

UNIVERSITÀ DEGLI STUDI DELL'AQUILA

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DIPARTIMENTO DI INGEGNERIA E SCIENZE  
DELL'INFORMAZIONE E MATEMATICA



## “Documentation of the DALI Interpreter”

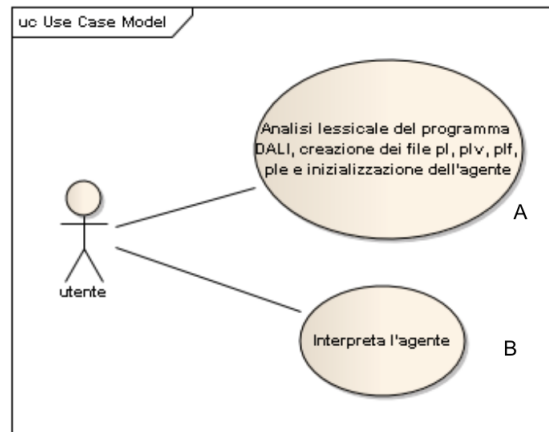
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# Detailed Documentation of the DALI Interpreter

The first step that was done, when the work of this thesis began, was to revise the whole implementation, generating at the same time a quantity of documentation such as to adequately support future interpretations and extensions of the interpreter. It has been produced a detailed description of the UML Diagram realized in 2010 by a student in his thesis (Lanti, D.: Reingegnerizzazione dell'interprete DALI ed integrazione con il Web Semantico,(Università degli Studi dell'Aquila,2010)), then the following images are extracted by it. The diagrams produced are UML Diagrams, in particular Use Case Diagram and Activity Diagram. These diagrams allow to describe a system by a behavioral point of view. The Activity Diagrams have been realized at different levels of detail: the main diagram shows all the macroscopic activity is actually an activity composed of many other activities(atomic or compound) and actions, and which are illustrated by other Activity Diagrams. In the picture 1 it is shown a tree of DALI Interpreter in which is represented the principal predicate that are called: the nodes represented with a circle described the predicates that are called other predicates, the nodes represented with a rectangle described the predicates that aren't called other predicates.



## 1. Use Case Diagram: Functionality of the DALI Interpreter



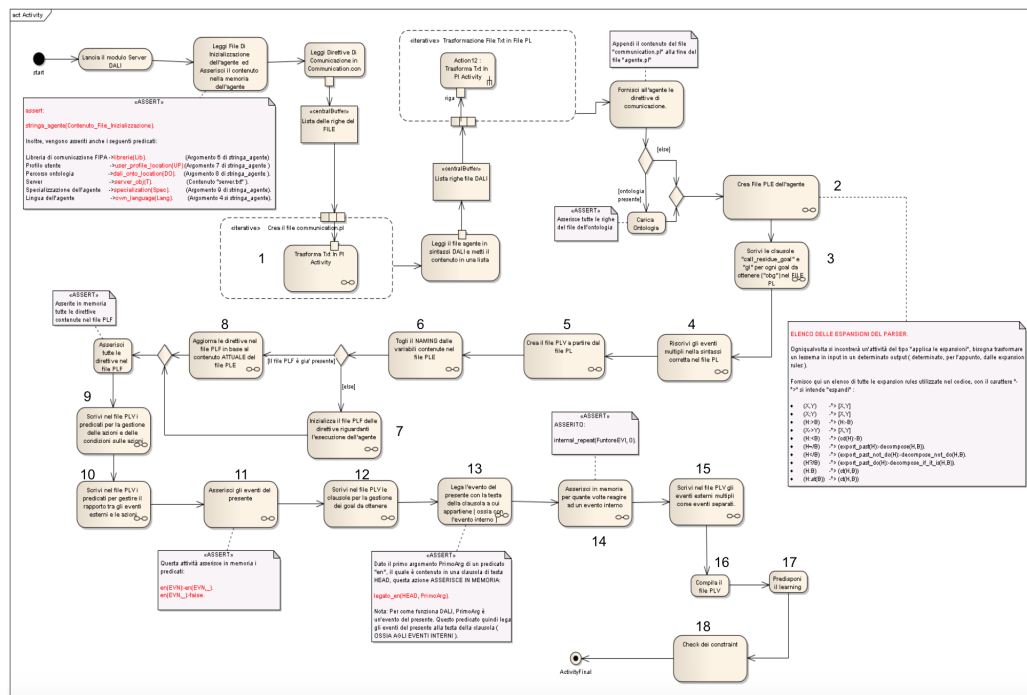
This is the Use Case of DALI Interpreter, where the only external actor to the system is the user. In this Use Case Diagram there are two many divisions that represent the funtionality of the system: the first is the lexical analysis of the DALI program, creation of files pl, plv, plf, ple and inizialization of the agent; the second is interpretations of the agent.

## 2. Activity Diagrams that explain the Use Case Diagram

The two macroscopic division described in the previous Use Case will be shown : **Inizializzazione of the agent (A)** and **Interpret the agent (B)**.



### A. Initialization of the agent



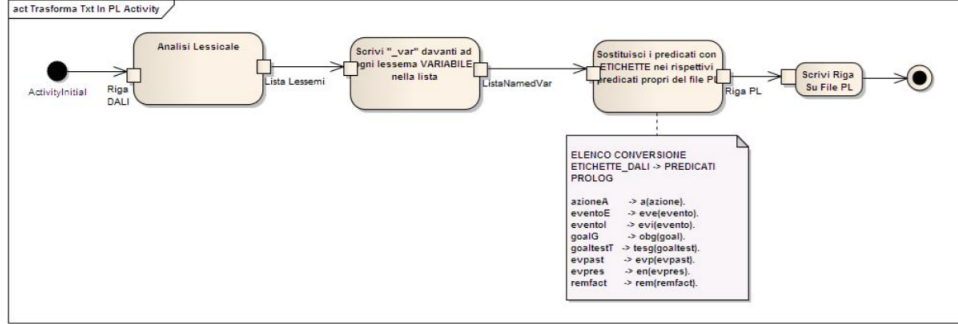
This is the Activity Diagram that explain the Use Case: **“The lexical analysis of the DALI program, creation of files pl, plv, plf, ple and initializzation of the agent”**. Once the file `startmash.sh` is executed, it starts the predicate `start0` from file `active _dali _wi.pl` in line 87. It is read the inizationalization file of the agent and the content is asserted in memory of the agent. The following predicates are asserted also:

- the library of FIPA communications (*Librerie(Lib)*, argument 6 of the string agent in line 91);
- user profile(*user \_profile \_location(UP)*, argument 7 of the string agent in line 91 );
- ontology (*dali \_onto \_location(DO)*, argument 8 of the string agent in line 91);
- the server(*server obj(T)*, that is the content of the file server.txt);

- specializzation of the agent(*specialization(Spec)*, argument 9 of the string agent in line 91);
- language of the agent (*own \_ language(Lang)*, argument 4 of the string agent in line 91).

The directive of communication is read in *communication.con* and a file *communication.pl* (see A.1) is created. The agent file in syntax DALI is read, the content of the file is saved in a list and subsequently the txt-file is transformed in pl-file. The communication directives to the agent are provided (the content of the file *communication.pl* is append to the end of the file *agent.pl*). The ontology is loaded, if it is present, then it creates the ple-file of the agent. The plv-file is created from pl-file. The NAMING of the variables is eliminated in ple-file, line 139. If the plf-file exists (predicate *file \_ exist(FilePlf)*, line 141), then the directives in plf-file are updated, based on the current content of ple-file. Else the plf-file is initialized of the directive about the execution of the agent(line 141). All directives are asserted in plf-file (line 143). The predicates for the management of the actions and conditions of the actions are written in plv-file (line from 154 to 156). The present events are asserted in line 158. The clause for the management of goals to be obtained is written to the plv-file (line 160). The present event is bound with the head of the clause to which it belongs(namely with the internal event). How many times to react to the internal event are asserted in memory (line 164). The multiple external events are written in plv-file as separated events (line 166). The plv-file is compiled (line 168) and the learning is predisposed (line 171), finally the constraints are checked in line 175.

## A.1. Transform the txt file into pl-file



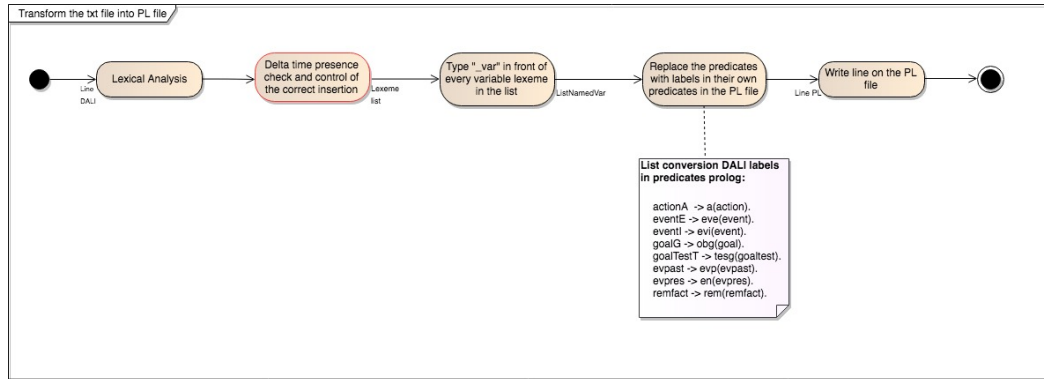
This is a sub-activity of **Inizializzazione of the agent**. The txt-file, writes in DALI syntax, is transformed to the pl-file. In file *active \_dali \_wi.pl*, line 105, the predicate *token(File)* is called. This predicate opens a file called *tokefun.pl* in line 34. From here the lexical analysis starts in line 94: a list of ASCII code is transformed in a list of lexemes. In line 188 it is checked if the character ASCII represents a variable with the predicate *variable(Me)*. If the result is true, the predicate *examine \_variable* is called in line 229, in front of each VARIABLE lexeme in the list it is written “*\_var*”. Subsequently from line 239 to 256 the events label is examined and the predicates with labels are replaced with a prolog predicates. In the following the list of conversion of events label is shown :

- *azioneA* → *a(azione)*.
- *eventoE* → *eve(evento)*.
- *eventoI* → *evi(evento)*.
- *goalG* → *obg(goal)*.
- *goaltestT* → *tesg(goaltest)*.
- *evpast* → *evp(evpast)*.
- *evpres* → *en(evpres)*.

- $remfact \rightarrow rem(remfact)$ .

This predicates are written in pl file.

## A.1.(Updated) Transform the txt file into pl-file



This is a sub-activity of **Inizializzazione of the agent**. The txt-file, writes in DALI syntax, is transformed to the pl-file. This diagram is the updated version of the previously diagram. In this version the verification of the insertion of the Delta time and the control of its correct syntax was introduced. It is shown the changes in file *tokefun.pl*. In line 159 it is asserted *eventi\_esterni(0)* and *assert(deltaT(0))* that will be used later :

```

take_meta(L,F):-assert(eventi_esterni(0)),
assert(deltaT(0)),assert(parenesi(0)),assert(buffer([])),
name(F,Lf),append(Lf,[46,112,108],Lft),
name(Nf,Lft),if(file_exists(Nf),delete_file(Nf),true),
last(L,U),
repeat,
member(Me,L),
examine_all(Me),
Me==U,!,if(clause(residue(R),_),(name(R1,R),
examine_all(R1)),true),
if(clause(buffer(ParsedC),_),
(retractall(buffer(_)),name(Parsed,ParsedC),
open(Nf, append, Stream, []), write(Stream, Parsed),
close(Stream)),(write('Errore take_meta'),nl)),
re_file(Nf).
  
```

In line 186 it is inserted:

```
examine_all1(Me):-if(member(Me,['(',')']),
conta_parentesi(Me),true),
if(tempo(Me), scrittura(Me),
(if(variabile(Me),examine_variable(Me),
if(label(Me),examine_label(Me),write_NovarNoLabel(Me)))))).
```

It is inserted an if-construct in order to check the predicate *tempo(Me)*. If this predicate is true then it is executed the predicate *scrittura(Me)*. From lines 203 to 211 it is inserted:

```
tempo(Me):- name(Me,L), nth0(0,L,E1,L1), E1==116,
numbertime(L1).
numbertime(L1):- nth0(0,L1,E1,L_rest),check_number(E1),
scorri(L_rest).
check_number(E1):- E1>47, E1<58.
scorri(L_rest):-if(L_rest=[],true, scorri_list(L_rest)).
scorri_list(L_rest):- nth0(0,L_rest,X,L2),
check_number(X),scorri(L2).
scrittura(Me):-name(Me,L),nth0(0,L,R,L3),
append(Parsed,L2,Parola),assert(buffer(Parola)),
clause(deltaT(X),true),retractall(deltaT(X)),
assert(deltaT(1)).
```

With the predicate *tempo(Me)*, where *Me* is the *tseconds*, it is checked if the first character is the “t”. If it is so it is executed *numbertime(L1)*, this predicate it verifies if the characters that following the first character “t” are numbers. If the syntax of *tseconds* is correct the predicate *scrittura(Me)* is executed and then it is proceed to write into files .pl and .plv ,relative to the agent, the following predicate:

**deltaT(seconds).**

and finally it is asserted the predicate *deltaT(1)*. In line 243 it is inserted:

```
appE(L,U):- conta_eventi_esterni,length(L,N),
nth1(N,L,U,R),append([101,118,101,40],R,L1),
re_write(L1),assert(evento_aperto).
```

and line 258:

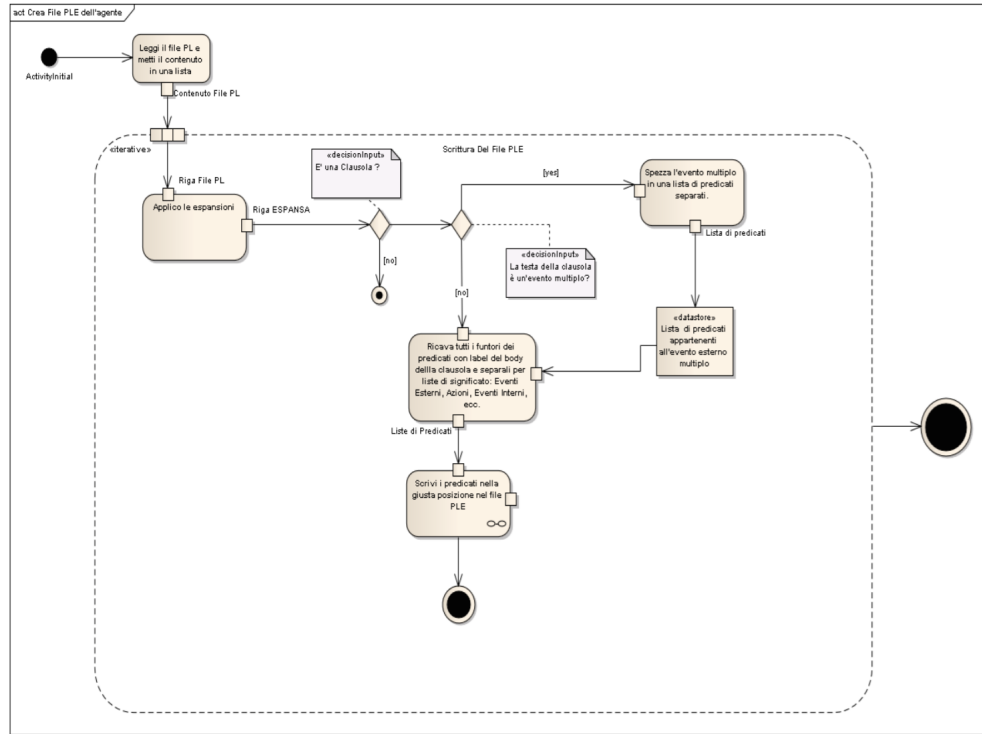
```

conta_eventi_esterni:- clause(eventi_esterni(X),true),
R is X+1, retractall(eventi_esterni(X)),
assert(eventi_esterni(R)).

```

The predicate *conta\_eventi\_esterni* allows to determine the number of the external event. It is introduced this predicate because if the number of the external event is greater than one, the tseconds must be present in the agent's file.

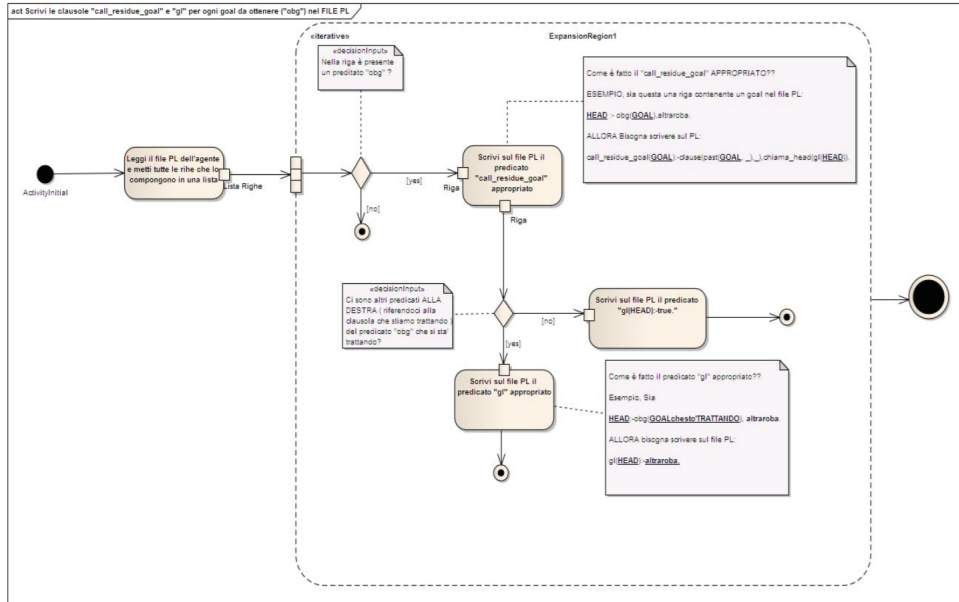
## A.2. Create ple-file of the agent



This is the sub-activity of the **Inizialization of the agent**. In file *active\_dali\_wi.pl*, line 131 the predicate *aprifile(FilePl)* is called, the pl-file is read and its content is put in a list. The predicate in line 215 is called, the expansion is applied. If the head of clause isn't a multiple event, all the functors of the predicates with the label are derived and are separated by list

of meaning. Finally the predicates are written in a right position in ple-file. If the head of clause is a multiple event, first the multiple event is broken into a list of separate predicates and then all the functors of the predicates with the label are derived and are separated by list of meaning. Finally the predicates are written in a right position in ple-file.

### A.3. Write the clause “ call \_residue \_goal ” and “gl” for each goal to get(“ obg ”) in the pl-file



This diagram represent a sub-activity of the **Inizialization of the agent**. In the file *active \_dali \_wi.pl*, line 133 the predicate *aprifile \_res(FilePl)* is called and the predicates for the manage to the “obtaining goals” are written in pl-file. The predicate is called in file *memory.pl*, line 12. The pl-file of the agent is read and all the lines that compose it are put in a list. In file *memory.pl*, line 41 it is checked if the predicate “obg” is present in the line. If this predicate is present, the predicate “call \_residue \_goal” is written in pl-file ,line 49. How is an appropriate “call \_residue \_goal” made?

