., 2006) and in our previous work (Zedadra et al., 2015a). We analyze and qualitatively compare our proposed approach to relevant related works through an extension of a reference comparison framework (Winfield, 2009). An in-depth performance evaluation based on simulation is given that allows us to quantitatively compare our approach with cooperative and non cooperative foraging approaches. Results show that C-SAF algorithm takes on average less time to finish the foraging, returns larger food amounts and provides optimal food-to-nest paths without the need to revisit a cell several times. Moreover, it shows to be scalable and presents considerable parallelism with growing agent number. Finally, it outperforms the other approaches in terms of efficiency and scalability.

The remainder of the paper is organized as follows: Section 2 presents the background concepts and definitions, and the related work on foraging tasks. Section 3 describes the C-SAF algorithm. A qualitative comparison between C-SAF and the Non-Cooperative Switching Algorithm for Foraging (NC-SAF), c-marking and Non-Cooperative c-marking (NC-c-marking) algorithms, is presented in Section 4. Section 5 details the different scenarios, obtained results and quantitative comparison between the four analyzed algorithms. In Section 6, we provide a discussion on how to go towards a real robotic implementation of our approach. Finally, we draw conclusions in Section 7 and delineate some future works.

## 2. Background and related work

### 2.1. Basic concepts

Here we define and clarify some key terms and concepts that will be used throughout this paper:

- **Agent-Based Modeling (ABM)** is a class of computational models for simulating the actions and interactions of autonomous agents (both individual or collective entities such as organizations or groups) with a view to assessing their effects on the system as a whole and analyzing possible emergent behaviors.
- **Swarm Intelligence** is the study of natural and artificial systems of multiple agents that adopt a distributed autonomous control. Instead, global intelligent behaviors emerge from a cooperating collection of simple individual behaviors (Bonabeau et al., 1999).
- **Stigmergy** is a particular form of indirect communication mediated by modifications of the environment used by social insects to coordinate their actions (Grassé, 1959).
- **Artificial Potential Field (APF)** is a wavefront of integer values written by agents in the environment, to mark the short distance between any cell and the nest (Simonin et al., 2014).
- **Search** is defined as the act of looking into or over carefully and thoroughly in an effort to find or discover something (Méndez and Bartumeus, 2014).
- **Foraging** is the act of searching for and collecting food at one or more storage points. It is a complex task that involves the coordination of multiple other tasks, such as searching, homing, and grabbing (Winfield, 2009).

### 2.2. Multi-Agent Foraging problem definition on 2D grids

In the following, we present a formal definition of the Multi-Agent Foraging problem on two dimensional grids with one nest. The environment  $E$  is represented as 2D grid of cells of size  $N \times N$ .  $E = E_{Free} \cup E_{Obstacles}$ , where  $E_{Free}$  denotes the subset of  $E$  containing all obstacle-free cells and  $E_{Obstacles}$  denotes the subset of  $E$  containing obstacles.  $E_{Free} = E_{Reachable} \cup E_{Unreachable}$ , where  $E_{Reachable}$  denotes the subset of  $E_{Free}$  containing reachable cells (i.e. all cells that are reachable by agents) and  $E_{Unreachable}$  denotes the subset of  $E_{Free}$  containing unreachable cells (i.e. cells enclosed in obstacles).  $E_{Obstacles}$  is a subset of  $E$  defined as a set of obstacle cells where  $Obs_i \subseteq E_{Obstacles}$ . Each cell  $c = (x, y) \in E$  has a maximum of four neighbors  $(x-1, y)$ ,  $(x+1, y)$ ,  $(x, y-1)$ ,  $(x, y+1)$ . Let  $c_0$  be the origin (or nest) and the starting cell for all the agents positioned at the center of the environment (coordinates (0,0)). Let  $E_{Visited}$  be the set of cells already visited where  $E_{Visited} \subseteq E_{Reachable}$  and  $E_{NotVisited}$  the set of cells not yet visited containing at the starting time all the cells of  $E_{Reachable}$  except  $c_0$  where  $E_{NotVisited} \subset E_{Reachable}$  and  $E_{Reachable} = E_{NotVisited} \cup E_{Visited}$ .  $A$  is a set of identical agents  $a_0 \dots a_n$ , where  $n$  is the total number of agents. Each  $a_i$  has initial heading ( $0^\circ = \text{up}$ ,  $90^\circ = \text{right}$ ,  $180^\circ = \text{down}$  or  $270^\circ = \text{left}$ ) and can perceive the four neighbors in up, right, down, and left directions.  $N$  food locations (referred to as food density), each with  $M$  food items (referred to as concentration) are spread on a set of cells included in the list  $C_{goal}$  (list of cells containing a food). Agents do not know the coordinates of the cells in  $C_{goal}$ . The goal is therefore to forage all food in  $C_{goal}$  by minimizing the  $T_{foraging}$  time, i.e. the overall time needed to complete the foraging task. To guarantee completeness of search and reachability of food, we assume that obstacles do not partition  $E$  and do not enclose any agent, nest or food cell. Moreover, the problem definition can be easily generalized from one to many nests.

The components of our Multi-Agent Foraging system (*Environment, Agent, Pheromone*) are modeled as follows:

- **Environment Model:** The environment is organized as a  $N \times N$  grid with several food locations, one or multiple sinks (or nests), obstacles and pheromone markings. The grid is divided into equal squares in a cartesian coordinate system. Grid maps are thought to an efficient metric for navigation in large-scale (Thrun and Bücken, 1996). Obstacles with rectangular or square shapes take place on some fixed cells, the nest is at the center (in a multi-sinks space, sinks take specific positions).  $N$  food with  $M$  items take fixed positions on the grid. In the offline version of searching, an agent has already the map of the world and plans its path before starting the search. Whereas, in the online version the agent plans the path directly when searching. The online version of the grid-based exploration has received considerable attention due to its applicability in practice since grids represent a natural discretization of planar environments (Thrun and Bücken, 1996). It is often used for solving tasks like path planning, localization, search, coverage and surveillance (Balch, 1996; Yean and Chetty, 2012; Lau et al., 2013; Panov and Koceska, 2014; Gabriely and Rimon, 2002; Choi et al., 2009). Apart from 2D-grids, that are the most used environments in the Multi-Agent Foraging problem, other approaches do exist. For instance, topological approaches such as those described in Mataric (1994); Gutjahr (2000) and Dorigo et al. (2006), represent robot environments by graphs. Nodes in such graphs correspond to distinct situations, places, or landmarks. They are connected by arcs if there exists a direct path between them.
- **Agent Model:** Agents are modelled as simple, reactive ant-like agents that can move in the four directions up, down, right and left,

corresponding to  $0^\circ$ ,  $90^\circ$ ,  $180^\circ$  and  $270^\circ$  headings, respectively. Agent can detect pheromone at a specific distance (at one cell step). At each move, agents deposit an amount of pheromone to mark the current cell as already visited.

- **Pheromone Model:** As in traditional ant system, agents in our model release pheromone on the ground on their visited cells or even in their four neighboring cells. The pheromone evaporates according to a specific rate. The total amount of pheromone that evaporates at the position  $i$  in time  $t$  is modeled as in Eq. (1):

$$\Gamma_i(t+1) = \Gamma_i(t) - p * \Gamma_i(t) \quad (1)$$

where  $p$  is a coefficient which represents the evaporation rate of trail between time  $t$  and  $(t+1)$ .

### 2.3. S-MASA algorithm

The S-MASA algorithm (Zedadra et al., 2014), was proposed for multi-target search and foraging tasks. It is designed according to the main ideas of the artificial ant system (Dorigo et al., 2000; Kuyucu et al., 2012; Calvo et al., 2011) and the water vortex dynamics. Basically, the system is a group of homogenous artificial agents (e.g. robots) moving and taking decisions based on the stimuli received from the environment. While agents navigate, they deposit pheromone to mark visited cells and can sense the existence of pheromone in their vicinity. The absence of pheromone in its right cell calls the agent to change its heading and to move to its right cell, else the agent will keep going forward in its current heading. We adopt in this algorithm the opposite logic of ant system theory. So, instead of following the cells that contain pheromone, our agents are guided to not yet visited cells (where no pheromone is present). S-MASA provides: (1) a quick search by the large dispersion of agents when avoiding already visited cells, optimizing by this their chance to find food, (2) optimal paths to return to nest result from the wavefront of pheromone concentration expansion created by agents simultaneously when exploring.

The Detect\_Pheromone\_Adjust\_Heading\_And\_Move() (detailed in Algorithm 1) is used to extract the stimuli and to decide about the next move to take. The agent changes its heading according to rules in Algorithm 1 if there is no pheromone in its right cell (Fig. 1) else it keeps its heading unchanged. The avoid\_obstacle() procedure is used to avoid obstacles. Specifically, agents turn around obstacles in the direction of already visited cells until encountering a not yet visited cell, which means the obstacle is avoided and agents can continue their search.

**Algorithm 1.** Detect\_Pheromone\_Adjust\_Heading\_And\_Move.

```

if ( $\nexists$  pheromone in right cell) then
  if (current_heading = 270) then
    | Set heading to 0;
  else
    | Set heading to current_heading + 90;
  | move forward of 1 cell in the heading direction;
else if ( $\exists$  pheromone in right cell) and (the front cell is not an
  obstacle) then
  | move forward of 1 cell in the heading direction
else if ( $\exists$  pheromone in right cell) and (the front cell is an
  obstacle) then
  | avoid_obstacle();

```

**Statement 2.1.** Considering a set of agents performing Algorithm 1 in a bounded 2D Grid obstacle-free environment (see Section 2.2), each cell will be visited.

**Proof.** According to Algorithm 1, an agent  $a_i$  moves to  $c_x \in E_{NotVisited}$  in its right direction if it exists, else it moves to  $c_x \in E_{Reachable}$  in its

front direction. Each move of  $a_i$  changes the state of  $c_x$  to visited, thus reducing  $E_{NotVisited}$  and extending  $E_{Visited}$ . Since  $E_{NotVisited}$  is limited, Algorithm 1 is complete because the number of moves is bounded by the number of cells in  $E_{NotVisited}$  and, at each step, a cell from  $E_{NotVisited}$  is moved to  $E_{Visited}$ ; at a given time  $t_i$ , all cells are in  $E_{Visited}$  and  $E_{NotVisited}$  is empty.

**Statement 2.2.** Considering a set of agents performing Algorithm 1 in a bounded 2D Grid obstacle environment (see Section 2.2), each reachable cell will be visited.

