

# Goldbach's Conjecture — A Route to the Inconsistency of Arithmetic

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**Abstract.** This paper proves an inconsistency in Peano arithmetic (PA). We express two propositions — a strengthened form of the strong Goldbach conjecture and its negation — using a specific set that varies depending on these propositions. On the other hand, we show that this set remains unchanged under the assumptions of the two statements. This causes a contradiction.

**Notations.** Let  $\mathbb{N}$  denote the natural numbers starting from 1, let  $\mathbb{N}_n$  denote the natural numbers starting from  $n > 1$  and let  $\mathbb{P}_3$  denote the prime numbers starting from 3.

Strengthened strong Goldbach conjecture (SSGB): *Every even integer greater than 6 can be expressed as the sum of two different primes.*

**Theorem.** *PA is contradictory, i.e. the statement FALSE can be derived.*

*Proof.* We define the set  $S_g := \{ (pk, mk, qk) \mid k, m \in \mathbb{N}; p, q \in \mathbb{P}_3, p < q; m = (p + q) / 2 \}$ .

SSGB is equivalent to saying that every integer  $n \geq 4$  is the arithmetic mean of two different odd primes and so it is equivalent to saying that all integers  $n \geq 4$  appear as  $m$  in a middle component  $mk$  of  $S_g$ . So, by the definition of  $S_g$  we have

$$\text{SSGB} \Leftrightarrow \forall n \in \mathbb{N}_4 \quad \exists (pk, mk, qk) \in S_g \quad n = m$$

$$\neg \text{SSGB} \Leftrightarrow \exists n \in \mathbb{N}_4 \quad \forall (pk, mk, qk) \in S_g \quad n \neq m.$$

The set  $S_g$  has the following two properties.

First, the whole range of  $\mathbb{N}_3$  can be expressed by the triple components of  $S_g$  ("covering"), because every integer  $x \geq 3$  can be written as some  $pk$  with  $k = 1$  when  $x$  is prime, as some  $pk$  with  $k \neq 1$  when  $x$  is composite and not a power of 2, or as  $(3 + 5)k / 2$  when  $x$  is a power of 2;  $p \in \mathbb{P}_3, k \in \mathbb{N}$ . So we have

$$(C) \quad \forall x \in \mathbb{N}_3 \quad \exists (pk, mk, qk) \in S_g \quad x = pk \quad \vee \quad x = mk.$$

A few examples of the covering:

$x = 19$ : (**19·1**, 21·1, 23·1), (**19·1**, 60·1, 101·1)

$x = 27$ : (**3·9**, 7·9, 11·9)

$x = 38$ : (**19·2**, 21·2, 23·2)

$x = 42$ : (**3·14**, 5·14, 7·14), (**7·6**, 9·6, 11·6)

$x = 4096$ : (3·1024, **4·1024**, 5·1024)

$x = 10000$ : (**5·2000**, 6·2000, 7·2000).

Second, according to the statement SSGB, all pairs  $(p, q)$  of distinct odd primes are used in the definition of the set  $S_g$  (“*maximality*”). So we have

**(M)**  $\forall p, q \in \mathbb{P}_3, p < q \quad \forall k \in \mathbb{N} \quad (pk, mk, qk) \in S_g$ , where  $m = (p + q) / 2$ .

The property (C) excludes the possibility that there is an  $n \geq 4$  different from all  $m$ , i.e.  $\neg$ SSGB, for the reason that  $n$  is different from all  $S_g$  triple components  $pk$  and  $mk$ .

The property (M) excludes the possibility that there is an  $n \geq 4$  different from all  $m$ , i.e.  $\neg$ SSGB, for the reason that  $n$  is the arithmetic mean of a pair of primes not used in  $S_g$ . Thus (M) excludes the possibility that  $\neg$ SSGB applies due to a missing prime number pair. The proof would no longer be possible if we left out any prime number pair in the formulation of SSGB and  $S_g$ .

Therefore, in the case of  $\neg$ SSGB, neither  $\neg$ (C) nor  $\neg$ (M) applies.

The basic idea is now the following.