**Example:** Let this a line containing a goal in pl-file:

$$HEAD : -obg(GOAL), other.$$

Then it must write on pl-file:

$$\begin{aligned} call\_residue\_goal(GOAL) : & -clause(past(GOAL, \_), \_), \\ chiama\_head(gl(HEAD)). \end{aligned}$$

Now if are there any other predicates to the right of the predicate “obg” being dealt with? If yes the predicate “gl” appropriate is written in pl-file .

How is an appropriate predicate “gl” made?

**Example:** Let

$$HEAD : -obg(GOALthatIsBeingDealtWith), other.$$

Then it must write on pl-file:

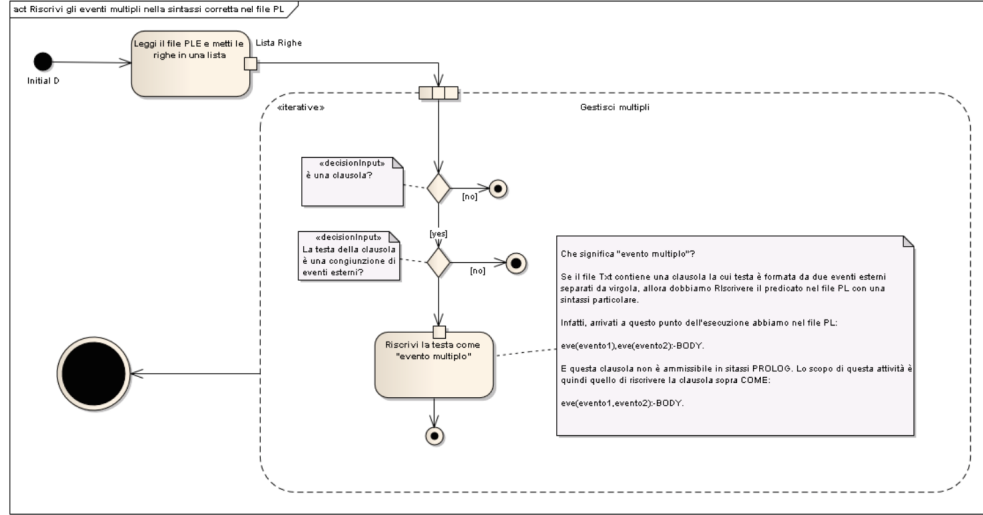
$$gl(HEAD) : -other.$$

If there aren't predicates to the right of the rule then this predicate is written in pl-file:

$$gl(HEAD) : -true.$$



## A.4. Rewrite multiple events in the correct syntax in the pl-file



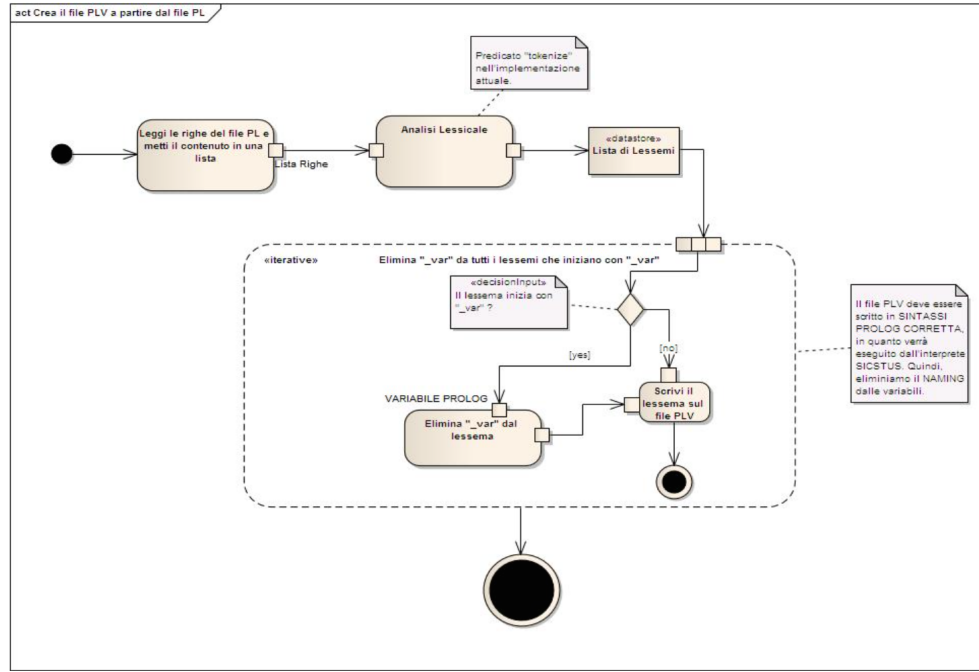
This diagram represent a sub-activity of the **Inizialization of the agent**. In the file *active \_dali \_wi.pl*, line 135 the predicate *carica \_file(FilePl)* is called and *leggi \_mul.pl* is opened the file in line 7. The ple-file is read and its content is put in a list. If it is a clause, then is the head of the clause a conjunction of external events? If yes, the head is rewritten as multiple event. What means multiple event? If the txt-file contains a clause whose head is formed by two external events, comma separated, then it is necessary rewrites this predicate in pl-file with a particular syntax. In fact at this point of the execution in pl-file it is:

$$eve(evento1),eve(evento2) : -BODY.$$

And this clause is not admissible in PROLOG syntax. The purpose of this activity is to rewrite the above clause as :

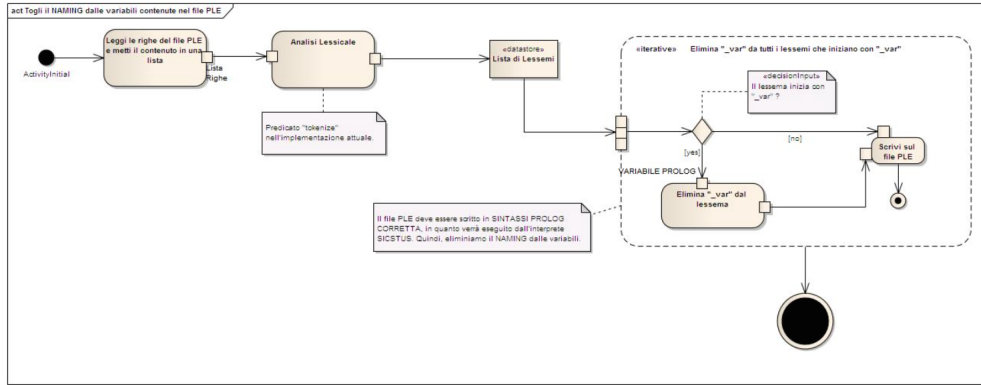
$$eve(evento1,evento2) : -BODY.$$

## A.5. Create the plv-file from pl-file



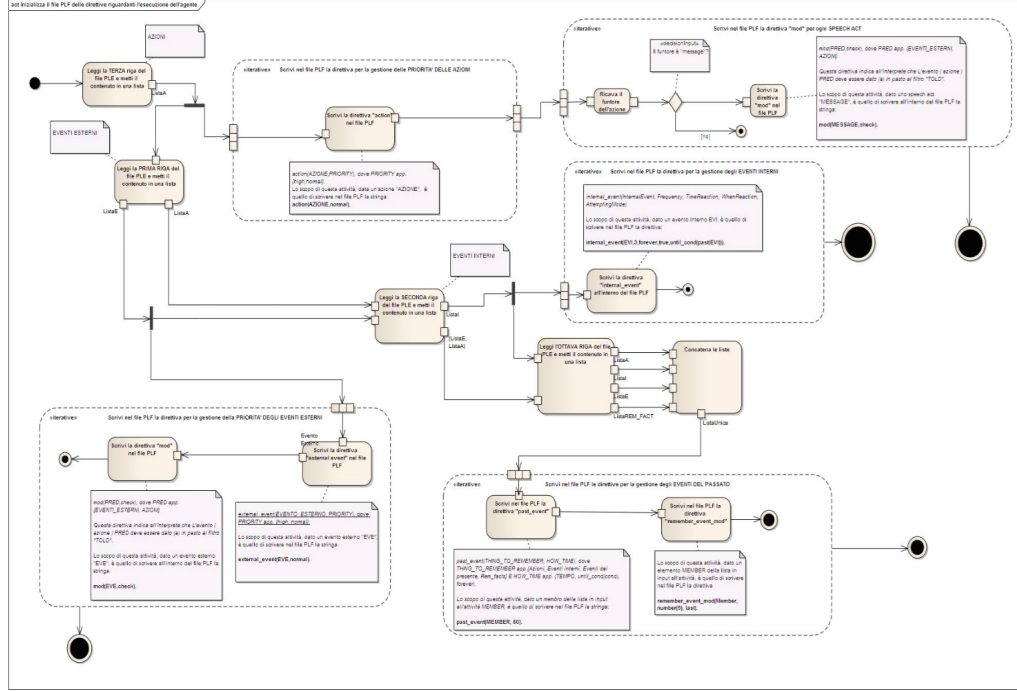
This diagram represent a sub-activity of the **Inizialization of the agent**. In the file *active \_dali \_wi.pl*, line 137 the predicate *togli \_var(F)* is called, and the file *togli \_var.pl* is opened in line 7. The ple-file is read and its content is put in a list. It is started the lexical analysis with predicate “tokenize”. Given the list of lexemes, “\_var” is eliminated from all the lexemes that begin with “\_var”. If the lexeme begins with “\_var”, file *togli \_var.pl* line 88, then *\_var* it is eliminated from lexeme and this is written in plv-file. Else it is written directly in plv-file. The plv-file must be write in a correct Prolog syntax because it will execute by SICSTUS Interpreter. Then the NAMING of the variable is eliminated.

## A.6. Eliminate the NAMING of the variable in ple-file



This diagram represent a sub-activity of the **Inizialization of the agent**. In the file *active\_dali\_wi.pl*, line 139 the predicate *togli\_var\_ple(Fe)* is called, and the file *togli\_var.pl* is opened in line 20. This predicate deals to eliminating the prefix “*var \_*” from variables to which the naming has been applied. The lines of ple-file is read and its contents is put in a list. It is started the lexical analysis with predicate “tokenize”. Given the list of lexemes, “\_var” is eliminated from all the lexemes that begin with “\_var”. If the lexeme begins with “\_var”, then \_var is eliminated from lexeme and this is written in ple-file. Else it is written directly in ple-file. The plv-file must be write in a correct Prolog syntax because it will execute by SICSTUS Interpreter. Then the NAMING of the variable is eliminated.

## A.7. Initialize the plf-file of the directive about the execution of the agent



This diagram represent a sub-activity of the **Inizialization of the agent**. It is in file *active\_dali\_wi.pl*, this diagram describes the else branch in the if-control in line 141. The predicate *“inizializza\_plf(FilePle)”* is called in line 454. In line 455 the first line in ple-file(external events) is read and the content is put in a list(Le), the second line in ple-file(internal events) and the content is put in a list(Li), the third line in ple-file( actions) and the content is put in a list(La) and the octave line in ple-file and puts the content in a list(Lr). This four lists are concatenated and there is only one list in output.

### A.7.1. Write in plf-file the directive for the management of external events

From line 468 to 479 the directives of PAST EVENTS are written in plf-file. The directive *“past\_event”* is written in plf-file. The purpose of this activity

is to write the string in the plf file:

**past \_event (MEMBER, 60)**

where MEMBER is a member of a list. Furthermore the directive “*remember \_event \_mod*” is written in plf-file. The purpose of this activity is to write the string in the plf-file:

**remember \_event \_mod (Member, number(5), last)**

where Member is a member of a list.

#### **A.7.2. Write in plf-file the directive for the management of action’s priority**

In line 500 the third line of ple-file is read and it is put the content in a list, and the directive “ action ” is written in plf-file. The purpose of this activity is to write the string in the plf-file:

**action (ACTION, priority)**

where priority could be high or normal. In line 609 with the predicate “*writ \_actn \_event* ” it is checked that if the functor is equal to message, the directive “mod ” is written in plf-file. The purpose of this activity, given a speech act “MESSAGE”, is to write the string in the plf-file:

**mod (MESSAGE, check).**

#### **A.7.3. Write in plf-file the directive for the management of external event’s priority**

In line 513 the first line of ple-file is read and it is put the content in a list, and it is written the directive “ mod ” in plf-file. The purpose of this activity, given an external event “EVE”, is to write the string in the plf-file:

**mod (EVE, check).**

In line 596 with the predicate “*writ \_est \_event* ” it is written the directive “external event ”. The purpose of this activity, given an external event “EVE”, is to write the string in the plf file:

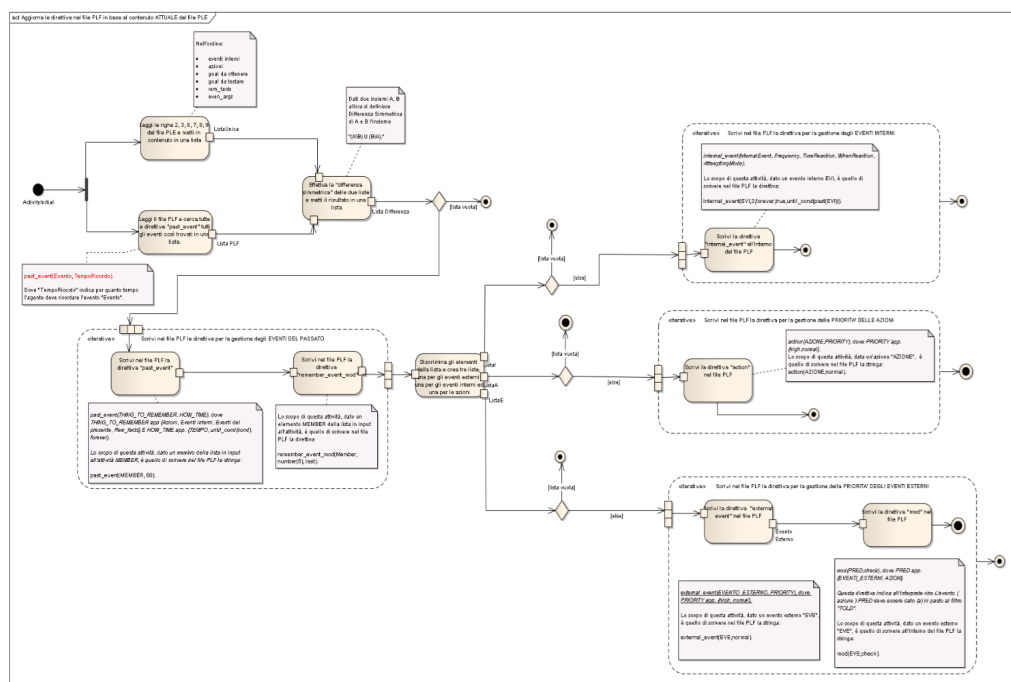
**external \_event (EVE, normal).**

#### A.7.4. Write in plf-file the directive for the management of internal events

In line 459 the predicate “*smon*<sub>ta</sub> \_ *interno*” is called in line 483. The second line of ple-file is read and the content is put in al list, and the directive “*internal \_ event*” is written in plf-file. The purpose of this activity, given an external event “EVE”, is to write the string in the plf file:

```
internal  event(EVI, 3, forever, true, until  cond(past(EVI))).
```

#### A.8. Updated the directives in plf-file based on the current content in ple-file



This diagram represent a sub-activity of the **Inizialization of the agent**. It is in file **active \_dali \_wi.pl**, this diagram describes the then-branch in the if-control in line 141. The predicate “*controlla \_ev \_all(F)*” is called in line 529. The line 2(internal event), line 3(actions), line 6(goal to get), line 7(goal to test), line 8(rem facts), line 9(event args) of ple-file are written

and the content is put in a list(L6). In line 542 the plf-file is read and all the directives “ past \_event” is searched and the events are put in a list(Ls) in line 549. In line 550 the symmetric difference among the two lists is carried out and the result is put in a list(Ldf).

#### **A.8.1. Write in plf-file the directive for the management of past events**

If this list ldf is not empty the directive for the management of the past event is written in the plf-file. The directive “ *past \_event*” is written. Given a member of the input list, the purpose of this activity is to write the string in the plf-file:

**past \_event(MEMBER, 60).**

Subsequently the directive “ remember \_event \_mod” is written in plf-file line 582. Given a member of the input list, the purpose of this activity is to write the string in the plf-file:

**remember \_event \_mod(MEMBER, number(5),last).**

Discriminate the element of the list, three lists are created: one for internal events, one for actions, one for external events.

#### **A.8.2. Write in plf-file the directive for the management of internal events**

The directive “ internal \_event” is written in plf-file, from line 603 to 607. Given an internal event EVI, the purpose of this activity is to write the string in the plf-file:

**internal \_event(EVI, 3,forever, true,until \_cond \_past(EVI)).**

#### **A.8.3. Write in plf-file the directive for the management of action’s priority**

The directive “ action ” is written in plf-file, line 609. Given an action, the purpose of this activity is to write the string in the plf-file:

**action(ACTION, priority).**

where priority could be high or normal.

#### A.8.4. Write in plf-file the directive for the management of external event's priority

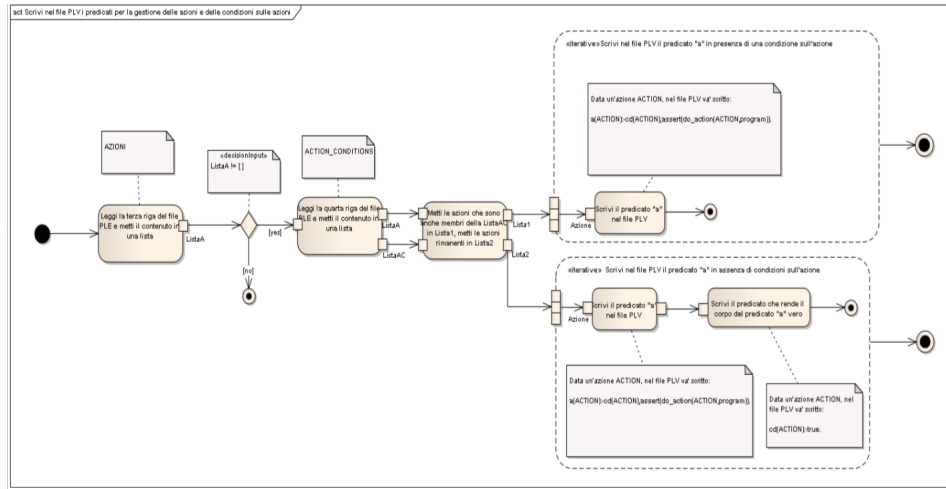
The directive “ external event ” is written in plf-file, line 596. Given an external event EVE, the purpose of this activity is to write the string in the plf-file:

**external \_event(EVE, normal).**

where priority could be high or normal. The directive “ mod ” is written in plf-file, line 600. Given an external event EVE, the purpose of this activity is to write the string in the plf-file:

**mod(EVE, check).**

#### A.9. Write in plv-file the predicates for the management of actions and action's conditions



In line 791 the third line of ple-file is read and the content is put in a list(X). If X is not empty then the fourth line in ple-file is read and it is put the content in a list (line 795). The actions that are also members of the list X are puts in list1 and the rimanent actions are puts in list2.