**Proof.** The strategy used by an agent  $a_i$  to avoid an obstacle is to get around the obstacle in the direction of already visited cells until encountering  $c_x \in E_{NotVisited}$ . According to Algorithm 1,  $a_i$  performs the following steps: (1) looks for  $c_x \in E_{NotVisited}$  in its right heading, if there exists, it moves to it; (2) else, it looks for  $c_x \in E_{Reachable}$  in its front heading, if such a cell exists it moves to it; (3) else, if the cell in front  $\in E_{Obstacles}$  then  $a_i$  avoids the obstacle as follows: (i) keeps moving to  $c_x \in E_{Reachable}$  in its right direction until it locates a  $c_x \in E_{Reachable}$  in front direction; (ii) it keeps then moving to  $c_x \in E_{Reachable}$  in its front direction until it locates a  $c_x \in E_{NotVisited}$  in right, front or left heading respectively. As  $Obs_i$  is finite, in a finite time the  $a_i$  avoids the obstacle and reaches a cell  $c_x \in E_{NotVisited}$ , it moves to such a cell and then restart from step (1). At each step,  $a_i$  changes the state of  $c_x$  to visited, thus reducing  $E_{NotVisited}$  and extending  $E_{Visited}$ . Since  $E_{NotVisited}$  is limited, the number of moves is limited too and because each move ends with the addition of a new cell to  $E_{Visited}$ , at a given time  $t_i$ , all cells are in  $E_{Visited}$  and  $E_{NotVisited}$  is empty, thus the search process ends.

### 2.4. Related work

The problem of coordination of multiple agents is considered complex (Speranzon, 2006; WeiXing et al., 2006). Pheromone-based coordination, inspired by the way social insects coordinate via leaving marks in their environment, provides fundamentals to design alternative strategies that overcome the difficulties (Barnes et al., 2006) imposed by mathematical formulations, agent and environment models (Jiang, 2006; Freeman et al., 2006). Pheromone-based coordination has been widely applied in exploration and surveillance tasks (Calvo et al., 2011; Kuyucu et al., 2012; Marjovi et al., 2009), distributed search and collective cleanup (Liu et al., 2010; Feinerman et al., 2012), multi-target search (Zedadra et al., 2014) and foraging (Fortino et al., 2014).

As we are interested in Multi-Agent Foraging, below we present relevant literature works on Multi-Agent Foraging. In Hoff et al. (2010), authors propose two decentralized foraging algorithms inspired by ants behavior. Robots are simple in sensing and processing power, they communicate with each other in neighborhood using simple InfraRed ring architecture. However, instead of using chemical pheromone, beacon robots are used to store virtual pheromone. The approach is very costly in number of robots (beacons) required to create the paths. The models proposed in Meng et al. (2012) are two ant colony foraging models adapting Panait's and Wilensky's models (Panait and Luke, 2004; Wilensky, 1997), by introducing direct interaction (via direct communication) besides indirect interaction (via pheromone). The proposed models improve the ability of the colony to find the shortest paths, however, agents need memory to store additional information about food location. The foraging problem was investigated in Hoff et al. (2013) through three algorithms (gradient, sweeper and adaptive). In the three algorithms, robots need to communicate directly their current position and food position to their neighbors. In the *gradient* algorithm, agents broadcast information about nest and food gradients, they switch from *beacon* to *walker* robots according to some criteria. While in the *sweeper* algorithm, agents use virtual forces to form a line expanding from nest to food

that sweeps the search world like the hand of clock; when food is found some robots remain as beacons while others act like walkers. The search becomes slower with long robot lines. The third algorithm (*adaptive*) helps the colony to switch between the two aforementioned algorithms: for close food they use gradient algorithm, for far distance food they use sweeper algorithm and when the number of robots is not sufficient to cover the whole world, robots switch to use random walk. Introducing preferences towards some tasks besides stimulus has some impact on the performance of the colony (Momen, 2013). Authors use a group of heterogeneous agents and a set of local rules to communicate with neighboring agents for each group. Agents switch between three tasks: *foraging*, *brood caring* and *resting* based on three thresholds for foraging, brood caring and resting. The model brings sufficient improvements but it is still very parameter dependent and several initial parameters need to be adjusted for best results. Authors in Pitonakova et al. (2014) compare simulated robotics swarm of individualist foragers and a swarm of collective foragers. It was shown that recruitment is useful when resources are hard to find, where agents tend to signal to each other the location of deposits and return home on beacons, they use odometry to recognize their location and the food location. Robots need internal memory to store additional information that helps them to react. In Lee et al. (2013) and Lee and Ahn (2011), a honey-bee inspired model for foraging was proposed to improve the energy efficiency by means of search space and labor division. They employ temporary stores at fixed locations, robots need to communicate directly the food location and use GPS to return home. The approach was tested on small scale systems (25 to 40 robots). Two decentralized foraging algorithms that imitate swarm behaviors using a limited number of robots were proposed in Chatunyakit et al. (2013). In both algorithms, robots are equipped with camera, infrared sensors, GPS and communication device. Each robot has a positioning system installed to determine its current location. Random motion is adopted by the two algorithms at first stages of search and robots transmit at each step their current positions between them. In Simonin et al. (2014), authors propose a Multi-Agent Foraging algorithm. It is a parameter free, distributed and asynchronous version of the wavefront algorithm (Barraquand et al., 1992). Agents while exploring the environment build simultaneously paths between food and the nest which results in computing a wavefront from the agent destination which involves building an ascending Artificial Potential Field (APF) incrementally. Due to the pseudo random walk, paths are not optimal and agents need to visit the same cell several times before the APF reaches its optimal value, which results in large foraging times.

C-SAF presents multiple benefits in comparison to the reported foraging works. The random walk used in most of the works as search strategy provides large foraging time and non-optimal

paths (Simonin et al., 2014). In our work, it is replaced with a new search algorithm called S-MASA. S-MASA allow agents to: (1) search quickly for food by allowing a large dispersion of agents by avoiding already visited cells as much as possible, (2) the vortex-like movements of agents allow building optimal paths by constructing a wavefront expansion of pheromone. Comparing to the related works, C-SAF is scalable, it could be used with large number of agents and large environment sizes. Agents in C-SAF are provided with very limited sensing (perceive the four neighboring cells) and computational capabilities. They communicate indirectly via pheromone and they need very small memory to store the amount of remaining food, and no specific material is needed (compasses for example). Agents return home by descending the negative gradient of pheromone wavefront created when searching (see Section 2.3). After food is located, agents attract each other to the location using pheromone trails.

Table 1 presents the comparison of the surveyed works with our algorithm C-SAF (Algo2) according to an extended taxonomy based on the one in Winfield (2009). Each one of the four major axis (*environment*, *robot*, *strategy* and *performance*) is characterized by a set of properties in the minor axis which can take specific values denoted in *value* axis in Table 1. Features of the *environment* concern: *search space* which can be bounded or unbounded, *source areas* that represent food locations which can be single or multiple, *object types* of limited or unlimited amount of items which take fixed or random positions (*object placement*) and a storage point where food is returned (*sinks*). The *robot* is defined by: a *number* of agents participating at the mission (single or multiple), their *type* either homogenous which means similar in functions and material or heterogeneous, the *object sensing* which is their range of sensors (limited or unlimited), their ability to know the current position (*localization*), their ability at exchanging information with others (*communication*) and their possession of limited or unlimited stores of energy (*power*).

The *strategy* includes the different sub-strategies used in foraging such as: *search*, *grabbing*, *transport*, and *homing* and if there exist mechanisms of *recruitment* or *coordination* to improve the performances. While in *performance*, we added in bold the different metrics used in the works compared in Table 1, in addition to those that exist in the reference taxonomy (Winfield, 2009).

### 3. Cooperative Switching Algorithm for Foraging (C-SAF)

The C-SAF Algorithm is a distributed Multi-Agent Foraging algorithm that creates optimal paths from food to nest by only depositing pheromone while searching. By using this algorithm,

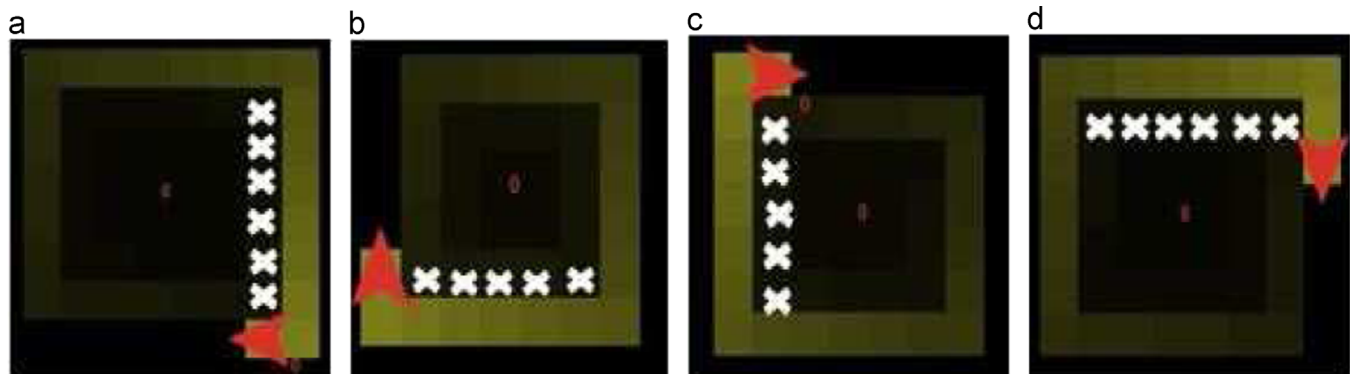


Fig. 1. S-MASA coordination rules which represent the changing of heading from: (a)  $180^\circ$  to  $270^\circ$ , (b)  $270^\circ$  to  $0^\circ$ , (c)  $0^\circ$  to  $90^\circ$ , (d)  $90^\circ$  to  $180^\circ$ , where white crosses represent already visited cells.

agents create simultaneously and synchronously a pheromone wavefront expansion from the nest to the food (see Statement 3.1) and can use the negative gradient to go back to nest (see Statements 3.2 and 3.3). It increases the cooperation between agents to transport food by diffusing pheromone when the food is located. Each agent can switch from a state to another according to surrounding events.