*There are two possibilities for  $S_g$ , exactly one of which must occur: Either there is an  $n \in \mathbb{N}_4$  in addition to all the numbers  $m$  defined in  $S_g$  or there is not. The latter is equivalent to SSGB and the former is equivalent to  $\neg$ SSGB.*

*Since, due to (M), an  $n \geq 4$  different from all  $m$  cannot be the arithmetic mean of a pair of primes not used in  $S_g$  and since, due to (C), this  $n$  equals a component of some  $S_g$  triple that exists by definition, the covering of  $\mathbb{N}_3$  by the  $S_g$  triples in the case  $n$  exists ( $\neg$ SSGB) is equal to that in the case  $n$  does not exist (SSGB). This causes a contradiction because in the case SSGB the numbers  $m$  defined in  $S_g$  take all integer values  $x \geq 4$  whereas in the case  $\neg$ SSGB they don't.*

The following steps are independent of the choice of  $n$  if, in the case of  $\neg\text{SSGB}$ , there is more than one that is different from all  $m$ . For example, the minimal such  $n$  works.

We split  $S_g$  into two complementary subsets in the following way. For any  $y \in \mathbb{N}_3$ , we write

$S_g = S_{g+}(y) \cup S_{g-}(y)$ , with

$$S_{g+}(y) := \{ (pk, mk, qk) \in S_g \mid \exists k' \in \mathbb{N} \quad pk = yk' \vee mk = yk' \vee qk = yk' \}$$

$$S_{g-}(y) := \{ (pk, mk, qk) \in S_g \mid \forall k' \in \mathbb{N} \quad pk \neq yk' \wedge mk \neq yk' \wedge qk \neq yk' \}.$$

We define

$$S_1 := \{ (pk, mk, qk) \in S_g \mid \text{SSGB} \}$$

$$S_2 := \{ (pk, mk, qk) \in S_g \mid \neg\text{SSGB} \wedge ((C) \wedge (M)) \}.$$

Under the assumption  $\text{SSGB}$  there is no  $n \in \mathbb{N}_4$  in addition to all  $m$  defined in  $S_g$ . Then, under this assumption, we can choose any  $y \in \mathbb{N}_3$  such that  $S_g = S_{g+}(y) \cup S_{g-}(y)$ . So,

$$(1.1) \quad \forall y \in \mathbb{N}_3 \quad \text{SSGB} \Rightarrow S_1 = S_{g+}(y) \cup S_{g-}(y).$$

Under the assumption  $\neg\text{SSGB}$  there is an  $n \in \mathbb{N}_4$  as described above. Then, since the possibilities  $\neg(C)$  and  $\neg(M)$  for the existence of  $n$  are ruled out,

$$(1.2) \quad \neg\text{SSGB} \Rightarrow S_2 = S_{g+}(n) \cup S_{g-}(n).$$

$S_{g+}(n) \cup S_{g-}(n)$  is independent of  $n$ . Hence, from (1.2) we get

$$(1.2') \quad \forall y \in \mathbb{N}_3 \quad \neg\text{SSGB} \Rightarrow S_2 = S_{g+}(y) \cup S_{g-}(y).$$

Now, we will make use of the following principle.

If two sets of (possibly infinitely many)  $x$ -tuples are equal, then the sets of their corresponding  $i$ -th components are equal;  $1 \leq i \leq x$ .

To this end, for each  $k \in \mathbb{N}$  we define

$$M_1(k) := \{ mk \mid (pk, mk, qk) \in S_1 \}$$

$$M_2(k) := \{ mk \mid (pk, mk, qk) \in S_2 \}.$$

Then, applying the principle above to the middle component of the triples  $(pk, mk, qk)$ ,  
 $((1.1) \wedge (1.2'))$  implies

$$(2.1) \quad \forall k \in \mathbb{N} \quad \forall y \in \mathbb{N}_3 \quad \text{SSGB} \Rightarrow M_1(k) = \{ mk \mid (pk, mk, qk) \in S_{g^+}(y) \cup S_{g^-}(y) \}$$

$\wedge$

$$(2.2) \quad \forall k \in \mathbb{N} \quad \forall y \in \mathbb{N}_3 \quad \neg \text{SSGB} \Rightarrow M_2(k) = \{ mk \mid (pk, mk, qk) \in S_{g^+}(y) \cup S_{g^-}(y) \}.$$

Setting  $M_1 := M_1(1)$  and  $M_2 := M_2(1)$ , we get

$$(2.1') \quad \forall y \in \mathbb{N}_3 \quad \text{SSGB} \Rightarrow M_1 = \{ m \mid (p, m, q) \in S_{g^+}(y) \cup S_{g^-}(y) \}$$

$\wedge$

$$(2.2') \quad \forall y \in \mathbb{N}_3 \quad \neg \text{SSGB} \Rightarrow M_2 = \{ m \mid (p, m, q) \in S_{g^+}(y) \cup S_{g^-}(y) \}.$$