### A.9.1. Write in plv-file the predicate in the presence of a condition on the action

Line 810, predicate *clause1*: given an action ACTION , it is written in plv-file:

$$a(ACTION) : -cd(ACTION), assert(do\_action(ACTION, program)).$$

### A.9.2. Write in plv-file the predicate in the absence of a condition on the action

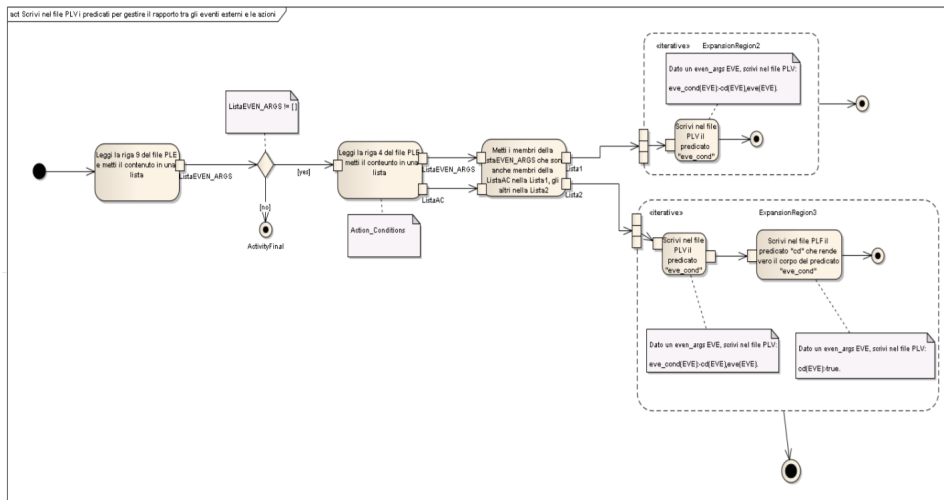
Line 815, predicate *clause2*: given an action ACTION , it is written in plv-file:

$$a(ACTION) : -cd(ACTION), assert(do\_action(ACTION, program)).$$

and it is written the predicate that makes true the body of predicate “a”.  
Given an action ACTION , it is written in plv-file:

$$cd(ACTION) : -true.$$

## A.10. Write in plv-file the predicates to manage the relationship between external events and actions



Line 825, predicate *cond \_esterni*. The line 9 of ple-file is read and the content is put in a list. If the list event \_args is not empty (line 828), the fourth line in ple-file is read and the content is put in a list, action \_condition list . It puts the member of list event \_args that are also member of action \_condition list in list 1, others in list 2. In line 846, predicate *clause1 \_esterni*. The predicate “ eve \_cond ” is written in plv-file. Given a even \_args EVE, in plv-file it is written:

$$eve\_cond(EVE) : -cd(EVE), eve(EVE).$$

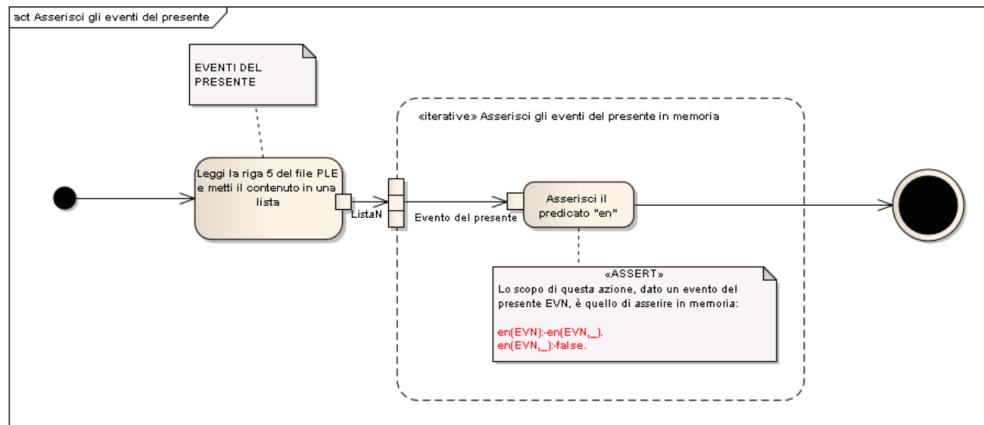
In line 851, predicate *clause2 \_esterni*. The predicate “ eve \_cond ” is written in plv-file. Given a even \_args EVE, in plv-file it is written:

$$eve\_cond(EVE) : -cd(EVE), eve(EVE).$$

Given a even \_args EVE, in plv-file it is written:

$$cd(EVE) : -true.$$

## A.11. Assert the present events

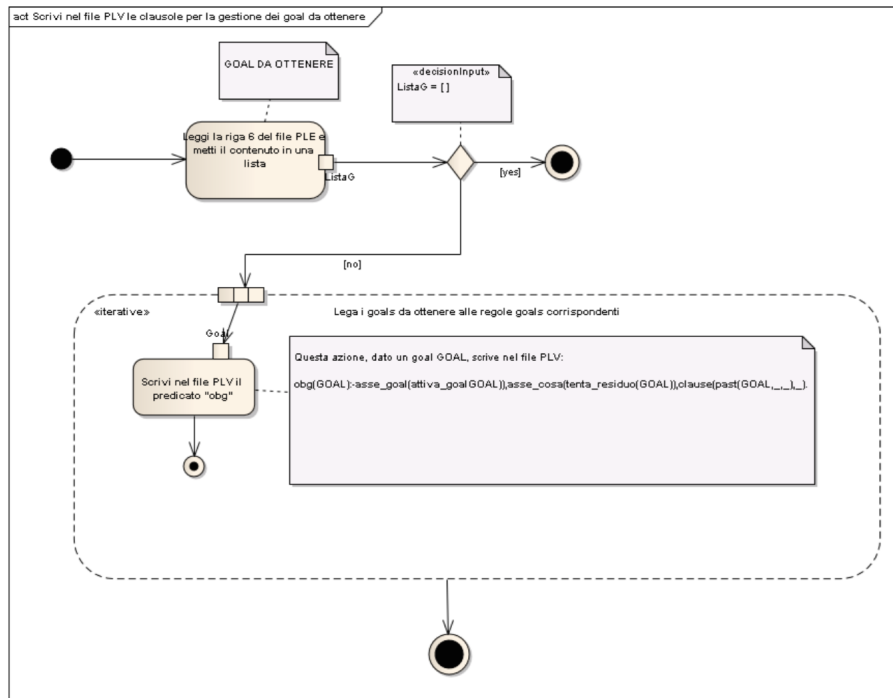


Line 860, predicate *assert\_evN*. The fifth line is read in ple-file and the content is put in a list. The predicate “en” is asserted. Given an present event EVN the purpose of this activity is to write the string in the plf-file:

**en(EVN):-en(EVN,\_).**

**en(EVN,\_):-false.**

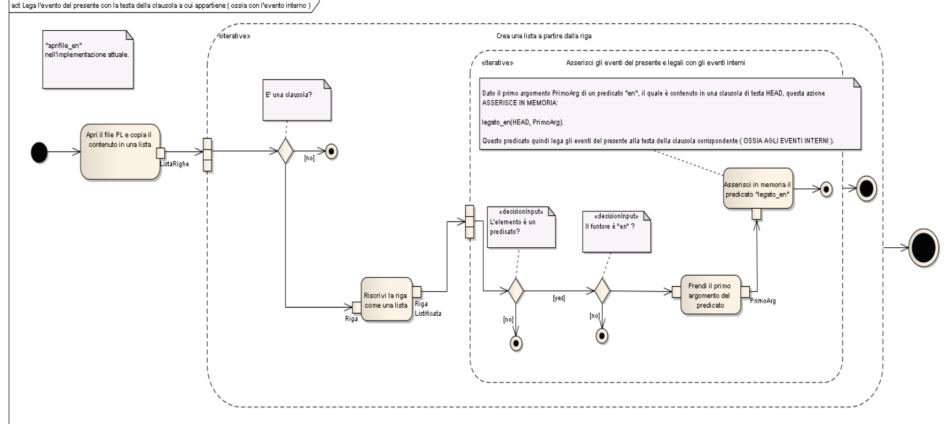
## A.12. Write in plv-file the clause to manage of goals to obtain



Line 874, predicate *obg\_goal*. The line 6 of the ple-file is read and the content is put in a list, list G. If this list is not empty then it is written in plv-file the predicate “ obg ”. This action, given a goal GOAL, is written in plv-file, line 887:

*obg(GOAL) : -asse\_goal(attiva\_goal(GOAL)),*  
*asse\_cosa(tenta\_residuo(GOAL)),clause(past(GOAL,\_,\_),\_).*

### A.13. Bind the present event with the head of the clause to which it belongs

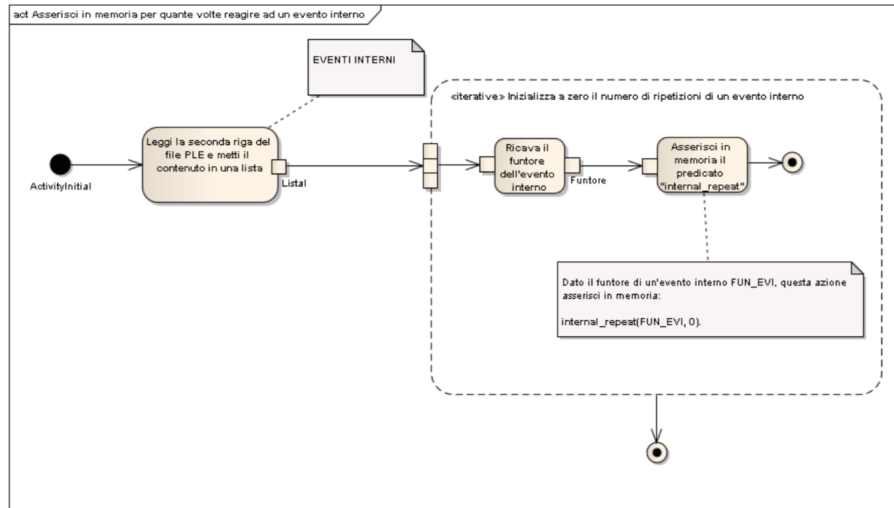


Line 162 the predicate “ `aprifile _en` ” is called in line 658. The pl-file is opened and the content is put in a list, this is a list of line. If a line is a clause , it is rewrite the line in a list. If the element is a predicate, and the functor is equal to “en” then it takes the first argument of the predicate and the predicate “ `legato _en` ” is asserted in memory . Given a first argument `PrimoArg` of a predicate “ `en` ”, it asserts in memory :

$$\text{legato\_en}(\text{HEAD}, \text{PrimoArg}).$$

This predicate binds the present events with the head of the corresponding clause.

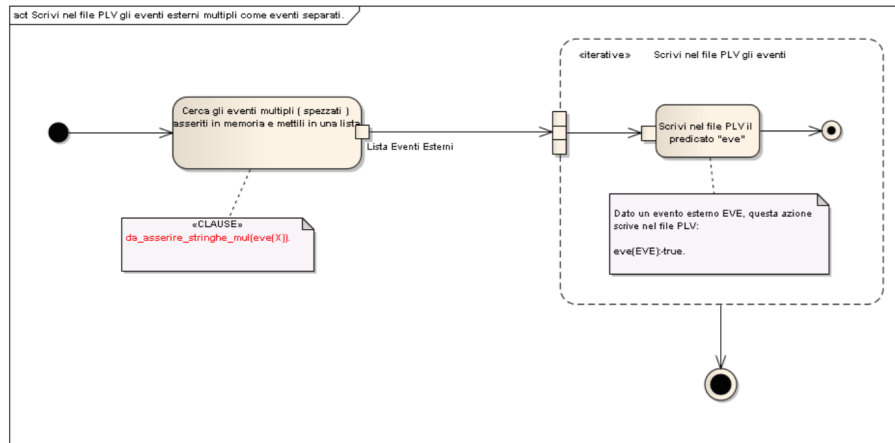
## A.14. Assert in memory how many times to react to an internal event



This diagram represent a sub-activity of the **Inizialization of the agent**. Line 164, the predicates “ ass \_internal \_repeat” is called in line 1554. The second line in ple-file is read and the content is put in a list. The functor of the internal event is obtained and the predicate “ internal \_repeat ” is asserted in memory . Given a functor of internal event FUN \_EVI, this action asserted in memory :

**internal \_repeat (FUN \_EVI,0)**

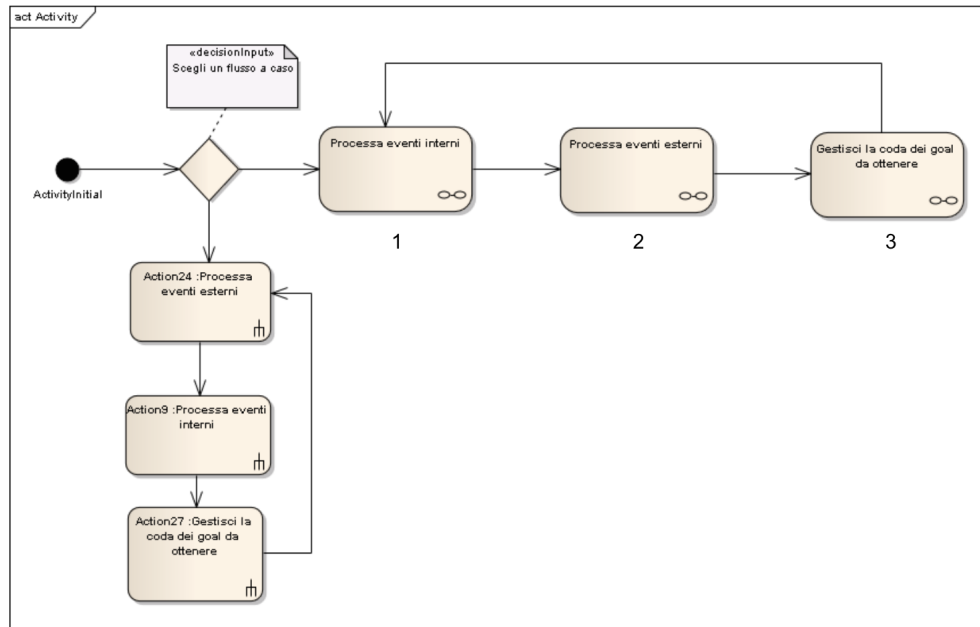
## A.15. Write in plv-file the multiple external events as separate events



This diagram represent a sub-activity of the **Inizialization of the agent**. Line 166 the predicate “ ass \_stringhe \_mul” is called in line 1566. The multiple events asserted in memory are created and it put this events in a lists (clause(da \_asserire \_stringhe \_mul)). The predicate “eve” is written in the plv-file. Given an external event EVE, this action writes in plv-file (line 1573):

**eve(EVE):- true.**

## B. Interpret the agent



This is a macro-activity associated at the Use Case “ **Interpret the agent**”. After the initialization the agent is active. In line 177 there is a predicate “go”, this predicate is called in line 1510.

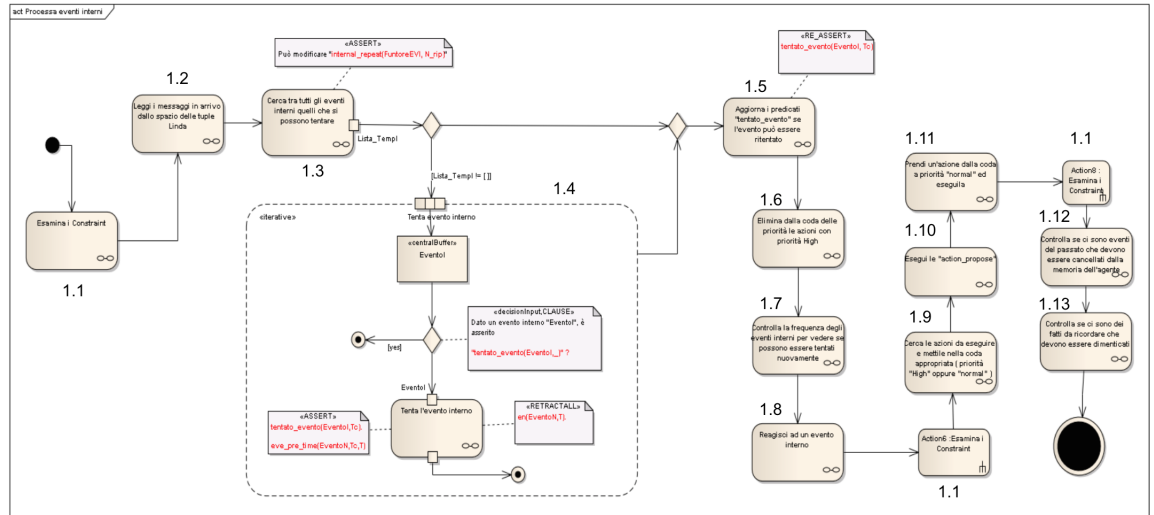
```
go:-repeat, (pari,side_goal ), fail,!.
```

The predicate “ pari ”, that represent this diagram substantially, is called in lines 1499-1500.

```
pari:-sleep(1), random(1,10,R), R1 is R mod 2,
if(R1=0,(internal,external), (external,internal)).
```

The branch (internal,external) will be analyzed then “ **Processa eventi interni** ” (B.1), “ **Processa eventi esterni** ” (B.2), “**Gestisci la coda dei goal da ottenere** ” (B.3). Each of these activities will be described with activities diagram.

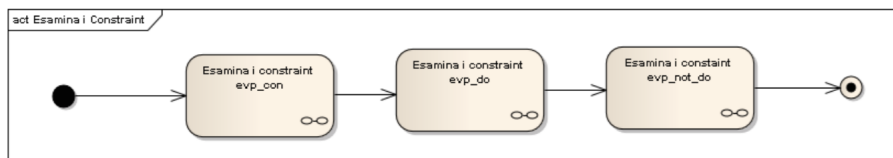
## B.1. Process Internal Events



This diagram is represented by the predicate “ **internal** ” in line 1502.

```
internal:- blocco_constr, ricmess, ev_int, ev_goal,
ev_int0, blocco_numero_ev_int, blocco_frequenza, ev_int2,
controlla_freq_tent, svuota_coda_priority,
controlla_freq_iv, scatena, blocco_constr, keep_action,
execute_do_action_propose, prendi_action_normal,
blocco_constr, controlla_vita.
```

### B.1.1. Examine the Constraints



This diagram is a sub-activity of the activity “**Process internal events**” and it is represented in line 1513 by the predicate “ blocco \_constr ” :

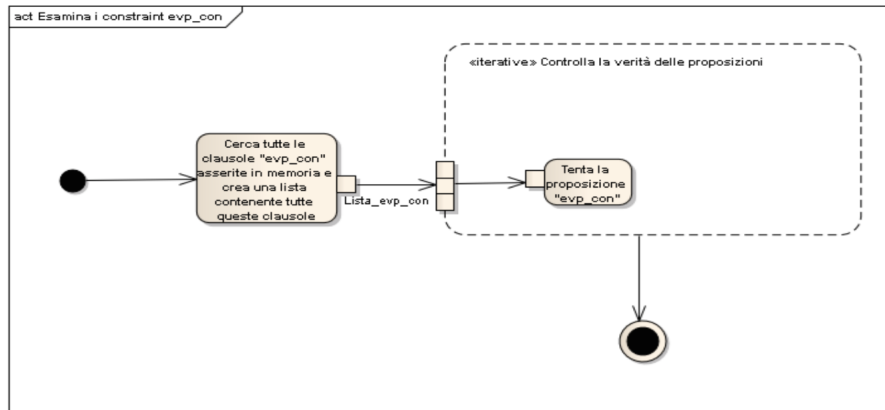


```

blocco_constr:- evaluate_evp_constr, evaluate_evp_do,
evaluate_evp_not_do.

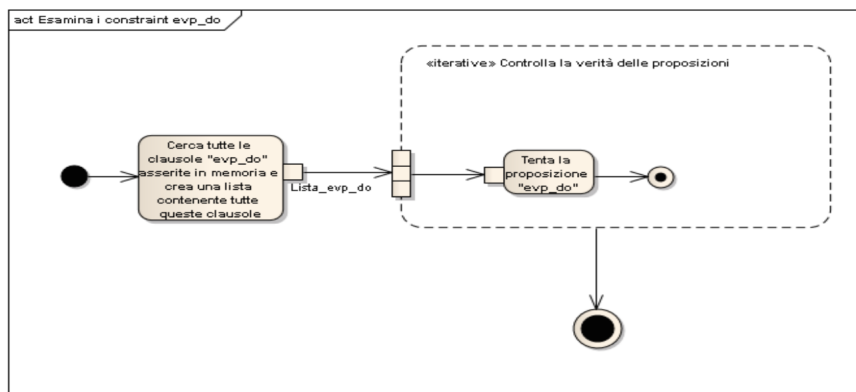
```

#### B.1.1.1. Examine the Constraints evp \_con



This diagram is sub-activity of “ **Examine the Constraint** ” and it represent to predicate “ **evaluate \_evp \_constr** ”. This predicate is called in file **examine \_past \_constraints.pl**, line 128. All the clause “evp \_con”, asserted in memory, is searched and a list with all these clause is created. Subsequently it is verified the veracity of propositions and the proposition “evp \_con” is attempted.