In particular the agent switches between several states: *Look-for-Food*, *Choose-Next-Patch*, *Pick-Food*, *Return-and-Color*, *Return-to-Nest*, *At-Home*, *Climb* and *Remove-Trail*. It starts with looking for food (*Look-for-Food* state). If there is no food, it moves in its environment using *Choose-Next-Patch* rules and deposits pheromones that evaporate with time. It changes automatically to *Look-for-Food* state and, if there is food, picks a quantity of it (*Pick-Food* state), and returns home by following the colored trail, if there exists one (*Return-to-Nest*). If there exists no trail, it diffuses the information to its neighbors by depositing recruitment pheromone and creates a trail by depositing return pheromone (*Return-and-Color*) while returning to home. When the

home is reached, it unloads the food (*At-Home*) and if the food is exhausted it removes the existing trail (*Remove-Trail*), else it climbs the trail to the food location (*Climb*). Fig. 2 shows the finite state machine of our foraging agent whereas the complete algorithm is given in Algorithm 2.

**Statement 3.1.** *In a bounded 2D Grid environment (see Section 2.2), a set of agents performing the C-SAF Algorithm builds food paths when searching.*

**Proof.** Let us define a *food path* as the set of cells free from obstacles relaying nest to food. Let  $W_i$  be the set of cells corresponding to the same cost (pheromone value). Several wavefronts exist  $W_0 \dots W_n$ , where  $W_0$  is the initial wavefront, which corresponds to the starting cell from where all agents start searching. An agent  $a_i$  keeps moving according to the strategy defined in Statement 2.2. At the first step, agents move to reachable non-obstacle cells from  $W_0$  and depose a pheromone creating the wavefront  $W_1$ . At the second step, agents move to reachable non-obstacle cells from  $W_1$  and depose a

**Table 1**

A comparison of Multi-Agent Foraging algorithms based on the enhanced taxonomy. Algo2, Algo3 and Algo5 are respectively C-SAF, NC-SAF, and NC-c-marking (see Sections 3 and 4).

Major axis	Minor axis	Value	Simonin et al. (2014)	Momen (2013)	Meng et al. (2012)	Hoff et al. (2013)	Algo 2	Algo 3	Algo 5
Environment	search space	constrained	X	X	X	X	X	X	X
		single limited		X	X				
	source areas	single unlimited				X			
		multiple limited	X				X	X	X
	sinks	single	X	X	X	X	X	X	X
	object types	multiple static	X	X	X	X	X	X	X
object placement	fixed location			X	X				
	random	X				X	X	X	X
Robot (s)	number	multiple	X	X	X	X	X	X	X
		homogenous	X			X	X	X	X
	type	heterogeneous		X	X				
		limited	X	X	X	X	X	X	X
	object sensing	non	X		X	X	X	X	X
		relative			X				
	communication	direct		X	X	X			
		indirect	X		X		X	X	X
power	unlimited	X	X	X	X	X	X	X	
Strategy	search	random walk	X	X	X	X			X
		trail following	X	X	X		X	X	X
		follow other robots in teams				X			
	grabbing	<b>S-MASA</b>					X	X	
		single	X	X	X	X	X	X	X
	transport	cooperative	X	X	X	X	X		
		single						X	X
	homing	on beacon				X			
		direct trail information			X	X			
	recruitment	follow trail	X	X	X		X	X	X
		direct				X			
	coordination	indirect	X	X	X		X	X	X
direct			X	X	X		X	X	
Performance	time	fixed					X	X	X
		minimum			X				
other metrics		<b>varying agent number</b>	X				X	X	X
		<b>varying environment size</b>	X				X	X	X
		<b>first time to find food</b>					X		
		<b>average hunger level</b>		X					
		<b>food found or not</b>					X		
	<b>rate of returned food</b>					X	X	X	

pheromone to create  $W_2$ . Such process iterates from  $i$  to  $i+1$  until all reachable non-obstacle cells have been reached. At each step agents relay cells from  $W_i$  to cells in  $W_{i+1}$ , thus from the starting cell ( $c_0$ ) there exists a trail that leads to the goal cell ( $c_{goal}$ ). As agents adopt a vortex-like movements, they need to turn around each other in the direction of non-visited cells in obstacle-free environments (or of the visited cells in obstacle environments). Thus the built food paths are non optimal.

**Statement 3.2.** In a bounded 2D Grid environment (see Section 2.2), a set of agents performing the C-SAF Algorithm builds homing paths by following the negative gradient created when searching (see Statement 3.1).

**Proof.** Let us define a *homing path* as the set of cells free from obstacles relaying food to nest. Let  $\mathbb{P}$  be initial pheromone deposited by an agent at time  $t_0$ , and  $\mathbb{P}_j$  its value after evaporation where  $j$  is the number of evaporations that  $\mathbb{P}$  had. At  $t_1$ , agents move to all reachable cells in one step from  $W_0$ , deposit  $\mathbb{P}$ , thus create the wavefront  $W_1$ . At  $t_2$ , agents move to non visited neighbors of  $W_1$ , deposit  $\mathbb{P}$  in current cells, thus create the wavefront  $W_2$  at the same time values of  $W_1$  take the first evaporation  $\mathbb{P}_1$  ( $\mathbb{P}_1 < \mathbb{P}$ ). At  $t_3$ , agents move to non visited neighbors of  $W_3$ , they take the value  $\mathbb{P}$ , while values of  $W_1$  take their second evaporation to get  $\mathbb{P}_2$ , and pheromones of  $W_2$  take their first evaporation to get  $\mathbb{P}_1$  where  $\mathbb{P}_2 < \mathbb{P}_1 < \mathbb{P}$  corresponding to  $W_1$ ,  $W_2$  and  $W_3$  respectively. At each time  $t_i$ , if there exists an obstacle then  $a_i$  keeps avoiding it in the direction of  $c_x \in E_{Reachable}$  until it encounters a  $c_x \in E_{NotVisited}$ . This process iterates from  $i$  to  $i+1$  until all reachable cells have been reached, agents can then follow the negative gradient from the greatest  $\mathbb{P}$  value to the lowest one ( $W_0$ ) to get home. The homing paths are shorter than food paths since agents choose at each step cells with the smallest  $\mathbb{P}$  value in the four neighboring cells.

**Statement 3.3.** The homing paths built according to Statement 3.2 are optimal.

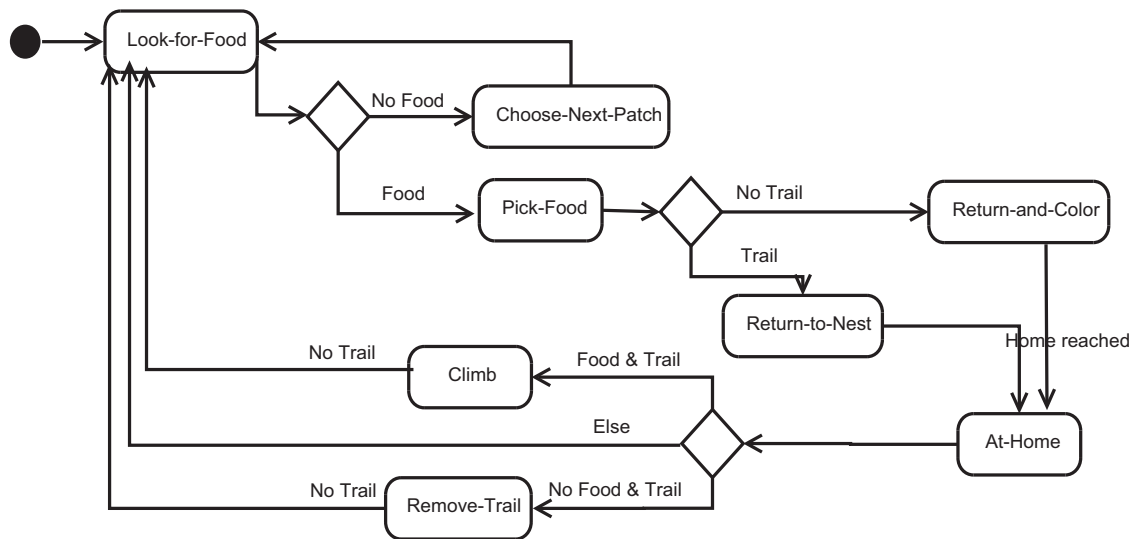
**Proof.** Optimality of the paths is considered with respect to the path length or the distance in number of cells between nest and food. We define an optimal path (shortest path) as the path with the smallest number of cells relaying nest to food. The vortex-like movements allow each agent  $a_i$  at any time  $t_j$  to get four different neighbors:  $c_{right}, c_{back} \in E_{Visited}$  and  $c_{front}, c_{left} \in E_{NotVisited}$  reachable at one step to the right, backward, forward and left directions respectively. Let  $\mathbb{P}_r$  be the pheromone in  $c_{right}$ ,  $\mathbb{P}_b$  the pheromone in  $c_{back}$  and  $\mathbb{P}_x$  the pheromone in the current cell  $c_x$  ( $c_{front}$  and  $c_{left}$  do not contain pheromones since they are not yet visited), since  $c_{back}$  is visited at  $t_{j-1}$  and  $c_{right}$  is visited at  $t_{j-n}$  where  $n$  is the number of steps between  $c_{right}$  and the current cell (following the principle of turns like vortex in movements). At  $t_j$  when  $c_x$  is reached,  $\mathbb{P}_r < \mathbb{P}_b < \mathbb{P}_x$  because  $\mathbb{P}_r$  evaporates  $n \times p$ ,  $\mathbb{P}_b$  evaporates  $1 \times p$  and  $\mathbb{P}_x$  does not evaporate yet (where  $n$  is the number of steps between  $c_{right}$  and  $c_x$  and  $p$  is the evaporation rate). According to this principle, when  $a_i$  returns home, it chooses  $c_x \in E_{Visited}$  with the smallest  $\mathbb{P}$  which is the  $c_{right}$  cell, this process iterates until home is reached, thus from a cell to another the optimal cost is chosen, therefore optimal paths are created.