Since for every  $y \in \mathbb{N}_3$   $S_{g^+}(y) \cup S_{g^-}(y)$  equals  $S_g$  by definition, for every  $y \in \mathbb{N}_3$   
 $\{ m \mid (p, m, q) \in S_{g^+}(y) \cup S_{g^-}(y) \}$  equals the set  $X := \{ m \mid (p, m, q) \in S_g \}$ . So, from  
 $((2.1') \wedge (2.2'))$  we obtain

$$(3) (SSGB \Rightarrow M_1 = X) \wedge (\neg SSGB \Rightarrow M_2 = X).$$

The set  $X$  is a free variable in (3) that is either equal to  $\mathbb{N}_4$  or to some non-empty proper subset  $Y$  of  $\mathbb{N}_4$ .

Now, we make use of the following rule.

Let  $P = P(A)$  be a proposition that depends on a set  $A$ . Then, for any set  $B$ ,

$$(\text{we have a proof of } P(A) \wedge \text{we have a proof of } A = B) \Rightarrow \text{we have a proof of } P(B).$$

In the special case that  $A$  is a free variable that is replaced by the value  $B$ , the above conjunct (we have a proof of  $A = B$ ) is trivially true.

Since the set  $X$  is a free variable in (3) and since we have a proof of (3), we can apply the above rule with  $P = (3)$ . If  $X = \mathbb{N}_4$  we use the rule with  $A = X$  and  $B = \mathbb{N}_4$ , and if  $X = Y$  we use it with  $A = X$  and  $B = Y$ . Then, since either  $X = \mathbb{N}_4$  or  $X = Y$ , from (3) we obtain

$$(3.1) \text{ we have a proof of } (SSGB \Rightarrow M_1 = \mathbb{N}_4 \wedge \neg SSGB \Rightarrow M_2 = \mathbb{N}_4)$$

✓

$$(3.2) \text{ we have a proof of } (SSGB \Rightarrow M_1 = Y \neq \mathbb{N}_4 \wedge \neg SSGB \Rightarrow M_2 = Y \neq \mathbb{N}_4).$$

This implies

$$(3.1') \begin{array}{l} (\text{we have a proof of } (SSGB \Rightarrow M_1 = \mathbb{N}_4) \\ \wedge \\ \text{we have a proof of } (\neg SSGB \Rightarrow M_2 = \mathbb{N}_4)) \end{array}$$

✓

$$(3.2') \begin{array}{l} (\text{we have a proof of } (SSGB \Rightarrow M_1 = Y \neq \mathbb{N}_4) \\ \wedge \\ \text{we have a proof of } (\neg SSGB \Rightarrow M_2 = Y \neq \mathbb{N}_4)) \end{array}.$$

Now, we will establish a contradiction to  $((3.1') \vee (3.2'))$ .

Under the assumption SSGB the set  $X = \{ m \mid (p, m, q) \in S_g \}$  is equal to  $\mathbb{N}_4$  and under  $\neg$ SSGB it is equal to  $Y \neq \mathbb{N}_4$ . Therefore,

**(4.1)** we have a proof of  $(\text{SSGB} \Rightarrow M_1 = \mathbb{N}_4)$

$\wedge$

**(4.2)** we have a proof of  $(\neg\text{SSGB} \Rightarrow M_2 = Y \neq \mathbb{N}_4)$ .

Then,  $((3.1') \vee (3.2'))$  together with  $((4.1) \wedge (4.2))$  implies

**(5.1)** we have a proof of  $(\neg\text{SSGB} \Rightarrow M_2 = \mathbb{N}_4)$

$\vee$

**(5.2)** we have a proof of  $(\text{SSGB} \Rightarrow M_1 = Y \neq \mathbb{N}_4)$ .

Because of  $((4.1) \wedge (4.2))$  and because

$\text{SSGB} \Rightarrow M_2 = \{ \} \neq \mathbb{N}_4$

and

$\neg\text{SSGB} \Rightarrow M_1 = \{ \} \neq Y$ ,

we have a proof that  $(M_2 = \mathbb{N}_4)$  is false and we have a proof that  $(M_1 = Y \neq \mathbb{N}_4)$  is false. So,  $((5.1) \vee (5.2))$  yields

**(6.1)** we have a proof of  $\text{SSGB}$

$\vee$

**(6.2)** we have a proof of  $\neg\text{SSGB}$ .

Since we have neither a proof of SSGB nor of  $\neg$ SSGB, both (6.1) and (6.2) are false.

Therefore, we obtain  $(\text{FALSE} \vee \text{FALSE})$  and thus FALSE.

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