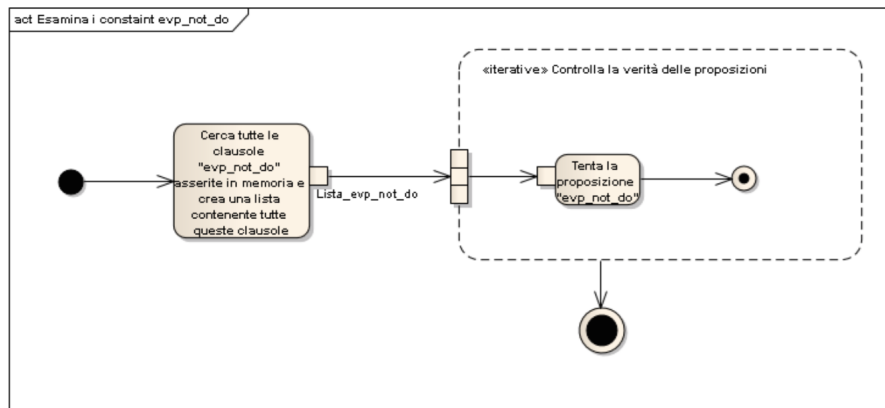
#### B.1.1.2. Examine the Constraints evp \_do



This diagram is sub-activity of “ **Examine the Constraint** ” and it represent to predicate “ **evaluate \_evp \_do** ”. This predicate is called in

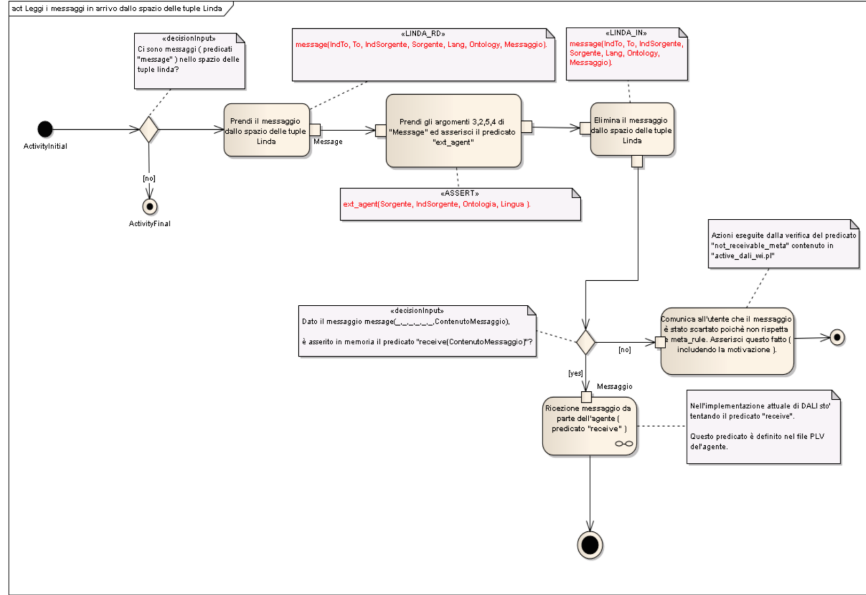
file **examine \_past \_constraints.pl**, line 324. All the clause “evp \_do”, asserted in memory, is searched and a list with all these clause is created. Subsequently it is verified the veracity of propositions and the proposition “evp \_do ” is attempted.

### B.1.1.3. Examine the Constraints evp \_not \_do



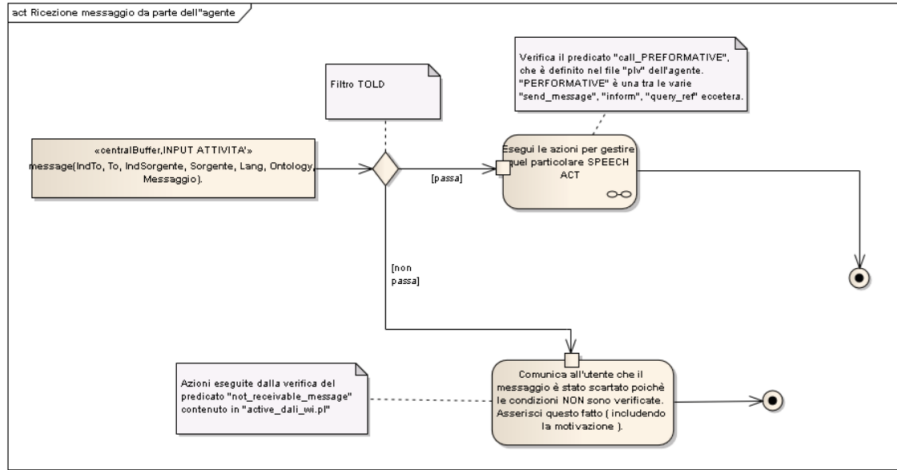
This diagram is sub-activity of “ **Examine the Constraint** ” and it represent to predicate “ evaluate \_evp \_not \_do ”. This predicate is called in file **examine \_past \_constraints.pl**, line 226. All the clause “evp \_not \_do”, asserted in memory, is searched and a list with all these clause is created. Subsequently the veracity of propositions is verified and the proposition “evp \_not \_do” is attempted.

### B.1.2. Read the arrived message from the space of Linda's tuple



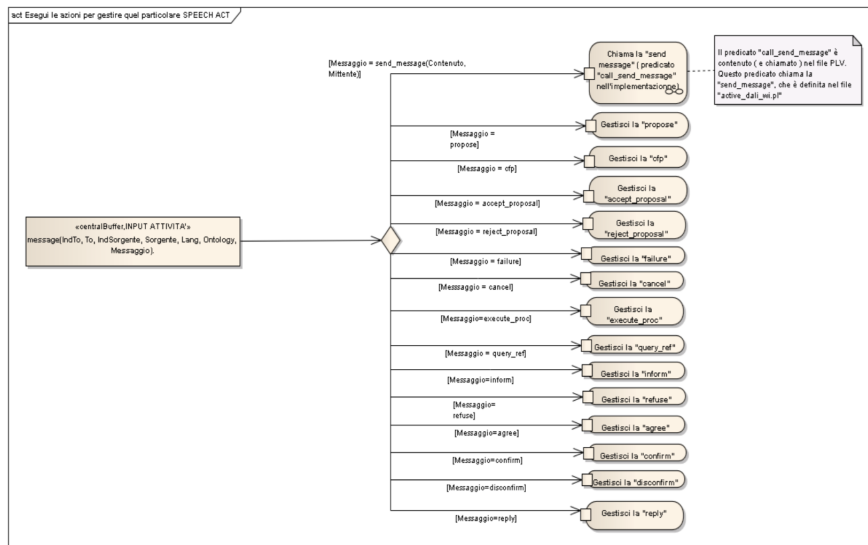
This diagram is a sub-activity of the activity “**Process internal events**” and it is represented in line 684 by the predicate “ricmess”. The question is: are there message( message predicates) in the space of Linda’s Tuple? If yes, in line 685 the predicate “rd \_noblock” is called in file **client.pl** line 118(under the directory PrologReference/library/linda). Subsequently the predicate “ricmess0” is called in line 689. It is taken the arguments 3,2,5,4 of Message and in line 690 calling to the predicate “ assecosa(ext \_agent(AgM, IndM, Ontology, Language))” where the predicate “ ext \_agent ” is asserted. In line 691 the predicate “in \_noblock” is called in file **client.pl** line 138(under the directory PrologReference/library/linda), where the message of the space of Linda’s Tule is eliminated. Line 693, given a message , is the predicate “ receive(ContenutoMessaggio) ” asserted in memory? If yes, the agent receives the message( predicate “ chiama \_con ”,line 695 - see diagram B.1.2.1). Else it is communicated to the agent that the message has been discarded as it does not respect the meta \_rule (predicate “ not \_receivable \_meta ”, line 693-740).

### B.1.2.1. Receiving the message from the agent



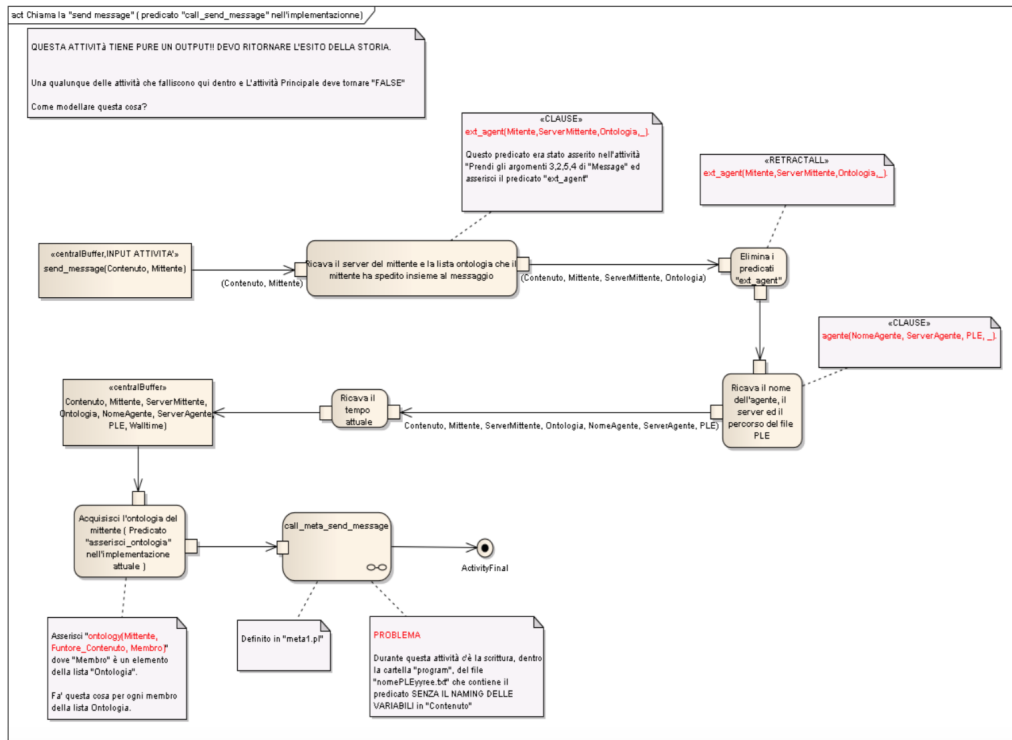
This diagram is a sub-activity of the activity “**Read the arrived message from the space of Linda’s tuple**” and it is represented in line 695 by the predicate “ chiama \_con ”. The TOLD filter is activated in the plv-file. If the filter has the positive result, the actions for manage the SPEECH ACT are executed( see diagram B.1.2.1.1.). Else it is communicated to the agent that the message has been discarded as it does not respect the meta \_rule (predicate “ not \_receivable \_meta ”, line 740).

#### B.1.2.1.1. The actions for manage the SPEECH ACT are executed



This diagram is a sub-activity of the activity “**Receiving the message from the agent**”. The predicate received in the plv-file can be contain various types of messages: propose, cfp, accept \_proposal, reject \_proposal, failure, cancel, execute \_proc, query \_ref, inform, refuse, agree, confirm, disconfirm, reply or send \_message(Content, Sender). The predicate “Call \_send \_message” is contained in plv-file. This predicate calls the “send \_message” line 698(B.1.2.1.1.1).

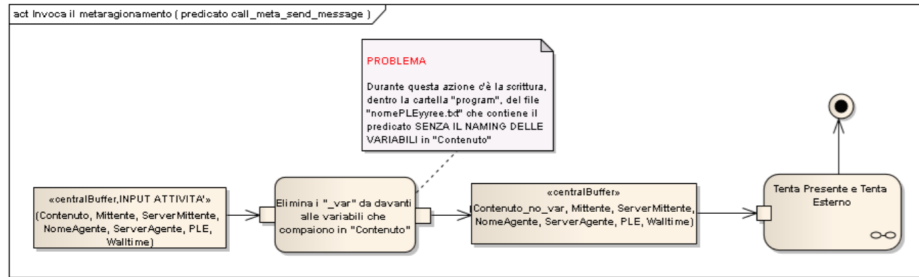
#### B.1.2.1.1.1 Call the “send message ”



This diagram is a sub-activity of the activity “**The actions for manage the SPEECH ACT are executed**”. Line 698, predicate “send \_message”, the sender’s server and the ontology that the sender has sent are derived(line 698). In line 699, the predicates “ext \_agent” is eliminated (retractall ext \_agent). Line 700, the name of the agent, the server and the path of ple-file are derived. In line 706 the time is derived with the predicate “statistics”.

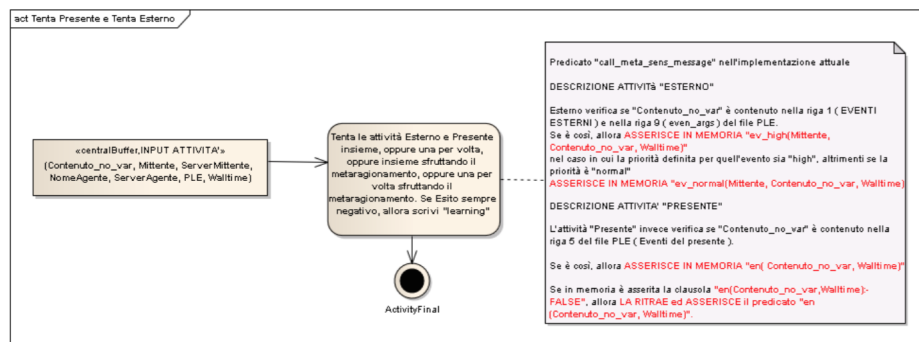
Line 702-729 the ontology of the sender is derived (predicate “asserisci \_ontologia”). In line 709 the predicate “call \_meta \_send \_message” is called. This predicate is defined in B.1.2.1.1.1.1.

#### B.1.2.1.1.1.1. Invoke the meta-reasoning



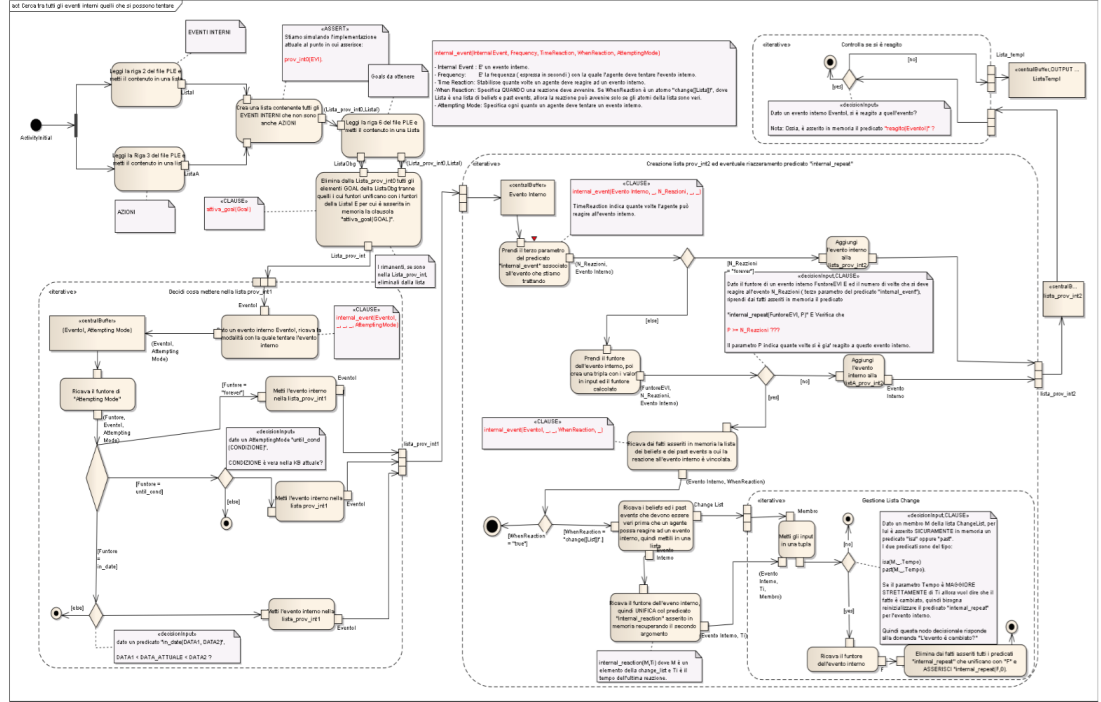
This diagram is a sub-activity of the activity “Call the send message”. The predicate “call \_meta \_send \_message ” is called in file **meta1.pl**, line 27. The “\_var ” is eliminated from the front of variables that appear in “Content ”. Subsequently the predicate “chiama \_send \_message” is called in line 33. Finally present and external activities are attempted (see diagram B.1.2.1.1.1.1.1.).

#### B.1.2.1.1.1.1.1. Present and external activities are attempted



This diagram is a sub-activity of the activity “Invoke the meta-reasoning”. The present and external events are attempted together, or one at a time, or together exploiting the meta-reasoning, or one at a time exploiting the meta-reasoning. If the result is negative, it is written “learning”(file meta1.pl, from line 29 to 53).