**Algorithm 2.** C-SAF Algorithm where pheromone is noted  $\mathbb{P}$ .

```

if AT-HOME then
  Unload food;
  if ( $\exists$  trail and food > 0) then goto CLIMB;
  elseif ( $\exists$  trail and food = 0) then goto REMOVE-TRAIL;
  else goto LOOK-FOR-FOOD;
if LOOK-FOR-FOOD then
  if (food > 0) then goto PICK-FOOD;
  else goto CHOOSE-NEXT-PATCH;
if CHOOSE-NEXT-PATCH then
  if (obstacle detected) then Avoid_Obstacle();
  else
    if (brown  $\mathbb{P}$  here) and (brown  $\mathbb{P}$  in right cell) then
      Diffuse( $\mathbb{P}$ ); move to food location using brown cells;
    else
      if (brown  $\mathbb{P}$  here) and (nobrown  $\mathbb{P}$  in right cell) then
        Remove brown trail;
      else
        Lay( $\mathbb{P}$ );
        Detect_And_Adjust_Heading( $\mathbb{P}$ );
        Update( $\mathbb{P}$ );
        Move();
  if PICK-FOOD then
    Pick up a given amount of food;
    Diffuse( $\mathbb{P}$ );
    if ( $\exists$  trail) then goto RETURN-TO-NEST;
    else goto RETURN-AND-COLOR;
if REMOVE-TRAIL then
  while  $\exists$  no trail do
    Move to neighboring colored cell with the greater value of  $\mathbb{P}$ ;
    Update its color to the default one (black);
  goto LOOK-FOR-FOOD;
if CLIMB then
  while  $\exists$  no trail do
    Move to neighboring colored cell with the greater value of  $\mathbb{P}$ ;
  goto LOOK-FOR-FOOD;
if RETURN-TO-NEST then
  while home not reached do
    Move to a colored neighboring cell with the lowest  $\mathbb{P}$  value;
  goto AT-HOME;
if RETURN-AND-COLOR then
  while home not reached then
    Move to a neighboring cell with the lowest  $\mathbb{P}$  value;
    Color that cell to a specific trail color (yellow);
  goto AT-HOME;

```



**Fig. 2.** State diagram showing the possible states of a foraging agent in C-SAF and NC-SAF algorithms. Black circle is the initial state, white diamonds are decision points and rectangles are states.

#### 4. Related algorithms: a comparison

The C-SAF algorithm (Section 3) allow recruiting as much as possible agents in trails to cooperate in exploitation of food. To show that cooperation may improve the performance of the individual agents or the overall behavior of the system, an in-depth quantitative comparison between the C-SAF (Algorithm 2), NC-SAF (Algorithm 3), c-marking (Simonin et al., 2014) (Algorithm 4) and NC-c-marking (Algorithm 5) is given in Section 5.

The NC-SAF algorithm uses the state machine depicted by Fig. 2 and the C-SAF set of behaviors. However, it does not allow cooperation between agents to transport the food when it is found. Instead, each found food is exploited and transported by its finder. The agent finder will not alert the other agents to the food location, hence they proceed to the search process rather than to the exploitation of the food collectively. It uses the same state machine as C-SAF, with changes in the *Choose-Next-Patch* state, where agents in this state do not diffuse brown pheromone to attract other agents. This procedure is used by the C-SAF algorithm to spread information to alert the other agents to the food location in order to exploit it collectively.

**Algorithm 3.** Variants of the NC-SAF Algorithm with respect to the C-SAF Algorithm (see Algorithm 2), where pheromone is denoted by  $\mathbb{P}$ .

```

if CHOOSE-NEXT-PATCH then
  if (obstacle detected) then
    | Avoid_Obstacle();
  else
    | Lay( $\mathbb{P}$ );
    | Detect_And_Adjust_Heading( $\mathbb{P}$ );
    | Update( $\mathbb{P}$ );
    | Move();
  
```

#### **if** PICK-FOOD **then**

```

  | Pick up a given amount of food;
  | if ( $\exists$  trail) then goto RETURN-TO-NEST;
  | else goto RETURN-AND-COLOR;
  
```

```

if (LOOK-FOR-FOOD or REMOVE-TRAIL or CLIMB or RETURN-TO-NEST or RETURN-AND-COLOR or AT-HOME) then
  | do the same as in algorithm 2;
  
```

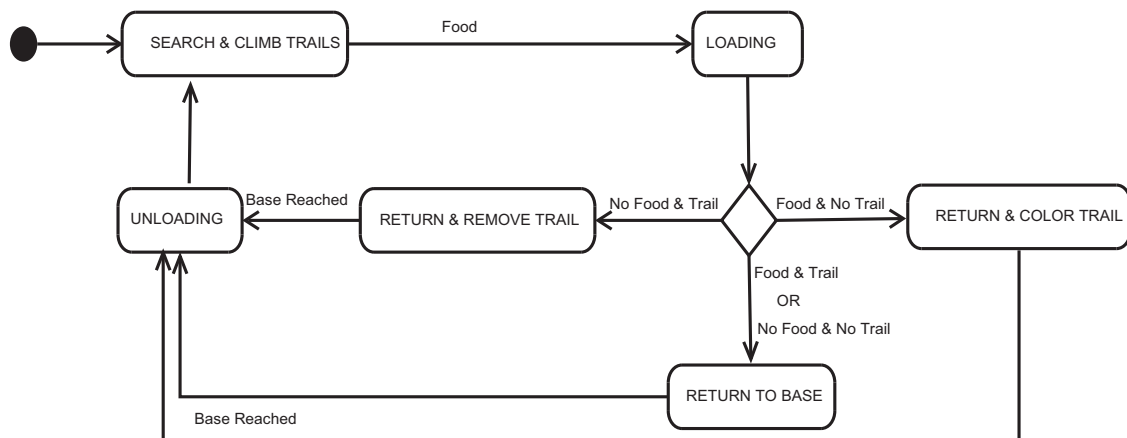
The c-marking algorithm (Algorithm 4, Simonin et al., 2014) is a parameter free foraging algorithm that is a distributed and asynchronous version of the wavefront algorithm (Barraquand et al., 1992). Agents while exploring the environment build simultaneously paths between food and the nest which results in computing a wavefront from the agent destination which involves building an ascending APF incrementally. Agents need to visit the same cell several times before the APF reaches its optimal value. A c-marking agent is always in one of the states depicted by Fig. 3 showing its finite state machine. It always starts from the state *SEARCH & CLIMB TRAIL* which is the default state for all agents. Agents use random walk for search. When food is located, an agent picks a given quantity and returns to nest by using a trail if it exists or creating one if it does not exist, by following the negative gradient. Other agents can be recruited in existing trails. When the food is exhausted the last agent must remove the trail.

NC-c-marking (Algorithm 5) is an alternative of c-marking algorithm that does not allow cooperation to exploit the food, and so there is no recruitment of agents in existing trails. Each agent which found a food can create a trail and after that climb it to return to the food and remove it at the end. The two algorithms use the same state diagram depicted by Fig. 3, where the difference is in the states *SEARCH & CLIMB* and *RETURN AND COLOR TRAIL*.

**Algorithm 4.** c-marking Algorithm (Simonin et al., 2014).

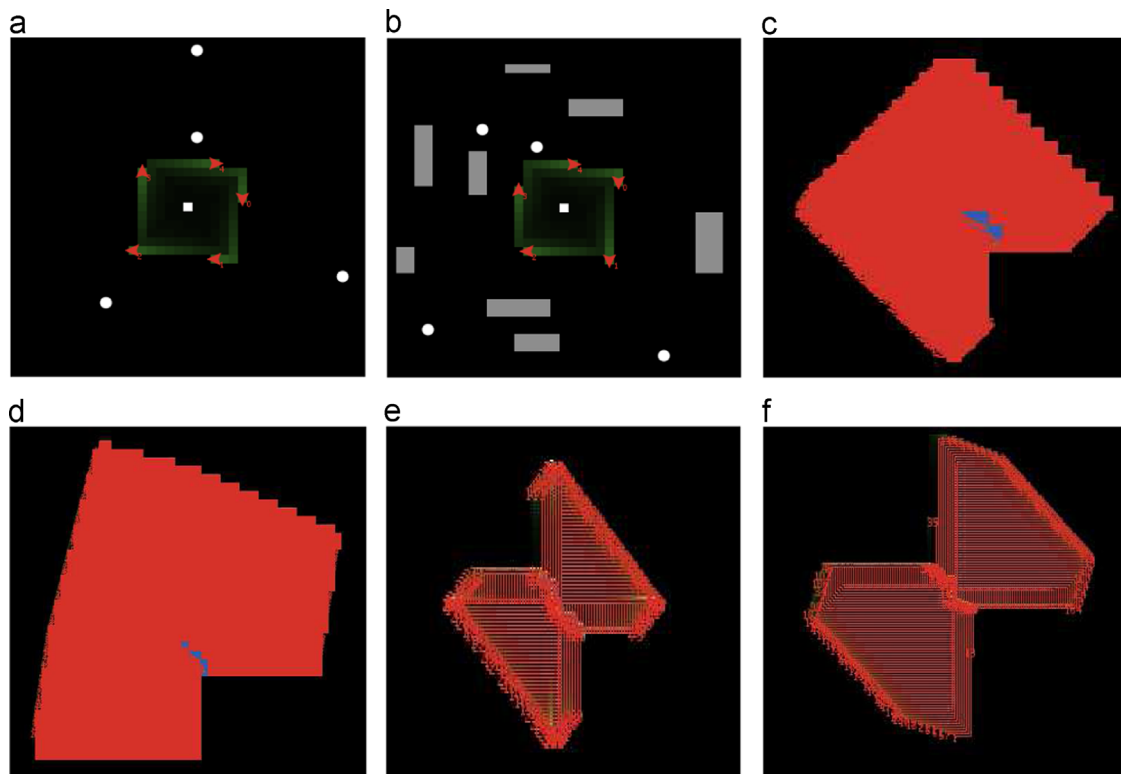
```

if SEARCH & CLIMB (Repeat) then
  if (food exist in neighboring cell) then move into that cell; goto LOADING;
  if (there exists no colored neighboring cell) then
    | goto EXPLORATION & APF CONSTRUCTION;
  else if (previous position was not a trail cell) then
    | Move to the highest valued colored neighboring cell;
  else
    | Move to a new cell with the current trail color;
  goto UPDATE-VALUE;
if EXPLORATION & APF CONSTRUCTION then
  if ( $\exists$  neighboring cells without a value) then move randomly to one of them;
  goto UPDATE-VALUE;
  else move randomly to one of neighbors; goto UPDATE-VALUE;
if UPDATE-VALUE then
  Compute  $val = 1 + \min(4 - neighborvalues)$ ;
  Write  $val$  in the current cell if its different from the starting cell;
if LOADING then
  Pick up a quantity  $Q_{max}$  of food;
  if (food is not exhausted) then if (the cell is colored) then
    | goto RETURN-TO-BASE;
  else
    | goto RETURN-AND-COLOR-TRAIL;
  else goto RETURN-AND-REMOVE-TRAIL;
if RETURN-AND-COLOR-TRAIL then
  Color the cell in a specific color;
  goto UPDATE-VALUE;
  if (base reached) then Unload food;
  goto SEARCH & CLIMB;
  else move to a neighboring cell with the smallest value;
if RETURN-AND-REMOVE-TRAIL then
  if (cell has the trail color) then remove this color and goto UPDATE-VALUE;
  if (base is reached) then Unload food;
  goto SEARCH & CLIMB elseif (there exists a neighboring cell with trail color) then
  move to neighboring trail cell
  else move to the smallest neighboring cell;
  
```



**Fig. 3.** State diagram showing the possible states of a foraging agent in c-marking (Simonin et al., 2014) and in NC-c-marking algorithms. Black circle is the entry point, white circle is a decision point and rectangles are states.





**Fig. 4.** World setups used in simulations: (a) Obstacle-free environment; (b) Obstacle environment, where arrows are agents, the white cell in the center is the nest, the white circles are food and the gray clusters are obstacles; (c), (d), (e) and (f) present C-SAF execution samples at different stages where Laden agents are in the middle line, and Walker agents spread themselves diagonally on sides, with respectively 1 and 2 food locations.

**Algorithm 5.** Variants of the NC-c-marking Algorithm with respect to the c-marking Algorithm (see Algorithm 4).

```

if SEARCH & CLIMB (Repeat) then
  if (food exist in neighboring cell) then
    | move into that cell; update marker to true;
    | goto LOADING
  if marker  $\bar{true}$  then
    | if there exists a colored cell with great value then
    | | move to it
    | else
    | | goto EXPLORATION & APF CONSTRUCTION
    | else goto EXPLORATION & APF CONSTRUCTION;
  if RETURN-AND-COLOR-TRAIL then
    | color the cell in a specific color; goto UPDATE-VALUE;
    | memorize the last step before the base cell;
    | if (base reached) then Unload food;
    | goto SEARCH & CLIMB;
    | else move to a neighboring cell with the smallest value;
  if (EXPLORATION & APF CONSTRUCTION or UPDATE-VALUE or
    LOADING or
    RETURN-AND-REMOVE-TRAIL) then
    | do the same as in algorithm 4;

```

## 5. Performance evaluation

In this section, we present a description of the proposed foraging Framework in Section 5.1 and evaluate the performance of the four algorithms (C-SAF, NC-SAF, c-marking and NC-c-marking)

in obstacle-free and obstacle environments. We present, in Section 5.2, the simulation parameters and the performance indices used to evaluate the algorithms. In Section 5.3, we describe four fundamental scenarios that are used for simulations, and report the importance and the benefits of such scenarios in understanding and comparing the algorithms. In the final Section 5.4, we present, discuss and compare the obtained results by the different algorithms in the four scenarios.

### 5.1. A simulation framework for Multi-Agent Foraging

The algorithms C-SAF, NC-SAF, c-marking and NC-c-marking have been implemented through our framework for Multi-Agent Foraging. The framework is composed of several components: the *search world* which is a set of *cells* modeled as stationary agents on which the actors of simulation are deployed, the *food* modeled as stationary agent with a fixed position and limited amount which can be exploited by foragers, *obstacles* and *nest* which take place at fixed positions, *foragers* are modeled as mobile agents that search for food using a searching algorithm (e.g. the S-MASA algorithm Zedadra et al., 2014, see Section 2.3) and transport it, when found, to home while depositing pheromones and *behavior* which includes a specific foraging algorithm (e.g. Algorithms 2–5). The framework is designed to be extensible as the definition of each object is independent from the others, and it can be reused in other foraging algorithms. It is also flexible as new algorithms or even behaviors can be defined and seamlessly used; hence, developers of other foraging algorithms can add new classes, remove classes which are not important in their systems, or modify existing methods (the *move* method of an agent for example), or even add other states (such as *recharging* to take into account energy-aware behaviors), that can be directly defined in the behavior class.

The framework is currently implemented in Netlogo (Wilensky, 1999), which is a multi-agent programmable modeling environment that allows us to rapidly prototype systems of situated agents evolving in a two dimensions world. It provides a sequential synchronization of execution where each agent must finish all the commands in the block before the next agent starts and the order in which agents are chosen to run the commands is random. Simulations are performed within a grid environment where cells can be either empty or occupied by food, agent, obstacle or the nest. The position of obstacles is fixed for all simulations in order to exclude its impact on the multi-agent system performance. The *Multi-Agent Foraging framework* which is notably platform-independent, is translated into a Netlogo platform-specific model and implementation. In Netlogo, each *Forager* is modeled as a *Turtle* having the ability to update information of neighboring or current cells when needed (such as depositing pheromones). Foragers communicate indirectly through their world. Communication is modeled directly in the behavior procedure. Agent behavior is modeled as procedures which can be called from any other procedure. *World* is modeled as a set of *Patches* (immobile agents) disposed on grid. Foragers, food, obstacles and nest are placed on the *world patches*. *Food* is modeled as a set of *Patches* which can take place at some positions on *world patches*.

Even though, in this paper, we implemented the Multi-Agent Foraging simulation framework in Netlogo, the framework can be easily ported onto other agent-based simulator (or general-purpose simulation environment) such as Repast (Collier, 2003), and ELDA Sim (Fortino and Russo, 2012).

### 5.2. Simulation parameters and performance indices

The two world setups (obstacle-free and obstacle environments), which are used for simulations, including positions of nest, food, obstacles and agents, are reported in Fig. 4. There are several simulation parameters that must be properly set: (i) the *World size* is the dimension of the search space, it is a grid of  $N \times N$  cells; (ii) the *Food density* is the number of food locations (sites), each site contains a limited amount of food. These locations take fixed positions in the environment; (iii) the *Food concentration* indicates the amount of food that each site contains (we refer to it as unit in the paper); (iv) the *Agent's capacity* is the amount of food (units) that an agent can transport at each time; (v) the *Agent's number* is the number of agents that participate in each simulation.

To evaluate the performance of the C-SAF algorithm and the other algorithms, three performance indices have been defined:

- *Average Foraging Time* – The amount of time in seconds needed to finish the foraging mission. It is when all the food sites are discovered and exhausted.
- *Total Food Returned* – The total amount of food returned to the nest by all the agents until a given foraging time.
- *Average Path Length* – Represents the path length between food and the nest, measured as number of cells that form the path.

### 5.3. Simulation scenarios

Simulations are based on the Netlogo-based Multi-Agent Foraging simulation framework (see Section 5.1). Each simulation is repeated 50 times in the four scenarios. The average value is then calculated from the 50 trials for each scenario. Four fundamental scenarios are used to test the performance of the algorithms (see Table 2):

- *Scenario 1* is set up to test which of the parameters presented in Section 5.2 can affect the average foraging time. It entails four

subscenarios. Subscenario 1 is designed to analyze the impact of the *Agent number* on foraging time. Achieving the foraging task can be very time consuming. Using ticks does not reflect the real time needed to accomplish the task, thus it will be difficult to show the gain of the proposed algorithms in according to computation time. To show the main differences when using each of the two units (virtual ticks and real seconds), we analyzed the average foraging time in subscenario 1 through both ticks and seconds. Subscenario 2 is to test how *World size* can affect the average foraging time and if there is any minimal size under which a number of agents can carry efficiently the foraging task. In order to identify whether or not distributing on single or multiple food locations is beneficial, we used subscenario 3 where we varied the *Food density* from 1 to 10. To test the effect of increasing the *capacity of agents*, we defined subscenario 4 in which the capacity is varying from 1 to 100 units.

- *Scenario 2* is defined to show the evolution pattern of returned food over time for each of the four algorithms in obstacle-free and obstacle environments.
- *Scenario 3* is designed to study the effect of the agent number on the *Path length*, whether or not using more agents will reduce the path length.

**Table 2**  
Parameters of Scenarios 1, 2, 3, and 4.

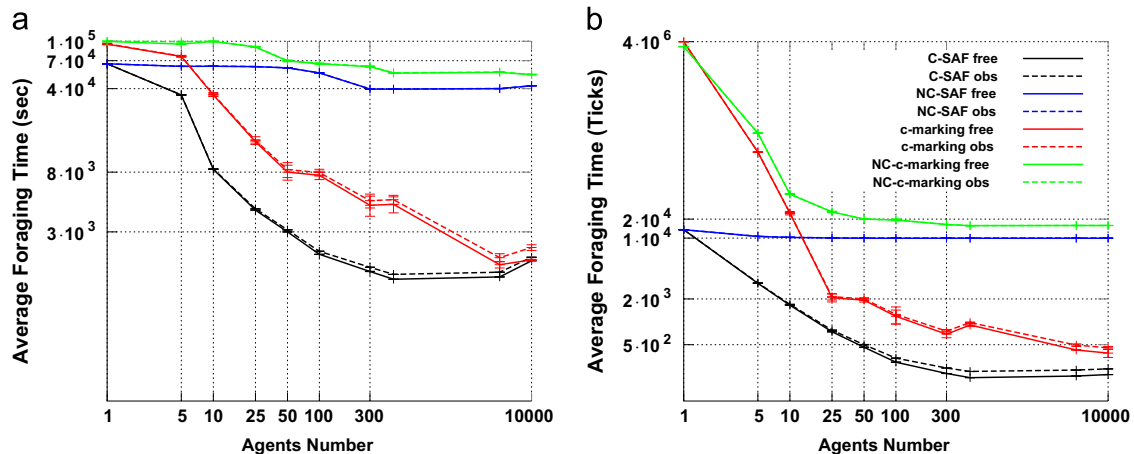
Parameter	Value
<b>Scenario 1</b>	<i>Average time analysis</i>
World size	100 × 100 – 1200 × 1200 cells
Number of agents	1–10000
Food density	1–10 sites
Food concentration	500–1500 units
Agent's capacity	1–100 units
<b>Scenario 2</b>	<i>Returned food analysis</i>
World size	1000 × 1000 cells
Number of agents	800
Food density	1 site
Food concentration	500 units
Agent's capacity	1 unit
<b>Scenario 3</b>	<i>Path length analysis</i>
World size	800 × 800 cells
Number of agents	1–1000
Food density	1 site
Food concentration	500 units
Agent's capacity	1 unit
<b>Scenario 4</b>	<i>Scalability analysis</i>
World size	100 × 100 – 1200 × 1200 cells
Number of agents	50–10000
Food density	1 site
Food concentration	500–1500 units
Agent's capacity	1–10 units

**Table 3**  
Definition of the terms used in the results tables and figures.

Term	Definition
Obstacle-free world (free)	Results in obstacle-free environment for the four algorithms
Obstacle world (obs)	Results in obstacle environment for the four algorithms
C-SAF	Results for Cooperative Switching Algorithm for Foraging
NC-SAF	Results for non-cooperative version of C-SAF
c-marking	Results for c-marking algorithm
NC-c-marking	Results for the non-cooperative version of c-marking algorithm
STD-DEV	Standard deviation of obtained results