### B.1.3. Search among all the internal events those that can be tried



This diagram is a sub-activity of the activity “**Process internal events**” and it is represented in line 1041 by the predicate “**ev \_int**”. The second line in ple-file is read and the content is put in a list. If this list is not empty then the third line of ple-file is read. The list that contains all the external events that do not the actions also, from line 1045 to 1064. Line 1502 the predicate “**ev \_goal**” is called in line 1066 (clause(prov \_int0(Me))). Line 1068 predicate “ev \_goal(L)”: the sixth line of ple file is read. The predicate “clause(obtgoal(X),true)” is called in line 1070. If the list X is empty the predicate “passa \_a \_prov \_int” is executed, else the predicate “ev \_goal21(X,L)” is called, where X is the list obtained from the file and L is the list containing the internal events. All the elements goal present in ListObg are eliminated from the list L, except those functors unify with functors of the ListI for which the clause “attiva \_goal(GOAL)” is asserted

in memory( line 1087). Line 1502, the predicate “ev \_int0” is called in line 1099. It is arrived to the line 1105 where the predicate “internal \_event(Me, \_, \_, \_, C), \_)” is called ( given an internal event, it is obtained the modality with which to try the internal event). In line 1106 the predicate “functor(C, F, \_)” is called and the functor of “Attempting mode” is obtained and the controls on the functors are started.

If the functor is equal to FOREVER, line 1153 is executed and the internal event is put in list *prov \_int1*.

If the functor is equal to UNTIL \_COND the predicate “esamina \_cd”, line 1112, is executed. If the condition is true in the current KB then line 1153 is executed and the internal event is put in list *prov \_int1*. Else the predicate “scarica(Me)” is executed in line 1155.

If the functor is equal to IN \_DATE, the predicate “esamina \_dt” is executed (from line 1115 to 1149 it checks on time is carried out) and the internal event is put in list *prov \_int1*. Else a message of error is written, line 1111.

In the second block a new list, *prov \_int2*, is created and there is a possible redirection to the predicate “internal \_repeat”.

The predicate “**blocco \_numero \_ev \_int**” , line 1502, is called in line 1159. If the list *prov \_int1* is not empty the predicate “blocco \_numero \_2” is called in line 1163. In line 1166 the third value of internal event, if the functor is equal to FOREVER, the event is added to the new list *prov \_int2*, line 1167. Else the predicate “check \_time \_rep” is executed in line 1170.

Given a functor of internal event F and the number of times you must react to the event N \_Reaction. The question is :  $P \geq N\_Reaction?$  , where P is the number of times that it is reacted at the internal event. If the answer is no, the event is added to the list *prov \_int2*, else the predicate “check \_cond \_repeat” is called in line 1177. The fourth parameter is derived ( $C_n$ , WhenReaction). It is derived the list of beliefs and past events to which the reaction to the internal event refers.

If the  $C_n$  is true then end, else the predicate “check \_cond \_repeat1” is called, line 1180.

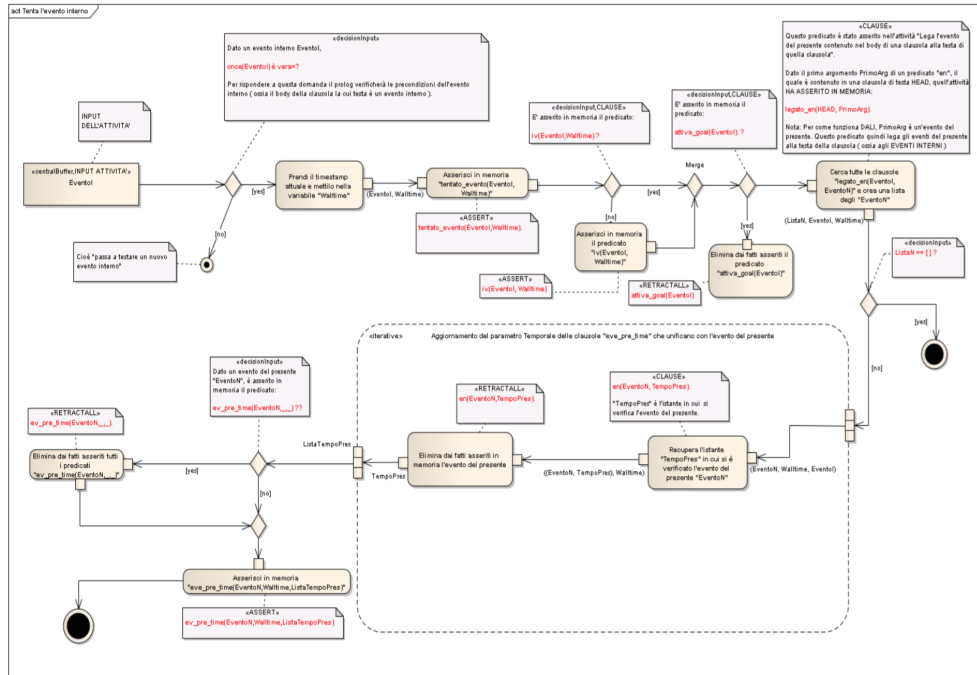
If the  $C_n$  is equal to change, it is derived the list of beliefs and past events that



must be true before the agent reacts to the internal event, and puts them on a list. If this list is not empty in line 1184 the predicate “check \_cond \_change” is called: the functor is derived by the internal event then UNIFIES with the predicate “internal \_reaction( $F, T_i$ )” (where  $F$  is an element of changeList and  $T_i$  is the time of the last reaction), recovering the second argument. Line 1197, predicate “ clause(isa( $M, \_, \_$ ))” is called. Given a member  $M$  of changeList, for this it is asserted in memory a predicate “isa” or “past”. If the time parameter is greater than  $T_i$ , then the fact is changed and it has to initialize the predicate “internal \_repeat”. Line 1193 the predicate “fact \_evp \_changed” is called: all predicates “internal \_repeat” that unifies with  $F$  are eliminated from asserted facts, and it is asserted the predicate “internal \_repeat( $F, 0$ )”.

In the third block the predicate “blocco \_frequenza”, line 1502, is called. The input is the list *prov \_int2*. In line 1212 the predicate “blocco \_frequenza1” is called, if the result of the findall,  $L$ , is not empty the predicate “blocco \_frequenza2” is called, and the clause(reagito( $Me, \_$ )) is called. Given an internal event  $I$ , did you react to that event? If yes it is ended, else the list ListaTempI is the output.

#### **B.1.4. Try the Internal Event**

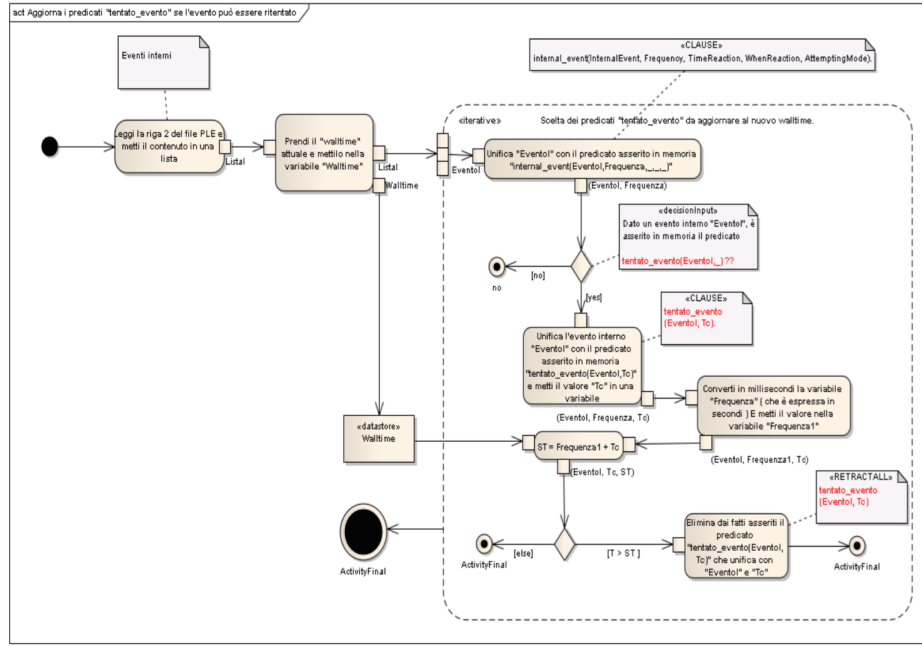


This diagram is a sub-activity of the activity “**Process internal events** ” and it is represented in line 1225 by the predicate “ **ev \_int2** ”. Given an internal event, is once(EventoI) true? If no a new internal event is tempted. Else the current time-stamp is put in variable,“walltime”. The predicate “ev \_int \_p \_no \_tent” is called, line 1235. The predicate “tentato \_evento ” is asserted in memory, line 1237. The question is : is the predicate “iv(EventoI, Walltime)” asserted in memory? If no this predicate is asserted, else the question is: is the predicate “attiva \_goal(EventoI)” asserted in memory? If yes this predicate is eliminated from asserted facts, line 1241.

Line 1242, the clause *legato \_en* is found and a list of “EventoN ” is created, ListN. If ListN is not empty there is an update of temporal parameter of clause “eve \_pre \_time ” that unifies with the present event. In line 1245 it is recovered the time in which the present event “EventoN” occurred. Line 1247 ,  $T_1$  is the instant in which the present event occurred. In line 1249 the present event is eliminated from the asserted facts. Given a present event “EventoN”, is the predicate “eve \_pre \_time (EventoN, \_\_, \_\_)” asserted in memory? If yes, all predicates “eve \_pre \_time (EventoN, \_\_, \_\_)” are

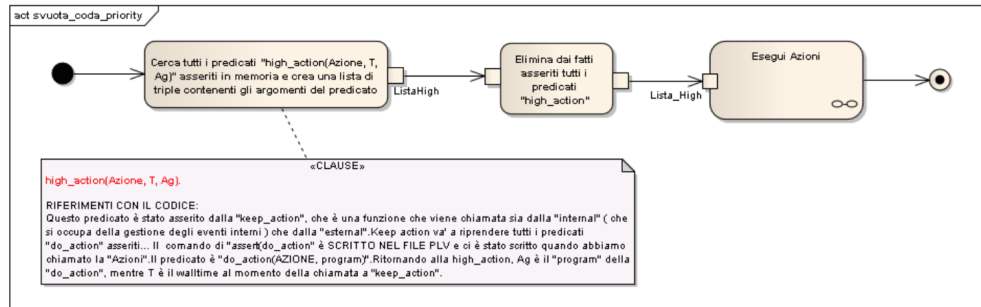
eliminated from the asserted facts, line 1255. Else the predicates “eve \_pre \_time (EventoN, Walltime , ListaTempoPres )” are asserted in memory, line 1256.

### B.1.5. Update the predicates “ tentato \_evento” if the event can be retried



This diagram is a sub-activity of the activity “**Process internal events**” and it is represented in line 1312 by the predicate “ **controlla \_freq \_tent**”. The current walltime is put in a list. The second line of ple-file is read and the content is put in a list. In line 1320 the event “EventoI” unifies with the predicate “internal \_events(EventoI, Frequenza,\_, \_,\_)”. Given an internal event “EventoI”, is the predicate “tentato \_evento(EventoI,\_)” asserted in memory? If yes, in line 1324 the internal event “EventoI” unifies with the predicate asserted in memory “tentato \_evento(EventoI,  $T_c$ )” and  $T_c$  is put in a variable. In line 1326 the variable “Frequenza” is converted in milliseconds and this value is put in a variable “Frequenza1”. In line 1327 it is executed  $ST = Frequenza1 + T_c$ . If  $T > ST$  then the predicate “tentato \_evento(EventoI,  $T_c$ )” that unifies with “EventoI” and “  $T_c$ ” is eliminated from the asserted facts, line 1329.

### B.1.6. The actions with priority “High” are deleted from priorities queue



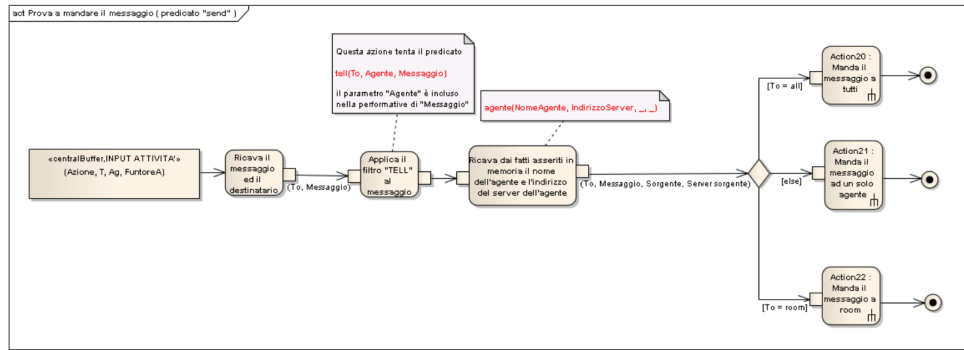
This diagram is a sub-activity of the activity “**Process internal events**” and it is represented in line 925 by the predicate “ **svuota \_coda \_priority**”. The predicate “**clause(high \_action(Azione, T, Ag))**” searches all predicates “**high \_action**” asserted in memory and creates a list of triplet containing the arguments of predicate. In line 934, predicate “ **svuota \_coda \_priority1**” is called: all predicates “ **high \_action**” are eliminated from asserted facts. In line 935 the actions are executed (B.1.6.1.).

#### B.1.6.1. Execute Actions



that the action is complete.

#### B.1.6.1.1. Try to send a message(predicate “ send ”)



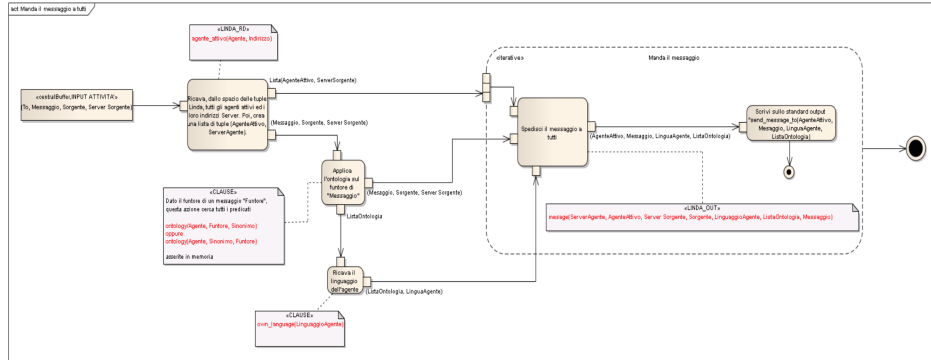
This diagram is a sub-activity of the activity “Execute Actions” and it is represented in line 974 by the predicate “check \_msg”. In line 974 the receiver(To) and the message(M) are obtained. The file agent.plv is opened and the Tell filter is applied to the message “send(To, send \_message(X,Ag)):-tell(To, Ag, send \_message(X)), send m(To, send \_message(X, Ag))”. The predicate “send \_m ” in line 984. The name of the agent and the address of the agent’s server are obtained by the facts asserted in memory. In line 985 there are checks on the type of receiver.

If the receiver is equal to all the message is sent to all(B.1.6.1.1.1.).

Else the message is sent to one agent(B.1.6.1.1.2.).

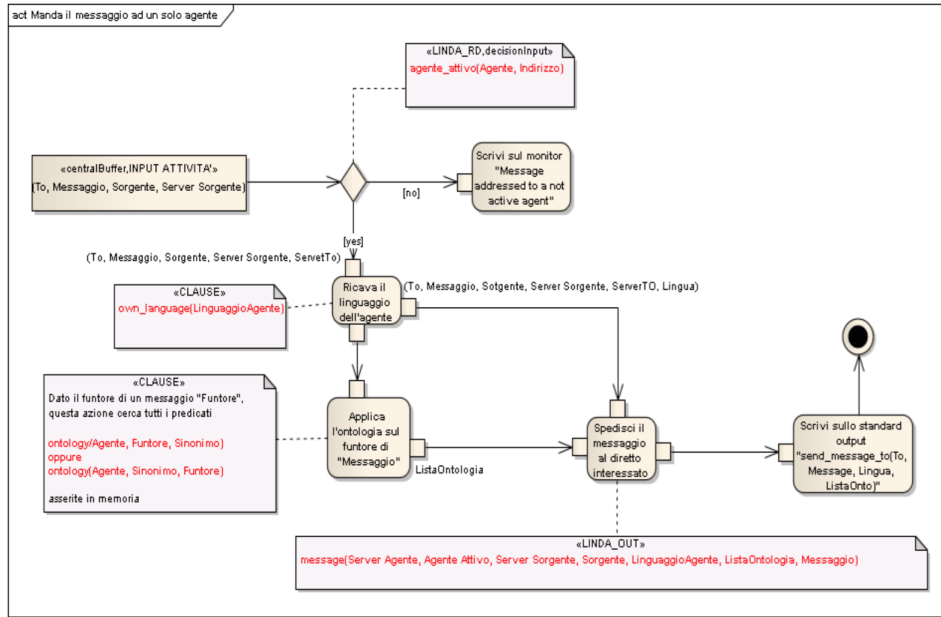
Finally if the receiver is equal to room the message is sent to room.

**B.1.6.1.1.1. The message is sent to all**



This diagram is a sub-activity of the activity “**Try to send a message(predicate “ send ”)**” and it is represented in line 1005 by the predicate “`manda _a _all(M, S, IndS)`”, where M is equal to message, S is the name of the agent and IndS is the address of the Server. From line 1005 to 1008, all the active agents and their server addresses are obtained by the space of the Linda’s tuple. After a list of tuples(`activeAgent`, `ServerAgent`) is created. In line 998 the predicate “`invia _terms _ontology(L)`” is called: given a functor of message, all predicates “ `ontology(Agent, Functor, Synonymous)`” or “ `ontology(Agent, Synonymous, Functor)`”, asserted in memory, are searched. In line 1010, the predicate “`own language( AgentLanguage)`” are called: the language of the agent is obtained. In line 1017, the predicate “`out(message(IndTo, To, IndS, S, Lang, L, M))`” is called (open the file `client.pl`) and the message is sent to all. In line 1018 it is printed the message “`send _message _to(To, M, Lang, L)`”.

#### B.1.6.1.1.2. The message is sent to one agent

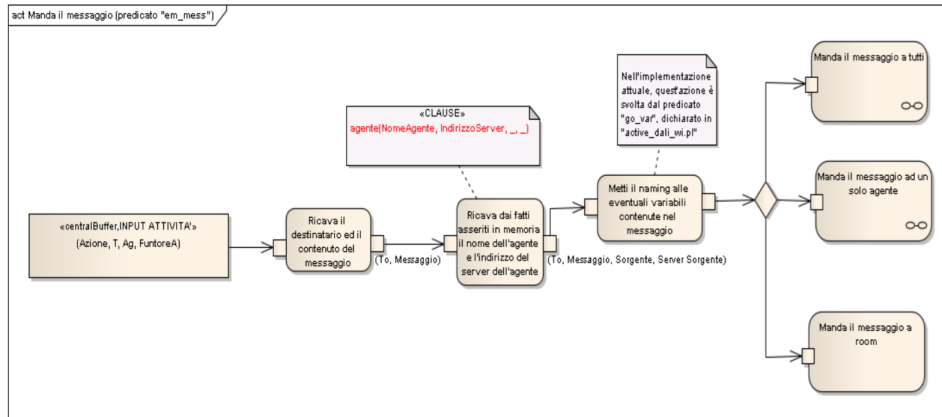


This diagram is a sub-activity of the activity **“Try to send a message(predicate “ send ”)”** and it is represented in line 987 by the predicate “`manda _a(To,M, S, IndS)`”, where To is the receiver M is equal to message, S is the name of the agent and IndS is the address of the Server. If the agent is active, the predicate “`rd _noblock(agente _attivo(To,IndTo))`” is called. If the result fail it is printed: “Message addressed to a not active agent”. Else the predicate “`emetti _ms`” is called in line 994: the language of the agent is obtained. The ontology is applied on the functor “message”. The message is sent to the interested agent(open the file client.pl, line 996). The predicate “`save _on _log _file`” is called in line 1993.