**Table 4**  
Average foraging time in ticks and seconds when varying agent number.

Agent number	1	10	100	300	500	5000	10000
<b>Obstacle-free world- Time in ticks</b>							
C-SAF	16029	1713	316	226	200	210	219
STD-DEV	0,1	5,2	3	0,4	0,4	0,9	0,4
NC-SAF	16029	12771	12451	12451	12451	12451	12451
STD-DEV	0,1	0,2	3	0	0	0	0
c-marking	4221763	26205	1232	723	953	454	414
STD-DEV	1020	1320	250	71	45	0,8	50
NC-c-marking	3657570	46122	21428	18676	18076	18144	18157
STD-DEV	1520	456	455	81	31	49	44
<b>Obstacle-free world- Time in seconds</b>							
C-SAF	65210	8569	1677	1206	1038	1086	1482
STD-DEV	5,1	1,2	1,5	5	2	3	5
NC-SAF	65210	62224	54735	40215	40019	40451	42719
STD-DEV	2,3	1,2	1,6	2	2	4	2
c-marking	95320	35799	7672	4309	4385	1373	1509
STD-DEV	1120	920	620	820	670	85	16
NC-c-marking	100345	100251	65300	61764	54517	55322	52906
STD-DEV	620	850	540	702	450	42	23
<b>Obstacle world- Time in ticks</b>							
C-SAF	16069	1753	356	266	240	250	259
STD-DEV	0,1	1,3	3	0,47	0,47	0,9	0,47
NC-SAF	16069	12811	12491	12491	12491	12491	12491
STD-DEV	0,1	3,2	5	0	0	0	0
c-marking	4221833	26275	1302	793	1023	524	484
STD-DEV	1310	520	321	41	15	20	20
NC-c-marking	3657615	46167	21473	18721	18121	18189	18202
STD-DEV	1120	890	530	61	11	29	24
<b>Obstacle world- Time in seconds</b>							
C-SAF	65312	8671	1779	1308	1140	1188	1584
STD-DEV	2,2	1,2	1,3	5	2	4	4
NC-SAF	65312	62327	54837	40317	40121	40553	42821
STD-DEV	3,2	1,3	1,2	4	4	8	3
c-marking	95721	36201	8073	4710	4786	1774	1910
STD-DEV	1359	1320	520	620	420	120	105
NC-c-marking	100540	100652	65701	62165	54918	55724	53307
STD-DEV	1020	985	545	423	145	12	16



**Fig. 5.** Average foraging time when varying the agent number: (a) Average foraging time measured in seconds, (b) Average foraging time measured in ticks, where the legend in (b) is the same as in (a).

- Scenario 4 is proposed to test the scalability of the algorithms and to analyze how much and in which way agent density variation can affect performances.

5.4. Performance and comparison of cooperative and non-cooperative algorithms

Here we show the results of C-SAF, operating in obstacle-free and obstacle environments, and compare it with NC-SAF, c-marking and

NC-c-marking algorithms. We fixed the position of food and obstacles in all simulations to exclude their impact on performances. Table 3 defines terms used in tables and figures about performance evaluation results.

5.4.1. Average time analysis (Scenario 1 results)

We first present in this section the performance of the four algorithms as the number of agents grows where the average foraging time is measured in seconds and in ticks respectively (see

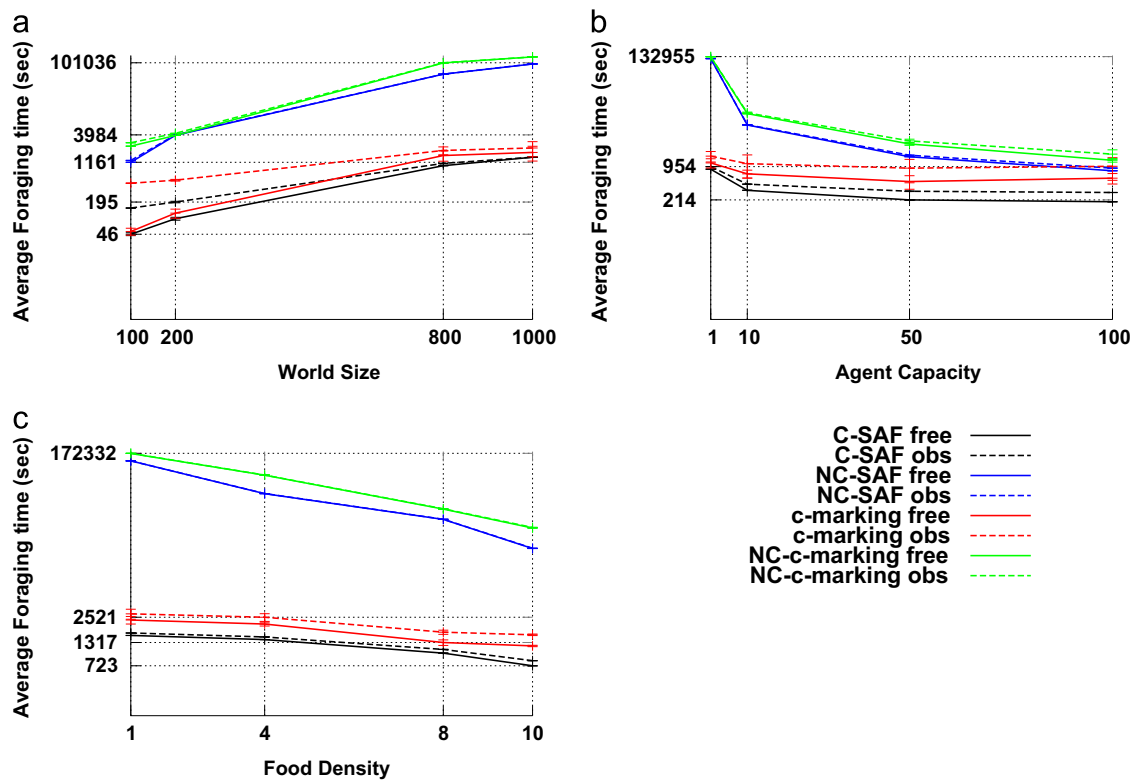


Fig. 6. Results of scenario 1 when varying: (a) world size, (b) agent capacity and (c) food density.

Table 5  
Average foraging time (sec) when varying world size.

World size	100 × 100	200 × 200	800 × 800	1000 × 1000
<b>Obstacle-free world</b>				
C-SAF	46	92	999	1355
STD-DEV	3	6	8	5
NC-SAF	1161	3881	60620	95975
STD-DEV	4	5	5	9
c-marking	52	118	1569	1803
STD-DEV	8	24	349	578
NC-c-marking	2389	3860	100634	131428
STD-DEV	2	5	234	845
<b>Obstacle world</b>				
C-SAF	149	195	1101	1458
STD-DEV	2	7	9	7
NC-SAF	1264	3984	60722	96077
STD-DEV	8	6	9	9
c-marking	454	519	1970	2204
STD-DEV	5	19	344	720
NC-c-marking	2790	4261	101036	131830
STD-DEV	2	9	533	734

Table 6  
Average foraging time (sec) when varying agent capacity.

Agent capacity	1	10	50	100
<b>Obstacle-free world</b>				
C-SAF	852	330	214	197
STD-DEV	4	1	4	2
NC-SAF	121941	6165	1490	792
STD-DEV	5	4	3	1
c-marking	1130	692	490	569
STD-DEV	230	123	144	132
NC-c-marking	132554	10313	2639	1272
STD-DEV	430	223	149	102
<b>Obstacle world</b>				
C-SAF	954	433	316	299
STD-DEV	4	3	1	1
NC-SAF	122043	6267	1592	894
STD-DEV	8	6	9	9
c-marking	1532	1093	892	970
STD-DEV	355	520	431	455
NC-c-marking	132955	10714	3040	1673
STD-DEV	355	320	231	355

Table 4 and Fig. 6, results are obtained by fixing the other simulation parameters as follows: world size=1200 × 1200 cells, food density=1, food concentration=1500 units, agent capacity=10 units, nest number=1). The algorithm was first evaluated with small number of agents (1, 5, 10, 25 and 50) and then this number was progressively increased to reach 10000. While the foraging time increases faster as the number of agents increases from 1 to 500. It is higher with 1 agent (65210 s), halved with 5 agents (35796 s) and even reduced more with 10 and 25 agents (8569 s and 3916 s respectively). It is halved from 50 to 300 and reduced considerably from 300 to 500. Above 500 agents, the foraging time decreases slower progressively until it reaches 1482 s, in obstacle-free environment, and 1584 s, in obstacle environment, with 10000 agents. Between 5000 and 10000 agents, in C-SAF there is a

slight increase of the foraging time; this happens only when using one food location (as in this case) as it is related to the length of search path that agents need to traverse to get to the food location (Fig. 4(c) and (d) shows an example of execution of such a case). Conversely, when we used multiple food locations (see Fig. 4 (e) and (f)), the foraging time keeps decreasing as also subscenario 4 (see below) will confirm. In c-marking algorithms, foraging time becomes faster as the number of agents increases. However, if the food position is out of the wavefront expansion the pseudo random walk can slow down the search task, thus the foraging time. Foraging time in non-cooperative protocols is very high compared to the cooperative ones, however, NC-SAF give better results than the c-marking one when the number of agents is low (1 and 5 agents) because in c-marking algorithm the wavefront created

with a low number of agents is small and the agents take much more time in search. The C-SAF algorithm gives better results than the other three cooperative and the non-cooperative protocols; it is much faster (up to 4 times more) than the c-marking algorithm; in fact, while C-SAF obtains its best performance with 500 agents (1140 s, in obstacle environment), c-marking obtains its best performance with 5000 agents (1744 s). This implies that C-SAF not only outperforms c-marking in terms of foraging time but also it would save a lot of resources so allowing for practical feasibility of the approach. The standard deviation is smaller in C-SAF algorithms and higher in c-marking algorithms where it decreases when increasing the number of agents. Results in obstacle environments are higher than in obstacle-free ones in the four protocols. Curves obtained in the four algorithms show similar trends when using seconds (Fig. 5(a)) or when considering ticks (Fig. 5 (b)). However, to reconcile the results to those that can be obtained in the real world, we use only seconds as measurement unit for the average foraging time in the other scenarios.

In Table 5 and Fig. 6(a), we can observe that foraging time is almost linear with respect to the growth of the world size in the four protocols. Results are obtained by fixing the other simulation parameters as follows: agent number=800, food density=1, food concentration=500 units, agent capacity=1 unit, nest number=1.