If the receiver is equal to room it is executed line from 1022 to 1027.

#### B.1.6.1.2. Send the message(predicate “ em \_mess ”)





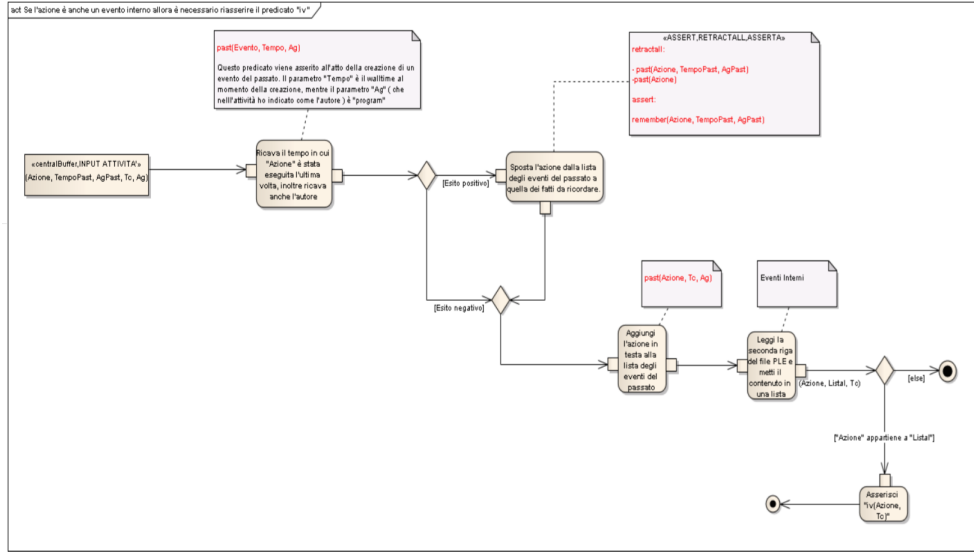
This diagram is a sub-activity of the activity **“Execute Actions”** and it is represented in line 977 by the predicate “em \_mess”. The content of message and the receiver of the message is obtained and the predicate “send \_m0” is called in line 982. The name of the agent and the address of the server are obtained by the facts asserted in memory. The predicate “go \_var” is called (file meta1.pl, line 104): the NAMING is put into the variables in the message. The predicate “send \_m” is called in line 984. In line 985 there are checks on the type of receiver.

If the receiver is equal to all the message is sent to all(B.1.6.1.1.1.).

Else the message is sent to one agent(B.1.6.1.1.2.).

Finally if the receiver is equal to room the message is sent to room.

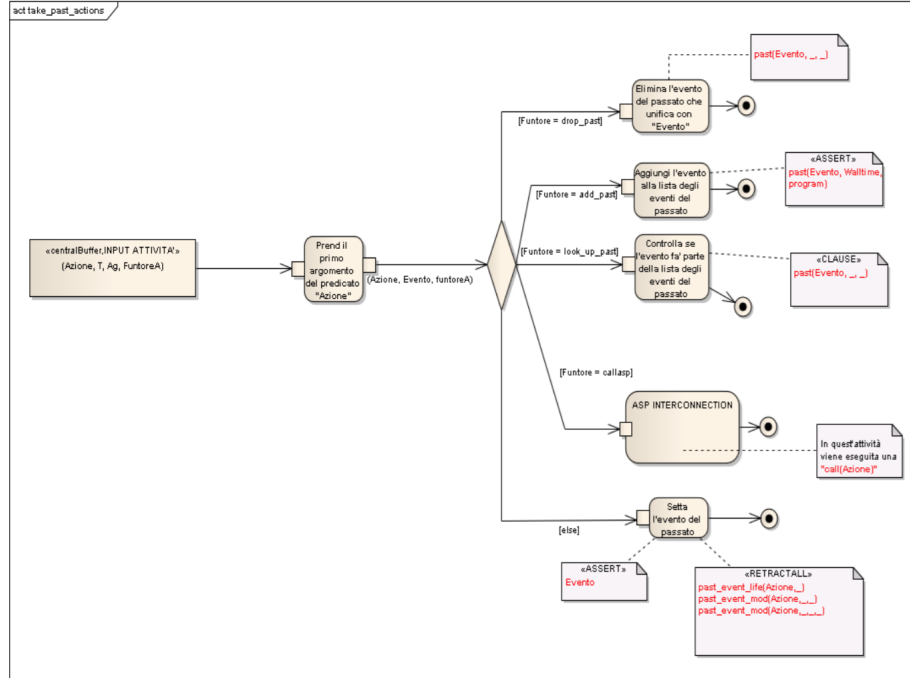
**B.1.6.1.3. If the action is the internal event also then it is necessary reasserted the predicate “iv”**



This diagram is a sub-activity of the activity “**Execute Actions**” and it is represented in line 976-1030 by the predicate “ver \_az \_int”. The predicate “divP” is called in line 1358. The predicate “past(Event,  $T_r$ , Ag)” is asserted when the past event is created. The parameter  $T_r$  is the walltime at the moment of creations, while A is program. It is obtained the time in which “Azione” was last performed and also derived the author. If the clause has the positive result the predicate “sposta \_in \_rem” is called in line 1360: the actions is moved from the list of the past events to the list of facts to remember. If the clause has the negative result the predicate “asserta(past(Me, $T_r$ ,  $A_g$ ))” is called in line 1358: the action is added in the head of the past events.

The predicate “ver \_az \_int” is called in line 1030. The second line in ple-file is called and the content is put in a list, L. If L is not empty the predicate “ver \_az \_int1” is called in line 1035 and the “iv(X,  $T_c$ )” is asserted, where X is Action.

#### B.1.6.1.4. Take \_past \_actions



This diagram is a sub-activity of the activity **“Execute Actions”** and it is represented in line 955-957 by the predicate “Take \_past \_actions”. The first argument( $X=Action$ ) is taken and the predicate “caso0 \_past( $F,E,X$ )” is called in line 958.

If the functor  $F$  is equal to drop \_past the predicate “drop \_evento( $X$ )” is called in line 967: the past event that unified with “Evento” is eliminated ( $retractall(past(X,_,_))$ ).

If the functor  $F$  is not equal to drop \_past the predicate “caso1 \_past( $F,E,X$ )” is called in line 959.

If the functor  $F$  is equal to add \_evento the predicate “add \_evento( $X$ )” is called in line 968: the event is added to the list of the past event.

If the functor  $F$  is not equal to add \_evento the predicate “caso2 \_past( $F,E,X$ )” is called in line 960.

If the functor  $F$  is equal to look \_up \_past the predicate “look \_up \_evento( $X$ )” is called in line 969: it is checked if the events belong to the list of past events.

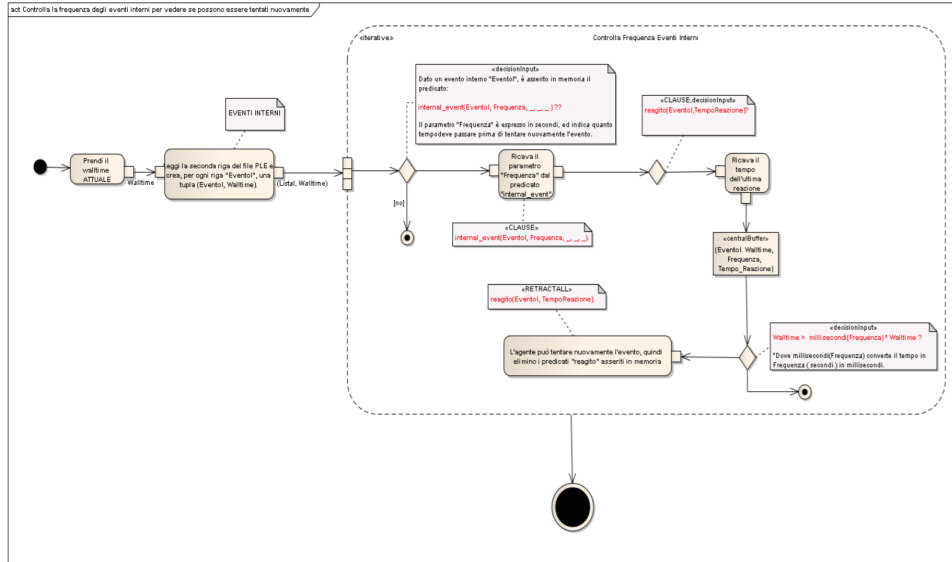
If the functor  $F$  is not equal to look \_up \_past the predicate “caso3 \_past( $F,E,X$ )”

is called in line 963.

If the functor F is equal to callaps the predicate “callasp” is called in line 1747: ASP-INTERCONNECTION. The predicate “call(Y)”, where Y is Action , is called.

Else the predicate “set \_past \_evento” is called in line 970. The past event is set: retractall of the predicates “past \_event \_life(X,\_)”, “past \_event \_mod(X, \_, \_)”, “past \_event \_mod (X, \_, \_, \_)”, where is Action.

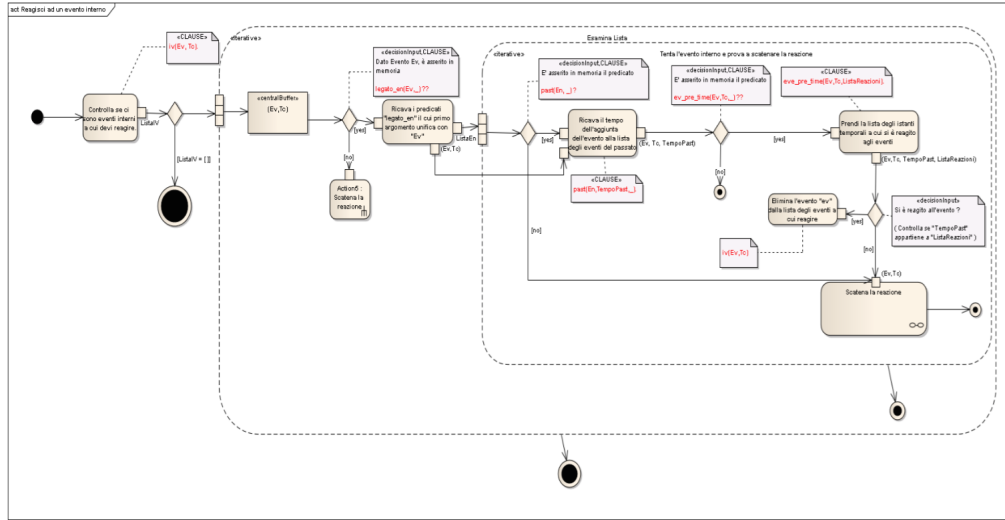
### B.1.7. Check the frequency of the internal events to see if they can be retried



This diagram is a sub-activity of the activity “**Process internal events**” and it is represented in line 1331 by the predicate “ **controlla \_freq \_iv**”. The current walltime is taken in line 1331. The second line of ple-file is read and for each line “EventoI” is created a tuple(EventoI, Walltime). If L is not empty the predicate “ **contr \_freq \_iv0** ” is called in line 1336. Given an internal event “EventoI”, the question is: is the predicate “**internal \_event (EventoI, Frequency, \_,\_,\_)**” asserted in memory? The Frequency parameter is expressed in seconds and indicates how much time must pass before attempting the event again. If yes, the parameter Frequency is ob-

tained from the “clause(internal \_event(X, Tp, \_, \_, \_))”, line 1339. The predicate “controlla \_freq \_iv1” is called in line 1340. If “clause(reagito(X, Tc))” is true the predicate “controlla \_freq \_int2” is called in line 1349: the time is transformed into milliseconds and  $ST = Tp1 + Tc$ . If  $T > ST$  then the predicate “canc \_e \_reagito” is called in line 1352: the agent can try the event again the the predicates “reagito”, asserted in memory, are eliminated .

### B.1.8. React to the Internal Event



This diagram is a sub-activity of the activity “**Process internal events**” and it is represented in line 1267 by the predicate “ **scatena** ”. It is checked if there are Internal events to react to and the predicate “ manage list int(L) ”, where L is a list of internal events, is called in line 1268. If L is not empty the predicate “ scat0 ” is called in line 1275. Given an event Ev, is the predicate “ legato \_en(Ev, \_) ” asserted in memory? If no the predicate “ scat1 ” is called in line 1290(B.1.8.1.).

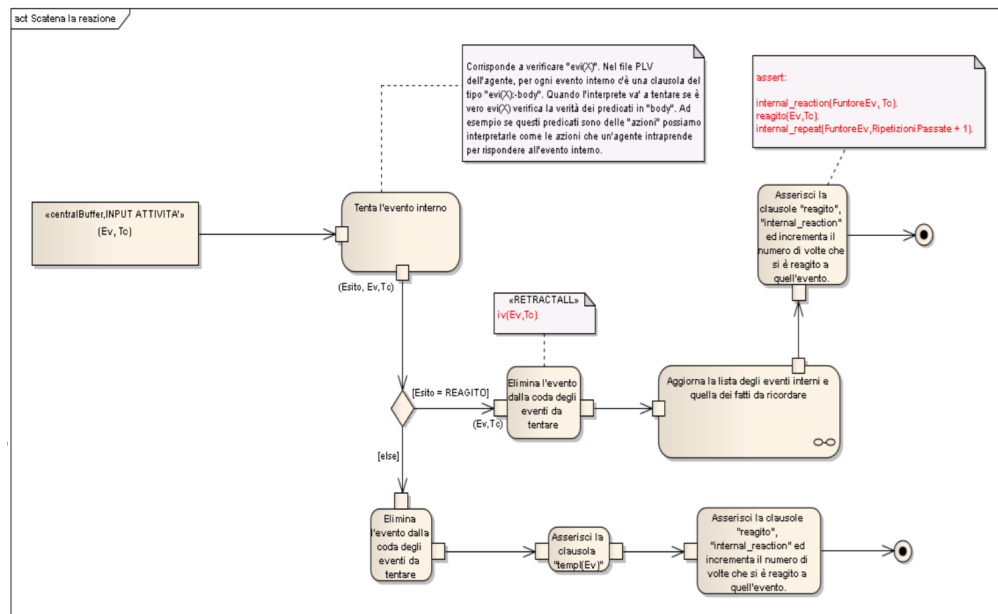
If yes the predicate “ scat0 \_en ” is called in line 1277: it is derived the predicates “ legato \_en ” whose first argument unifies with Ev. The question is: is the predicate “ past(X, T1, \_) ” asserted in memory? If no it is called line 1290(B.1.8.1.).

If yes it is obtained the time of adding the event to the list of past events and the predicate “ check \_app ” is called in line 1284. The question is : is the predicate “ eve \_pre \_time ” is asserted in memory? If yes the list of time instants to which it has reacted is obtained in line 1284. The predicate “ check \_member ” is called in line 1284.

If no the predicate “ scat1 ” is called in line 1290(B.1.8.1.).

If yes the predicate “ritrai \_ev \_pres” is called in line 1288: the event is eliminated from the list of internal event to which it has reacted.

### B.1.8.1. Trigger the reaction

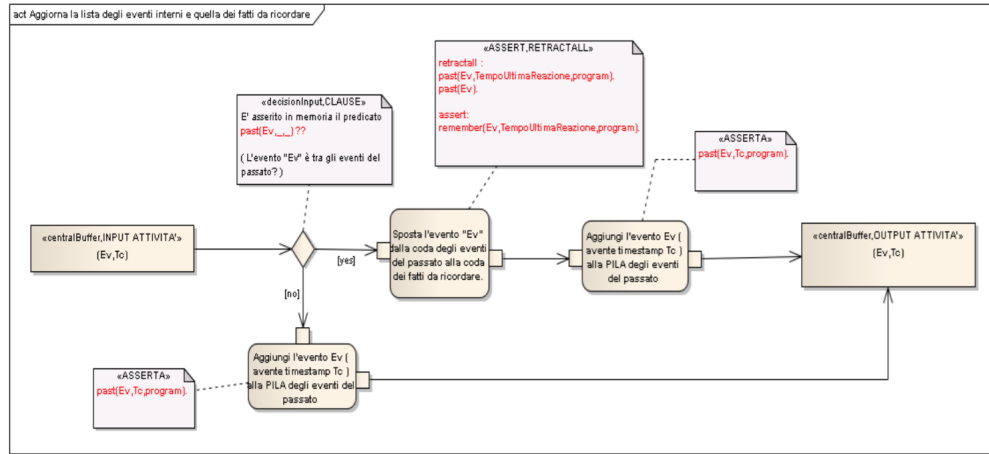


This diagram is a sub-activity of the activity “**React to the Internal Event**” and it is represented in line 1290 by the predicate “ scat1 ”. The predicate “evi(X)” is verified and then the internal event is triggered. If the result is equal to REAGITO, the predicate “scatena \_reagisci” is called in line 1296: the event is eliminated from the queue of the events to try. The predicate “divP” is called in line 1296: the list of the internal events and the list of the facts to remember are updated (B.1.8.1.1.). The clause “reagito” and

“internal reaction” are asserted and the number of times it has reacted, is increased line 1297-1299.

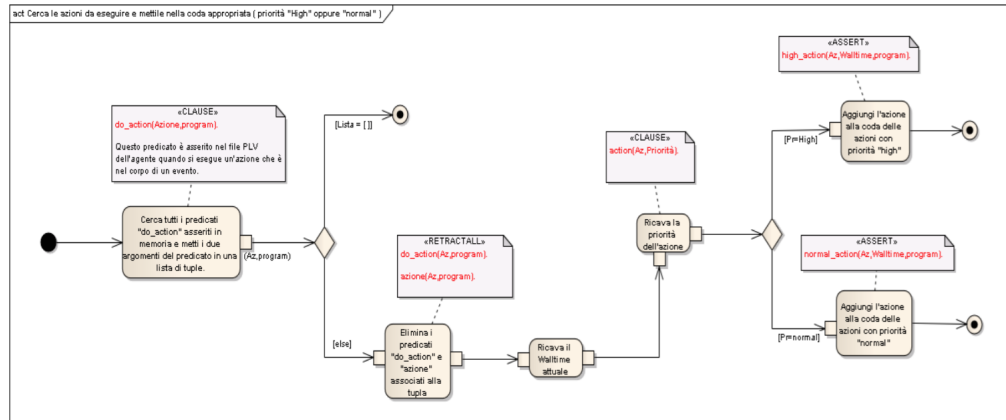
If the result is not equal to REAGITO, the predicate “no \_scatena” is called in line 1291: the event is eliminated from the queue of the events to try. The clause “reagito” and “internal reaction” are asserted and the number of times it has reacted, is increased line 1293.

#### B.1.8.1.1 The list of the internal events and the list of the facts to remember are updated



This diagram is a sub-activity of the activity “**Trigger the reaction** ” and it is represented in line 1296 by the predicate “ **scatena \_agisci** ”. The question is: is the predicate “ **past(Ev, \_, \_)** ” in line 1358 asserted in memory? If yes the predicate “**sposta \_in \_rem**” is called in line 1360: the event Ev is moved from queue of past events to the queue of the facts to remember and the event Ev is added to the stack of the past events. If no the predicate “**asserta(past(Ev, Tc,program))**” is called in line 1358: the event Ev is added to the stack of the past events.

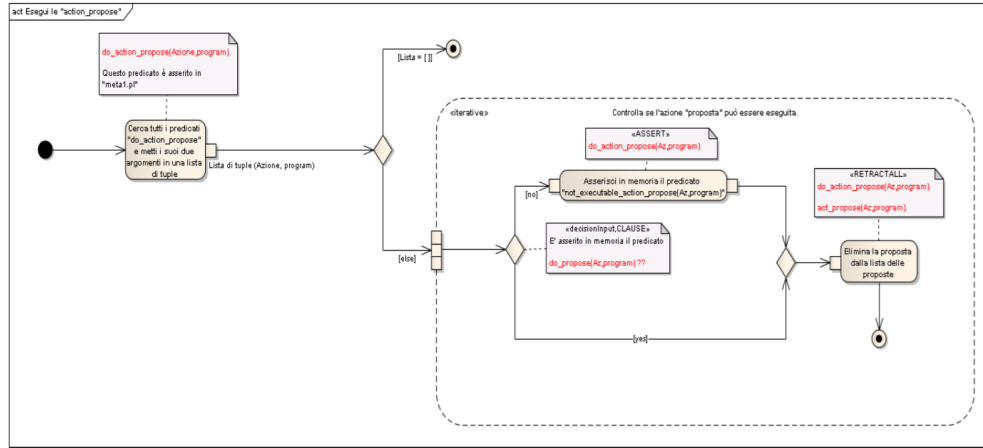
#### B.1.9. Search for the actions to be performed and place them in the appropriate queue



This diagram is a sub-activity of the activity “**Process internal events**” and it is represented in line 906 by the predicate “ **keep \_action** ”. All predicates “do \_action ” asserted in memory, are found and the two arguments of the predicate are put in a list of tuples(line 906). If the list is not empty the predicate “keep \_action1” is called in line 908: the predicates “do \_action” and “action” associated with the tuples, are eliminated. The current Walltime is obtained in line 918. The predicate “clause(action(Az,Mod), \_)” is called in line 919: it is obtained the priority of the actions(Mod). If Mod is equal to high then the action is added to the queue of the actions with priority high, line 920. If Mod is equal to normal then the action is added to the queue of the actions with priority high, line 920.



### B.1.10. Execute the action `_propose`

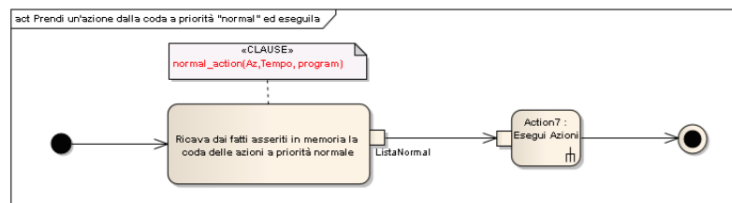


This diagram is a sub-activity of the activity “**Process internal events**” and it is represented in line 711 by the predicate “**execute \_do \_action \_propose**”. All predicates “do \_action \_propose” are found and the two arguments of the predicate are put in a list of tuples (line 906). This predicate is asserted in file meta1.pl. Is the predicate “do \_propose(Az, Ag)” asserted in memory?

If yes the retractall of “do \_action \_propose(Az, Ag)” and “act \_propose(Az, Ag)” are executed, line 722-723.

If no it is asserted in memory “not \_executable \_action \_propose”, line 721.

### B.1.11. Take an action from the queue with “normal” priority

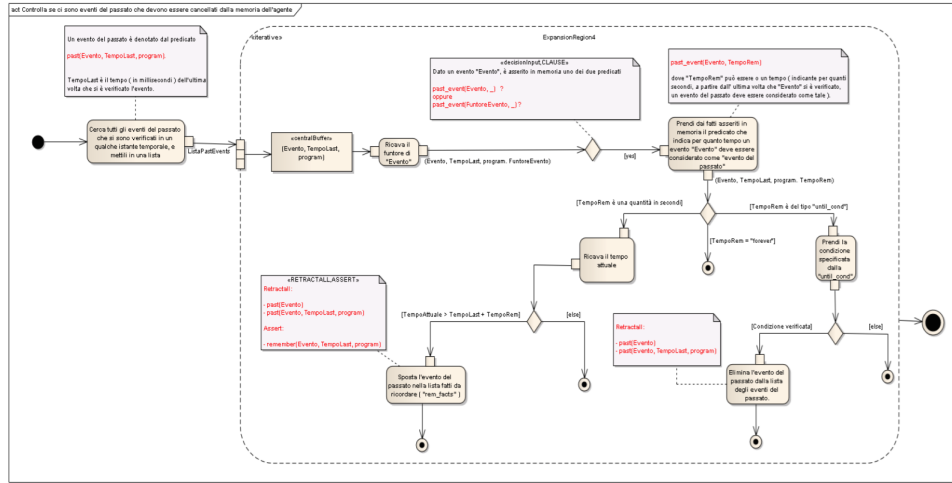


This diagram is a sub-activity of the activity “**Process internal events**” and it is represented in line 940 by the predicate “**prendi \_action \_normal**”. The queue of actions of normal priority is obtained from the facts asserted in memory. And subsequently the predicate “fa” is called in line 947 (see subsection B.1.6.1.).

### B.1.12. Predicate controlla \_vita

The predicate “ **controlla \_vita**” is described in line 1515 and is a sub-activity of the activity “**Process internal events**”. This predicate is composed by two predicate: “ **controlla \_vita \_past \_base**” (B.1.12.1.) and “ **controlla \_vita \_remember**” (B.1.12.2.).

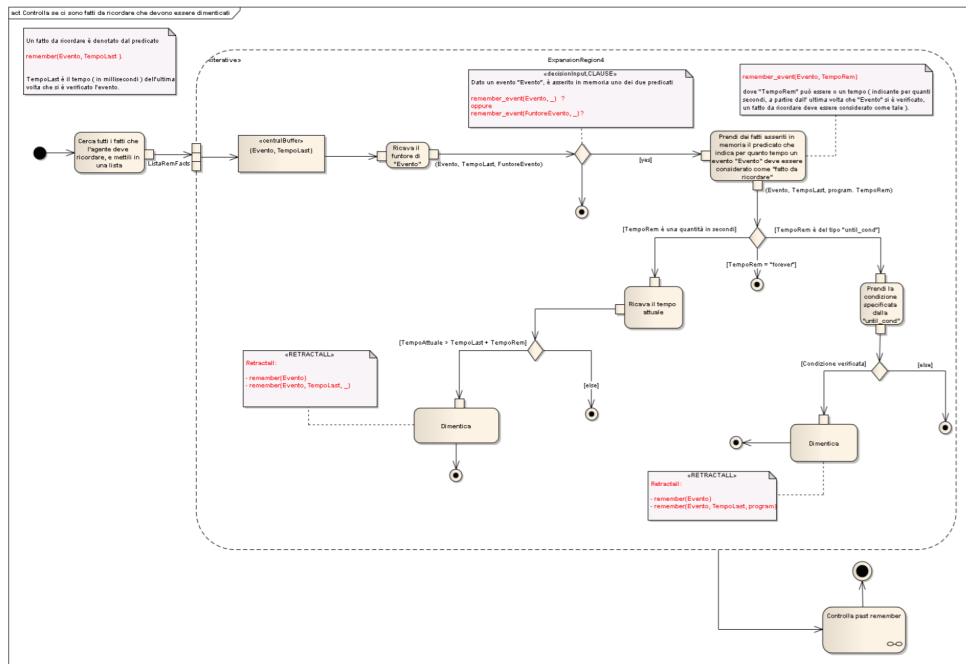
**B.1.12.1. Check if there are past events which be deleted from agent’s memory.**



This diagram is represented in line 1366 by the predicate “ **controlla \_vita \_past \_base**”. All the past events that have occurred at same point in time are searched for and put them in a list. The predicate “ **controlla \_vita \_past** ” is called in line 1368 and the functor of event is obtained in line 1371. The question is: given an event is one of the two predicates, “ **past \_event(Event, \_)**” or “ **past \_event(FuncionEvent, \_)**”, asserted in memory? If yes the predicate “ **contrr \_past \_event1**” is called in line 1376 and the predicate “**controlla \_vita \_past1**” is called in line 1377. If the functor is equal to forever then end, else the predicate “**controlla \_vita \_past2**” is called in line 1378: if the functor is equal to until the predicate “**processa \_cd \_past**” is called in line 1390. The past event is eliminated from the list of past events, line 1394. Else the predicate “**processa \_tempo \_past**” is called

in line 1380: the past event is moved to the list of facts to remember (“rem  
\_fact”).

**B.1.12.2.** Check there are any facts to remember that must be forgotten

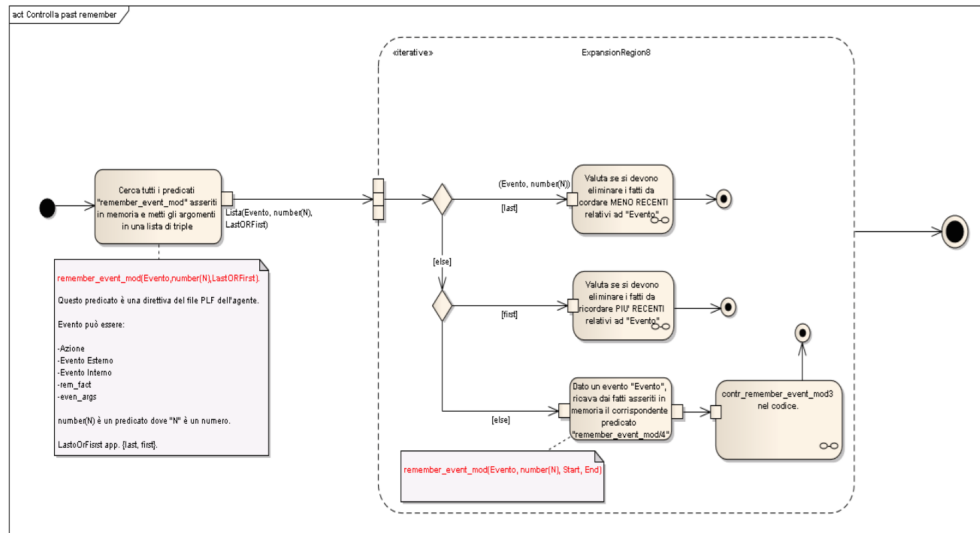


This diagram is represented in line 1400 by the predicate “**controlla \_vita \_remember**”. The predicate “controlla \_vita \_remember \_base” is called in line 1402: all the facts that the agent must be remember are searched for and put them in a list. The functor of “Event” is obtained in line 1410. The question is: given an event “Event” is one of the two predicates, “remember \_event(Event, \_)” or “remember \_event(FunctorEvent, \_)”, asserted in memory? If yes, the predicate indicating how long an event “Event” must be considered to “fact to remember”, is taken from the facts asserted in memory. The predicate “contrr \_remember \_event1” is called in line 1411-1412 and the predicate “controlla \_vita \_remember1” is called in line 1416. If the functor is equal to forever the end else the predicate “controlla \_vita \_remember2” is called in line 1417. If the functor is equal to until the predicate

“processa \_cd \_remember” is called in line 1427(retractall remember(Event), retractall remember(Event, TempLast, \_)). The predicate “processa \_tempo \_rem” is called in line 1419. If the current time is greater than of the sum between TempoLast plus TempoRem then the predicate “processa \_tempo \_canc item \_rem” is called in line 1424.

The predicate “controlla \_past \_remember” is called in line 1435(B.1.12.2.1.).

#### B.1.12.2.1. Controlla \_past \_remember



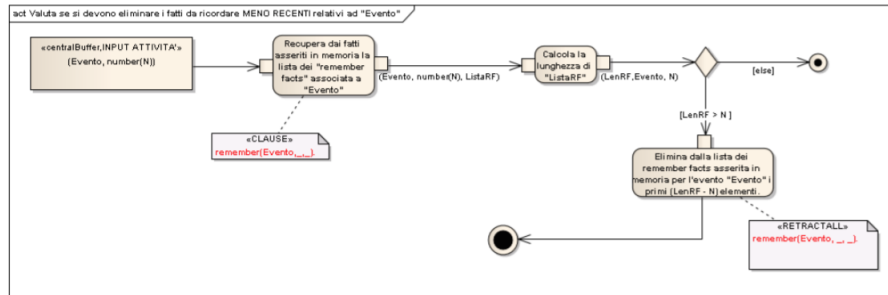
This diagram is a sub-activity of the activity “**Check there are any facts to remember that must be forgotten**” and it is represented in line 1435 by the predicate “**controlla \_past \_remember**”. All predicates “remember \_event \_mod” asserted in memory are searched and put the arguments in a list of tuples. This predicates is a directive of plf-file of the agent. If the list is not empty then the predicate “controlla \_past \_mod \_rem0” is called in line 1442.

The predicate “clause(remember event mod(Me, number(N), last), )” is true then the predicate “contrr \_remember \_event \_mod1” is called in line 1451: it is evaluated if you need to eliminate the oldest events to remember about “Event”(B.1.12.2.1.1.).

The predicate “clause(remember \_event \_mod(Me, number(N), first), \_)” is true then the predicate “contrr \_remember \_event \_mod2” is called in line 1469: it is evaluated if you need to eliminate the younger events to remember about “Event”(B.1.12.2.1.2.).

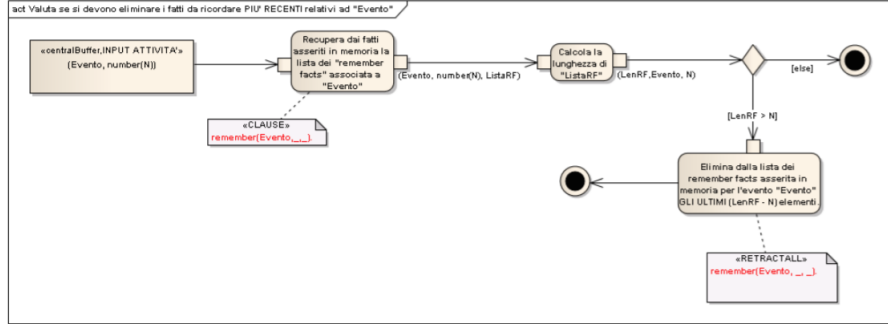
Else the predicate “clause(remember event mod(Me, number(N), T1, T2), )” is true then given an event “Event”, it obtains from the facts asserted in memory the corresponding predicate “remember event mod/4”. The predicate “contrr \_remember \_event \_mod3” is called in line 1487.

#### B.1.12.2.1.1. Evaluate if you need to delete the oldest events to remember about “Event”



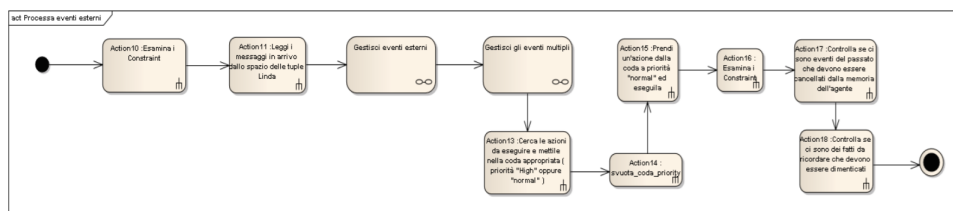
This diagram is a sub-activity of the activity “**Controlla \_past \_remember**” and it is represented in line 1451 by the predicate “**contrr \_remember \_event \_mod1**”. The list of the “remember facts” associated to “Event” is obtained from facts asserted in memory (clause (remember(P,T,S),\_), line 1451). The length of ListRF (List of remember facts) is calculated, line 1452. If the length of ListRF is greater than N then the predicate “eliminate \_not \_last \_items \_remember” is called in line 1456: it is eliminated from the remember-facts’ list, asserted in memory for the event “Event”, the first LenRF - N elements.

#### B.1.12.2.1.2. Evaluate if you need to delete the younger events to remember about “Event”



This diagram is a sub-activity of the activity “**Controlla \_past \_remember**” and it is represented in line 1469 by the predicate “**contrr \_remember \_event \_mod2**”. The list of the “remember facts” associated to “Event” is obtained from facts asserted in memory (clause (past(P,T,S),\_), line 1469 ). The length of ListRF (List of remember facts) is calculated, line 1470. If the length of ListRF is greater than N then the predicate “eliminate \_not \_first \_items \_remember” is called in line 1471: it is eliminated from the remember-facts’ list, asserted in memory for the event “Event”, the last LenRF - N elements.

## B.2. Process External Events



This diagram is represented by the predicate “ **external** ” in line 1505.

```

internal:- blocco_constr, ricmess, processa_eve,
examine_mul, keep_action, svuota_coda_priority,
prendi_action_normal, blocco_constr, controlla_vita.
  
```

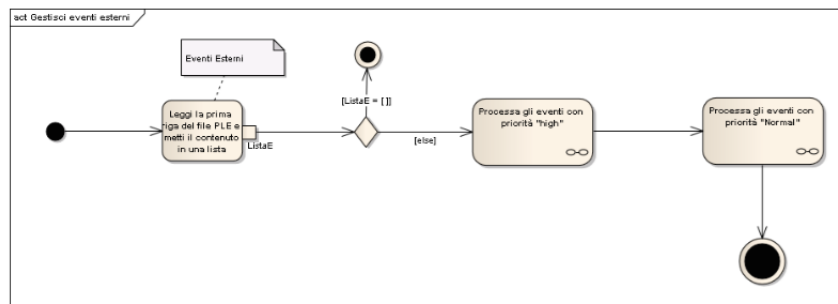
### B.2.1. Predicate blocco \_constr

The predicate “**blocco \_constr**” is the same predicate defined in section B.1.1.

### B.2.2. Predicate ricmess

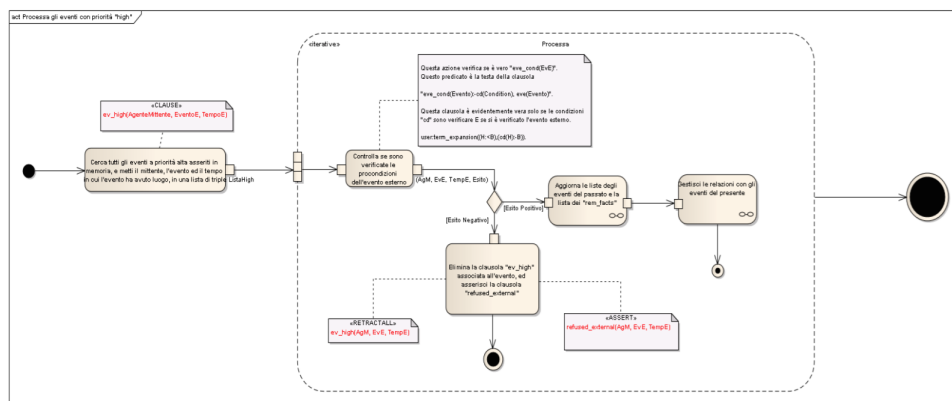
The predicate “**ricmess**” is the same predicate defined in section B.1.2.