**Table 7**  
Average foraging time (sec) when varying food density.

Food density	1	4	8	10
<b>Obstacle-free world</b>				
C-SAF	1580		1002	723
STD-DEV	2	1	2	1
NC-SAF	141931	60965	31452	14842
STD-DEV	8	3	3	2
c-marking	2350	2120	1317	1205
STD-DEV	230	103	94	21
NC-c-marking	171931	98210	40895	25122
STD-DEV	630	223	143	102
<b>Obstacle world</b>				
C-SAF	1682	1522	1104	825
STD-DEV	2	2	3	1
NC-SAF	142033	61068	31554	14944
STD-DEV	10	2	3	3
c-marking	2751	2521	1718	1606
STD-DEV	355	250,1	91,8	30
NC-c-marking	172332	98612	41297	25523
STD-DEV	730	323	103	95

**Table 8**  
Returned food over time until the foraging completion for all the four algorithms.

Time	500	1355	1450	1803	2035	95,975	96,345	131,428	152,412
<b>Obstacle-free world</b>									
C-SAF	41	500	500	500	500	500	500	500	500
STD-DEV	1	1	0	0	0	0	0	0	0
NC-SAF	2	5	6	8	9	500	500	500	500
STD-DEV	2,5	1,5	0	0	0	0	0	0	0
c-marking	0	194	251	500	500	500	500	500	500
STD-DEV	0	2	2	35	0	0	0	0	0
NC-c-marking	0	2	2	3	4	335	336	500	500
STD-DEV	0	0,7	0,7	1	2	3	7	1	0
<b>Obstacle world</b>									
C-SAF	35	367	500	500	500	500	500	500	500
STD-DEV	1	2	1	0	0	0	0	0	0
NC-SAF	0	0	1	2	4	466	500	500	500
STD-DEV	0	1	1	2	1	2	1	0	0
c-marking	0	180	241	489	500	500	500	500	500
STD-DEV	0	3	3	23	13	0	0	0	0
NC-c-marking	0	0	0	1	2	331	332	493	500
STD-DEV	0	0	0	0,7	0,7	10	4	16	1

World size can slow down the foraging time in C-SAF algorithm if it cannot contain a large number of agents, so reducing agent cooperation in food transportation. In the c-marking algorithm, when increasing the world size and the number of agents is not sufficient to sustain the wavefront expansion, the paths are not optimal and the foraging time increases. However, in this scenario the number of agents is sufficient to create the wavefront to food so helping at reducing foraging time. The non-cooperative protocols provide very high foraging time in comparison to the cooperative ones and C-SAF provides the smallest foraging time in all world configurations. Standard deviation is smaller in C-SAF protocols and increases in c-marking protocols from smaller to higher values because of the random walk. Results in obstacle environments are higher than in obstacle-free environments.

Increasing the agent capacity helps at reducing the average foraging time in the four algorithms (see Table 6 and Fig. 6(b)), results are obtained by fixing the other simulation parameters as follows: world size=800 × 800 cells, agent number=500, food density=1, food concentration=500 units, nest number=1). However, C-SAF algorithm takes much less average time to finish the foraging regarding the other three algorithms. The average time for c-marking agents is better than the one obtainable with the non-cooperative algorithms, NC-SAF and NC-c-marking, but it is much worse compared to C-SAF. The time decreases due to the reduction in the number of turns needed to transport all the food (e.g. agents need 10 turns to transport 10 units when capacity is 1, but they need 1 turn when capacity is 10 units).

Finally, in subscenario 4 (Table 7 and Fig. 6(c)), results are obtained by fixing the other simulation parameters as follows: world size=1000 × 1000 cells, agent number=500, food concentration=500 units, agent capacity=1 unit, nest number=1), we varied food density from 1 to 10, which means that we got multiple food locations and the total amount of food is distributed among them. The average foraging time decreases as the food density increases until it is halved in C-SAF and c-marking. However, the average time increases dramatically when the food locations are far away from the nest, especially in c-marking algorithms where the paths are too long and non optimal. Similar performance results are obtained in obstacle environment, with additional time needed to avoid obstacles.

5.4.2. Returned food analysis (Scenario 2 results)

C-SAF algorithm has a very fast time to locate and exhaust all the food (1355 s in obstacle-free environment), which is not the

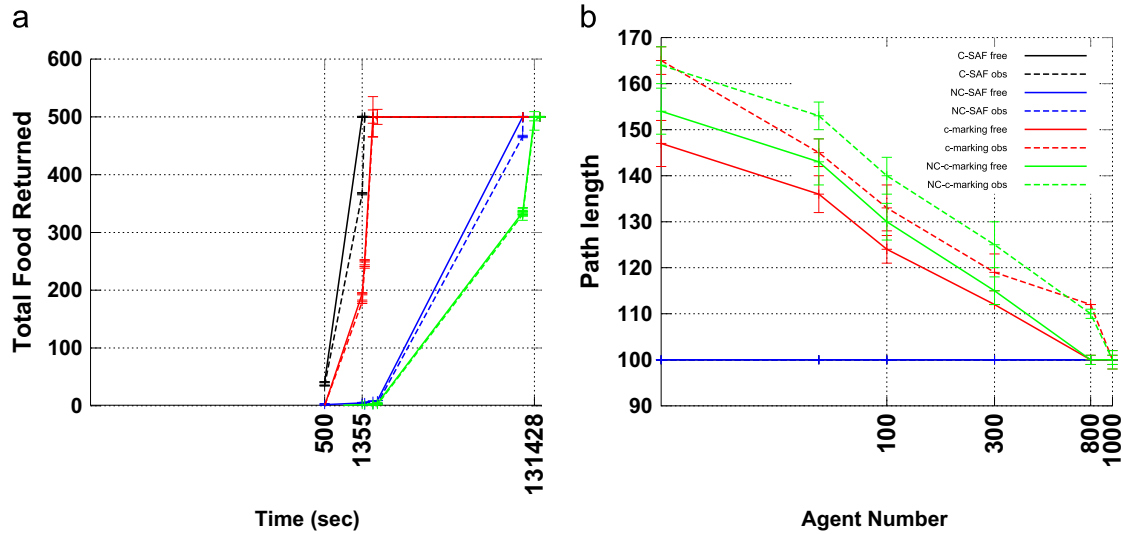


Fig. 7. (a) Returned food analysis; (b) Path length analysis. The legend in (a) is the same as in (b).

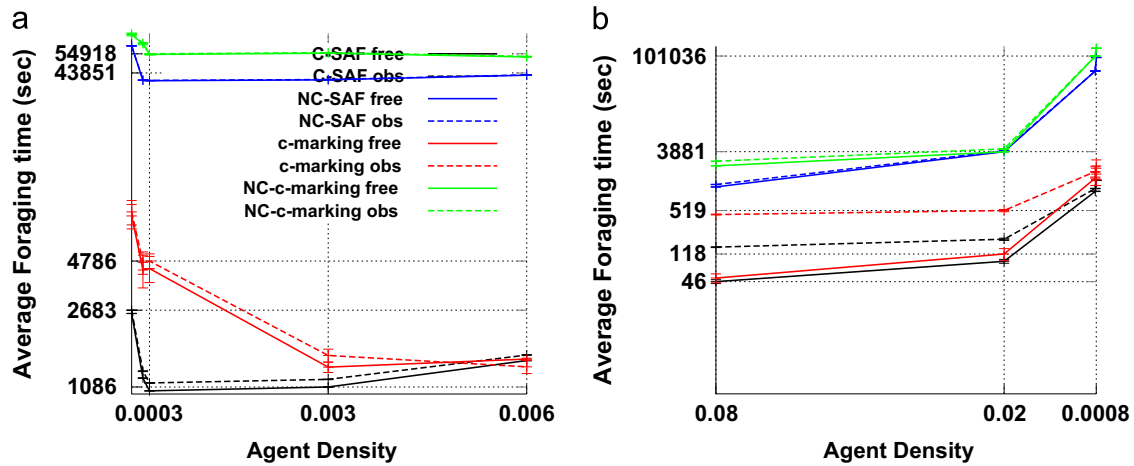


Fig. 8. Scalability analysis: (a) increasing agent density by fixing world size and varying agent number; (b) decreasing agent density by fixing agent number and varying world size. The legend in (b) is the same as in (a).

Table 9  
Average path length when varying agent number.

Agent number	10	50	100	300	800	1000
<b>Obstacle-free world</b>						
C-SAF	100	100	100	100	100	100
STD-DEV	0	0	0	0	0	0
NC-SAF	100	100	100	100	100	100
STD-DEV	0	0	0	0	0	0
c-marking	147	136	124	112	100	100
STD-DEV	5	4	3	0,7	1	1
NC-c-marking	154	143	130	115	100	
STD-DEV	5	5	4	3	1	2
<b>Obstacle world</b>						
C-SAF	100	100	100	100	100	100
STD-DEV	0	0	0	0	0	0
NC-SAF	100	100	100	100	100	100
STD-DEV	0	0	0	0	0	0
c-marking	165	145	133	119	112	100
STD-DEV	3	3	5	4	0,7	2
NC-c-marking	164	153	140	125	110	100
STD-DEV	4	3	4	5	1	1

case for the non-cooperative protocols that take about 95975 and 131428 s (NC-SAF and NC-c-marking in obstacle-free environment, respectively). C-SAF retains higher performance than c-marking:

about 500 s better. In this scenario, we reported a sample of the total amount of food returned overtime without the analysis of the parameters that can affect its evolution. However, this scenario is very much related to Scenario 1 and, for all the test configurations of Scenario 1, we had the same performance profile for the four protocols. Table 8 and Fig. 7(a), show the results obtained by the four algorithms in obstacle-free and obstacle environments (results are obtained by fixing the other simulation parameters as follows: world size=1000 × 1000 cells, agent number=800, agent capacity=1 unit, food density=1, food concentration=500 units, nest number=1).