### B.2.3. Manage the External events



This diagram is a sub-activity of the activity “**Process External Events**” and it is represented in line 746 by the predicate “**processa \_eve**” : the first line in ple-file is read and it is put them in a list, Es. If Es is not empty the predicate “processa \_eve \_high”(B.2.1.1.) and “processa \_eve \_normal” (B.2.1.2.) are executed .

#### B.2.3.1. Process the events with high priority

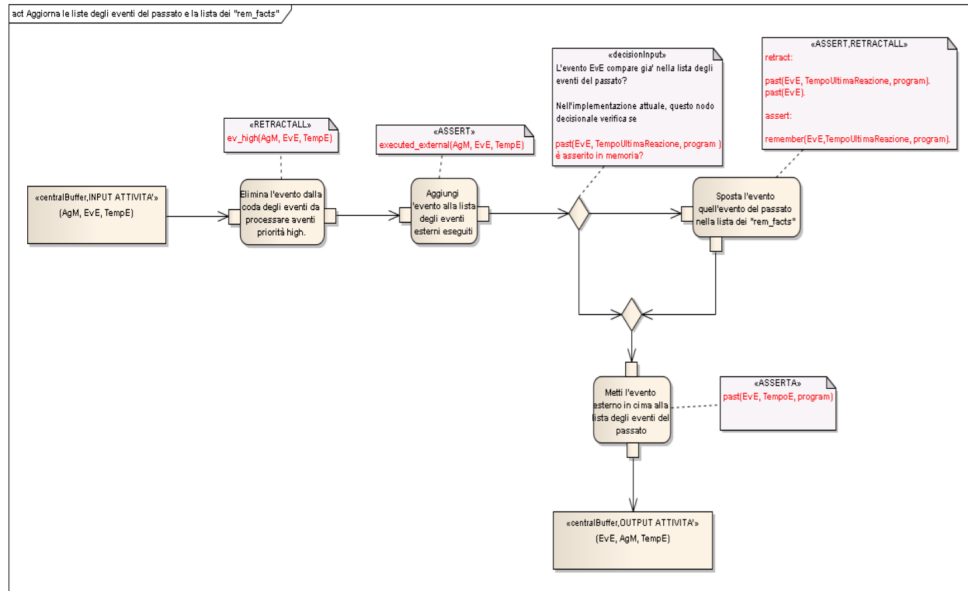


This diagram is a sub-activity of the activity “**Manage External events**” and it is represented in line 749 by the predicate “**processa \_eve \_high**” . All events with high priority asserted in memory are searched and it puts the event, the sender, and the time in which it took place, in a list of tuples, ListHigh. The predicate “processa \_eve \_high1” is called in line 750.

In line 755 it checked if the condition of external event is verified: this action verifies if the condition “eve \_cond(EVE)” is true (this predicate is in plv-file). If the result is negative the predicate “no \_proc \_eve \_high” is called in line 757: it asserts the clause “refused \_external”.

If the result is positive the predicate “processa \_eve \_high/3” is called in line 758: it updates the list of past events and the list of “rem \_facts” (B.2.3.1.1.) and it manages the relationship with the present events(B.2.3.1.2.).

#### B.2.3.1.1. Updated the list of past events and the list of “rem \_facts”

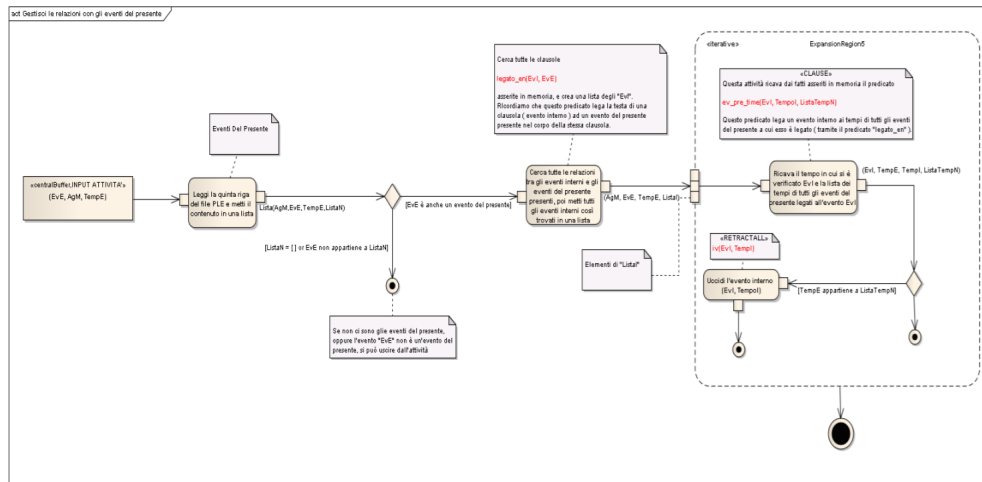


This diagram is a sub-activity of the activity “**Process the events with high priority**” and it is represented in line 758 by the predicate “**processa**



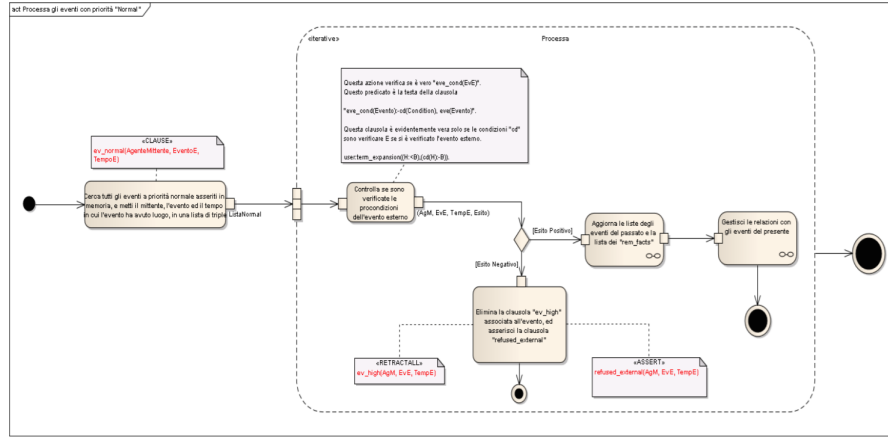
“\_eve\_high/3”. The event is eliminated from the queue of the event with priority high to be processed, line 758. The event is added to the list of external event executed, line 758. The predicate “divP” is called in line 1358. The question is: is the event already on the list? If yes the predicate “sposta\_in\_rem” is called in line 1360: it moves this past event from the list of “rem\_facts”. If no the external event is put on the top of the past event’s list, line 1358.

#### B.2.3.1.2. Manage relationship with the present events



This diagram is a sub-activity of the activity “**Process the events with high priority**” and it is represented in line 758 by the predicate “**processa \_eve\_high/3**”. The fifth line of ple-file is read and it puts the content in a list. If this list is empty or the event doesn’t belong to the list, the activity ends. If the event is also a present event the predicate “processa \_eve3” is called in line 781. It searched all relations between the internal events and the present events. It puts all internal events in a list. If the “clause(ev\_pre\_time)” is true in line 783, the predicate “processa \_eve4” is called in line 784: the internal event is killed(uccidi \_iv/2), line 786.

### B.2.3.2. Process the events with normal priority

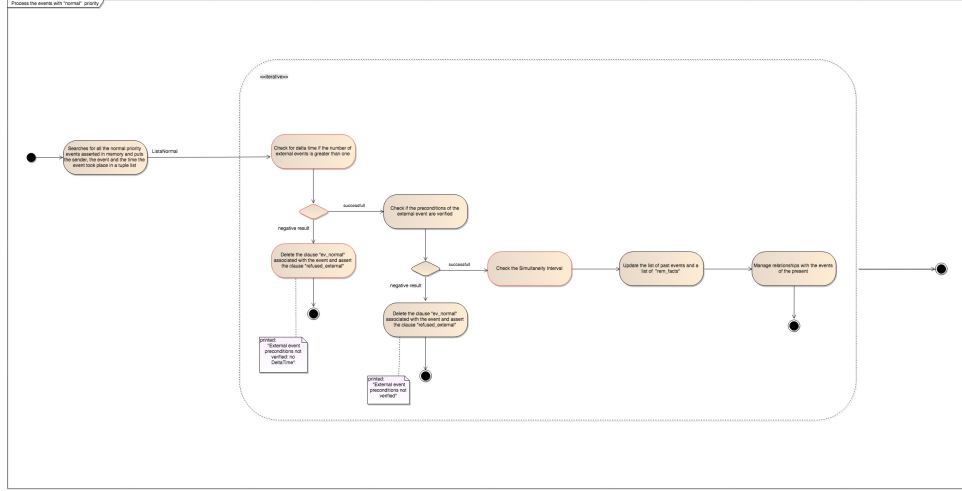


This diagram is a sub-activity of the activity “**Manage External events** ” and it is represented in line 763 by the predicate “ **processa \_eve \_normal** ” . All events with normal priority asserted in memory are searched and it puts the event, the sender, and the time in which it took place, in a list of tuples, ListNormal. The predicate “processa \_eve \_normal1” is called in line 763.

In line 766 it checked if the condition of external event is verified: this action verifies if the condition “eve \_cond(EVE)” is true (this predicate is in plv file). If the result is negative the predicate “no \_proc \_eve \_normal” is called in line 768: it asserts the clause “refused \_external”.

If the result is positive the predicate “processa \_eve \_normal/3” is called in line 775: it updates the list of past events and the list of “rem \_facts” and it manages the relationship with the present events.

### B.2.3.2.(Updated) Process the events with normal priority



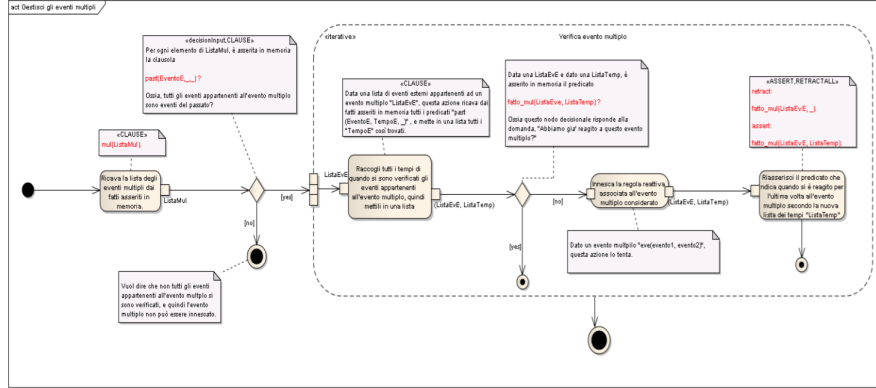
This diagram is a sub-activity of the activity “**Manage External events**” and it is represented in line 763 by the predicate “**processa \_eve \_normal**” . All events with normal priority asserted in memory are searched and it puts the event, the sender, and the time in which it took place, in a list of tuples, ListNormal. The predicate “processa \_eve \_normal1” is called in line 764.

In line 765 the predicate “*check \_presence*” is called. If the result is false then the predicate “no \_proc \_eve \_normal \_no \_time” is called in line 771. Else in line 766 it checked if the condition of external event is verified: this action verifies if the condition “eve \_cond(EVE)” is true (this predicate is in plv file).

If the result is negative the predicate “no \_proc \_eve \_normal” is called in line 768: it asserts the clause “refused \_external”.

If the result is positive the predicate “processa \_eve \_normal/3” is called in line 775: the predicate “**simultaneity \_interval**” is called in line 2167 and then it updates the list of past events and the list of “rem \_facts” and it manages the relationship with the present events.

## B.2.4. Manage multiple events



This diagram is a sub-activity of the activity **“Process External Events”** and it is represented in line 232 by the predicate **“examine \_\_mul”**. The predicate **“examine1 \_\_mul”** is called in line 233: the list of the multiple events is obtained from the facts asserted in memory(`clause(ListaMul)`). The predicate **“keep \_\_past”** is called in line 239: for each element of `ListaMul`, is the clause **“past(EventE, \_\_, \_\_)”** asserted in memory? If yes it is collected all time when the events belonging to the multiple event have occurred. The predicates **“chiama \_\_cong”** in line 245, **“controlla \_\_time \_\_ep/2”** in line 247 and **“cont \_\_tep1/2”** in line 251 are called. Given a list `ListEVE` and a list `ListTemp` is the predicate **“fatto \_\_mul/2”** asserted in memory? That is, did you react to this multiple event? If the answer is no, it triggers the corresponding reactive rule and finally it reasserts the predicate that indicates when the last time you reacted to the multiple event according to the new time list **“ListTemp”**.

## B.2.5. Predicate keep \_\_action

The predicate **“keep \_\_action”** is the same predicate defined in section B.1.9.

## B.2.6. Predicate svuota \_\_coda \_\_priority

The predicate **“svuota \_\_coda \_\_priority”** is the same predicate defined in section B.1.6.

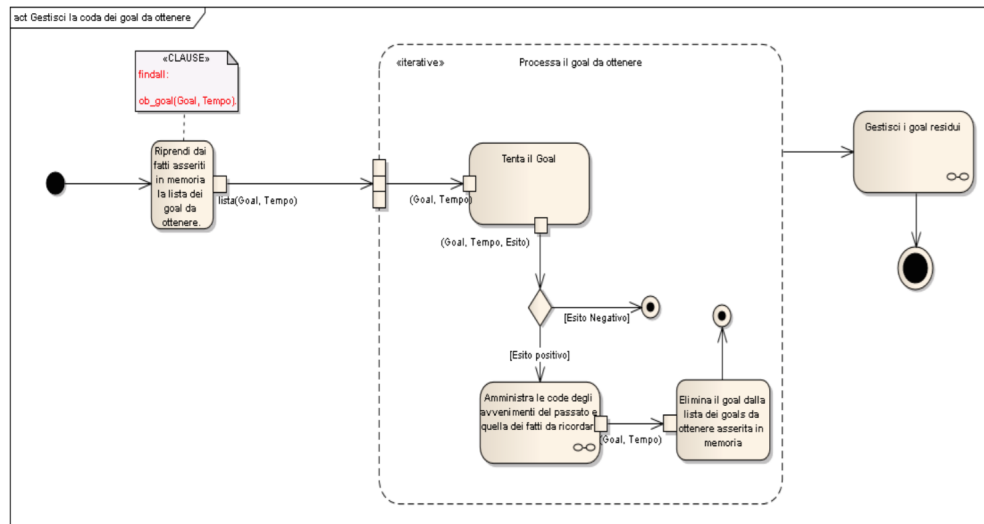
### B.2.7. Predicate `prendi _action _normal`

The predicate “`prendi _action _normal`” is the same predicate defined in section B.1.11.

### B.2.8. Predicate `controlla _vita`

The predicate “`controlla _vita`” is the same predicate defined in section B.1.12.

## B.3. Manage queue of the goals to get



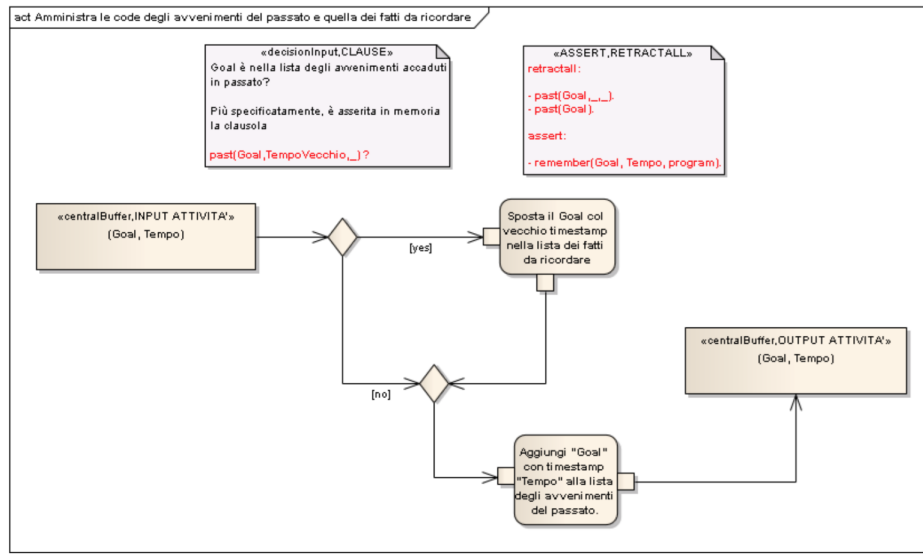
This diagram is represented by the predicate “ `side _goal` ” in line 1507.

```
side_goal:- obtaining_goals , residue_goal.
```

The predicate “ `obtaining _goals` ” is called in line 1259: the list of goal to obtain is recovered from the facts asserted in memory. The predicate “ `process _ob _goal` ” is called in line 1260: the goal is tempted. If the predicate “`once(E1)`” has a positive result the predicate “`fai _div _evp`” is called in line 1261 (B.3.1.) : manage the queue of past events and the facts to

remember. The list of goals asserted in memory is eliminated: “retractall(ob \_goal(El,T))”, line 1261. Finally the predicate “residue \_goal” is called in line 1533(B.3.2.).

### B.3.1. Manage queue of past events and the facts to remember

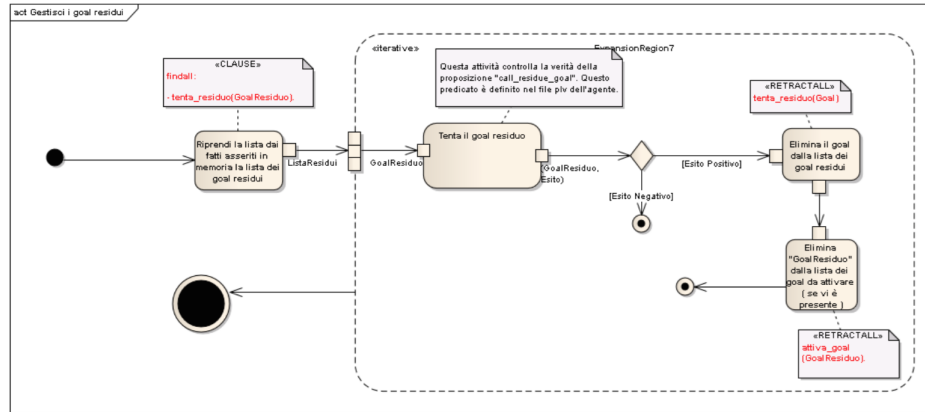


This diagram is a sub-activity of the activity “**Manage the queue of the goals to get**” and it is represented in line 1261 by the predicate “ **fai \_div \_evp** ” . The predicate “divP/3 ” is called in line 1358. Is the clause “ clause(past(Goal,Time,\_)) ” asserted in memory?

If no the predicate “ asserta(past(Goal, Time, \_)) ” is executed in line 1358.

If yes the predicate “sposta \_in \_rem/5 ” is called in line 1360: the goal with timestamp is added to the list of past events.

### B.3.2. Manage residual goals



This diagram is a sub-activity of the activity “**Manage the queue of the goals to get**” and it is represented in line 1533 by the predicate “**residue \_goal**” . The predicate “**residue \_goal1**” is called in line 1534: the list of facts asserted in memory is obtained from the list of residual goals. The truth of the preposition “**call \_residue \_goal**” is checked in line 1538 (this predicate is defined in pl file of the agent).

If the result is positive then the predicate “**resid \_goal**” is called in line 1540: the goal is eliminated from the list of residual goals. The predicate “**canc \_goal/1**” is called in line 1363: the residual goal is eliminated from the list of goals to activate.