#### 5.4.3. Path length analysis (Scenario 3 results)

Regardless of the number of agents, C-SAF and NC-SAF protocols produce optimal paths (with 100 cells length) simultaneously when exploring their environment. In c-marking algorithms, the path length is non optimal and decreases when increasing the number of agents until it reaches the optimal value with 800 agents in obstacle-free environment (1000 agents in obstacle environment) (see Table 9 and Fig. 7(b), results are obtained by fixing the other simulation parameters as follows: world size=800 × 800 cells, agent capacity=1 unit, food density=1, food concentration=500 units, nest number=1). It is worth noting that agents in the c-marking algorithm need to visit the same

**Table 10**  
Scalability analysis: increasing agent density by fixing world size and varying agent number.

Agent density	0.00003	0.0002	0.0003	0.003	0.006
Agent number	50	300	500	5000	10,000
World size	1200 × 1200	1200 × 1200	1200 × 1200	1200 × 1200	1200 × 1200
<b>Obstacle-free world</b>					
C-SAF	2580	1206	1038	1086	1482
STD-DEV	4	5	2	3	5
NC-SAF	60302	40215	40019	40451	42719
STD-DEV	4	2	2	4	2
c-marking	8142	4309	4385	1373	1509
STD-DEV	1160	820	670	85	16
NC-c-marking	68900	61764	54517	55322	52906
STD-DEV	1050	702	450	42	23
<b>Obstacle world</b>					
C-SAF	2683	1308	1140	1188	1584
STD-DEV	5	5	2	4	4
NC-SAF	60404	40317	40121	40553	42821
STD-DEV	4	4	4	8	3
c-marking	8544	4710	4786	1774	1910
STD-DEV	1220	620	420	120	105
NC-c-marking	69302	62165	54918	55724	53307
STD-DEV	789	423	145	12	16

**Table 11**  
Scalability analysis: decreasing agent density by fixing agent number and varying world size.

Agent density	0.08	0.02	0.001	0.0008
World size	100 × 100	200 × 200	800 × 800	1000 × 1000
Agent number	800	800	800	800
<b>Obstacle-free world</b>				
C-SAF	46	92	999	1355
STD-DEV	3	6	8	5
NC-SAF	1161	3881	60620	95975
STD-DEV	4	5	5	9
c-marking	52	118	1569	1803
STD-DEV	8	24	349	578
NC-c-marking	2389	3860	100634	131428
STD-DEV	2	5	234	845
<b>Obstacle world</b>				
C-SAF	149	195	1101	1458
STD-DEV	2	7	9	7
NC-SAF	1264	3984	60722	96077
STD-DEV	8	6	9	9
c-marking	454	519	1970	2204
STD-DEV	5	19	344	720
NC-c-marking	2790	4261	101036	131830
STD-DEV	2	9	533	734

cell several times to get to its optimal APF value. The position of the food can affect the path length: it will be shorter if food is close to nest and longer otherwise. However, in c-marking algorithms the number of agents, world size, food position and pseudo random walk can affect the path length. Results obtained with C-SAF and NC-SAF in obstacle-free and obstacle environment are very close and their curves overlap in Fig. 7(b).

#### 5.4.4. Scalability analysis (Scenario 4 results)

In this section, we analyze the scalability of the proposed algorithm through two subscenarios on the basis of the *Agent Density*, defined as *Agent Number* divided by *World Size*: (1) increasing the agent density from 0.00003 to 0.006, by varying the agent number and fixing the world size; (2) decreasing the agent density from 0.08 to 0.0008, by fixing the agent number and varying the world size. With reference to subscenario 1, results are shown in Table 10 and Fig. 9(a). In the C-SAF case, the foraging time quickly decreases with the increase of agent density until 0.0003 value; from such value, the foraging time starts increasing

**Table 12**  
Comparison highlights of the most important results between C-SAF and the other algorithms in obstacle environment.

Algorithms	Performance indices			
	Foraging time	Returned food	Path length	Scalability
C-SAF	1140 s with 500 agents	500 units in 1450 s	100 cells with 10 agents	2683 s with low density=3E-6; 149 s with high density=8E-2;
c-marking	1774 s with 5000 agents, equals to C-SAF with 200	500 units in 2035 s, 241 units in 1450 s	100 cells with 1000 agents; 165 cells with 10 agents	8544 s with low density=3E-6; 454 s with high density=8E-2;
NC-SAF	40,121 s with 500 agents	500 units in 96,345 s, 1 unit in 1450 s	100 cells with 10 agents	60,404 s with low density=3E-6; 1264 s with high density=8E-2;
NC-c-marking	53,307 s with 10,000 agents	500 units in 152,412 s, 0 unit in 1450 s	100 cells 1000 agents; 164 cells with 10 agents	69,302 s with low density=3E-6; 2790 s with high density=8E-2;

slightly. A similar trend happens for the c-marking algorithm even though with worse performances. The main reason for this is the long path that agents need to traverse to reach the food location.

In subscenario 4 of Scenario 1 (see Section 5.4.1), we showed that using multiple food locations helps at halving the average foraging time, thus avoiding the limitation in scalability when increasing too much the agent density. Regarding subscenario 2, the obtained results are shown in Table 11 and Fig. 8(b). Here we can see that time increases with the decrease of the agent density: it is proportional to the growth of the world size. Obviously, increasing the world size requires more time for exploration and transportation of food. In such scenario, the other simulation parameters are set as follows: agent capacity=10 units, food density=1, food concentration=500 units, nest number=1.

#### 5.4.5. Summary of results

In the four different analyzed scenarios (average time, returned food, path length, and scalability analysis), we have shown that C-SAF algorithm outperforms the considered cooperative (c-

marking) and non cooperative (NC-SAF and NC-c-marking) algorithms. In Table 12, we report a summary of the comparison between C-SAF and the other algorithms, highlighting the most important performance differences in the four analyzed scenarios. In particular, C-SAF provides on average less time and resources (in terms of number of agents) to finish the foraging task (see first column of Table 12), larger amount of returned food at a given time (see second column of Table 12), and optimal paths (see third column of Table 12). The differences between C-SAF and the other algorithms mainly stems from the more rapid exploration of the environment and the more rapid transportation of food to the nest. Concerning space exploration, our agents use the S-MASA algorithm (see Section 2.3) that, based on a guided search to non-visited cells in vortex-like movements, provides both a large dispersion, therefore a quicker search, and optimal paths to return home. Regarding food transportation, C-SAF agents deposit pheromone to attract other agents in vicinity when a food is found and, due to the vortex-like movements, a large number of agents can be recruited and cooperate to transport the food to the nest, thus accelerating the exploitation task. Moreover, the C-SAF algorithm has also shown a good scalability (see forth column of Table 12).

## 6. Toward a robotic implementation

Stigmergic communications via pheromone trails have shown to efficiently coordinate a team of robots and to allow them to quickly explore a given terrain (Beckers et al., 1994; Svennebring and Koenig, 2004). Most of the works on stigmergic coordination are based on unrealistic assumptions (Kuyucu et al., 2012). This kind of assumptions can help in analyzing and improving pheromone-based algorithms in simulations before their real robotic implementation. One of the key difficulties in the real development of such mechanisms is the implementation of the pheromone itself and how it can interact with agents. To date, several techniques for virtual marking have been proposed.

- Robots can mark physically their trail via physical marks using: *alcohol* (Sharpe and Webb, 1998), *virtual marks* (Svennebring and Koenig, 2004) or *RFID tags* (Mamei and Zambonelli, 2005).
- Robots communicate a pheromone model using wireless network (Vaughan et al., 2000).
- Robots transmit virtual pheromone using infrared communication (Payton et al., 2001).
- Robots can switch to beacons or even use immobile beacons to communicate pheromone like information (Barth, 2003).

Despite the proposed approaches, the implementation of pheromone is still in its early stages and most of the works are available in research laboratories. One of the works that we believe it constitutes an important way to the real implementation of our approach is the one in Ranjbar-Sahraei et al. (2012), where pheromone is defined as an electrical marker that can be placed at given positions in the environment, it can fully evaporate after time. From a theoretical point of view, two changes can take place. First, the robots should take specific initial locations on the four headings (0°, 90°, 180°, 270°) in order to avoid the bottleneck situation around the nest. Second, additional approaches to avoid collisions between robots when using the same homing trail are needed (e.g. leaving the trail for Laden robots). A case study that we plan to consider in the future, is deploying a group of mobile robots, based on a pheromone mechanism similar to the one proposed in Ranjbar-Sahraei et al. (2012), which can perform tasks such as searching for objects. We intend to use the agent programming methodologies and tools from Fortino et al. (2015) and Fortino and Russo (2012), to perform a rapid prototyping of our

approach in the PROFETA language (Fortino et al., 2013) in order to investigate which design and implementation choices are required before the real implementation.

## 7. Conclusion

In this paper, we have presented a distributed foraging algorithm called C-SAF. C-SAF agents are simple, reactive, with small memory and limited perception capabilities. They communicate by depositing pheromone, i.e. stigmergic communications. We have also presented a flexible simulation framework for Multi-Agent Foraging that can be reused and/or extended to address the needs of other foraging algorithms.

Four fundamental scenarios have been modeled and simulated to test the effect of the basic parameters (agent number, world size, food density, food concentration, agent capacity) on the performances of the four cooperative and non-cooperative algorithms (C-SAF, NC-SAF, c-marking and NC-c-marking) in terms of the defined performance indices, i.e. average foraging time, amount of returned food and average path length. In the first scenario, we have varied the basic parameters to test their effect on the average foraging time measured in real seconds; in the second scenario, we have shown the evolution pattern of returned food over time; in the third scenario, we varied the number of agents to observe its impact on the length of created paths; finally in the fourth scenario, we analyzed the scalability of the algorithm when varying agent density. In all the four scenarios, the C-SAF provided the best results in terms of average foraging time, total amount of returned food and average path length, both in obstacle-free and obstacle environments. The average foraging time was up to 4 times less than the one obtainable by the c-marking algorithm which is a main reference in the literature. Moreover simulation results demonstrate the benefit of cooperation in Multi-Agent Foraging when the team of agents is efficiently coordinated.

The C-SAF outperforms the c-marking one in three fundamental points: (1) the quick search provided by the S-MASA algorithm, (2) the optimal homing paths created simultaneously and synchronously while exploring and, (3) the high level of cooperation to exploit and transport food by diffusing pheromone to neighboring cells. But of course the C-SAF algorithm has some degradation in performances in some cases: (1) the presence of obstacles with complex shape, (2) when food is close to boundaries a small number of agents will contribute to transport because most of them already reach the boundaries and (3) for large-scale worlds, visiting all the environment to reach the food that is very far from the nest, could duplicate the area to search and foraging time increases.

In the near future, we intend to test our approach using real robots, after its implementation and validation supported by the ELDAMeth methodology and the PROFETA language. Moreover, we will provide the implemented Multi-Agent Foraging framework in open-source to support researchers interested in modeling and understanding particular Multi-Agent Foraging systems. In fact, by providing an open-source common framework, we may also facilitate comparison among many different foraging systems.

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